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A NON-LINEAR VECTOR FINITE DIFFERENCE SCHEME

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A Non-linear vector finite difference scheme \*)

by

Paul J. Schweitzer \*\*)

ABSTRACT

A non-linear vector iterative scheme of linear fractional form is investigated. Explicit expressions are given for each vector iterate and for the limiting vector. These equations arise in the bottleneck analysis of closed networks of queues.

KEY WORDS & PHRASES: *vector difference equations, queueing networks.*

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\*) This report will be submitted for publication elsewhere.

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## 1. INTRODUCTION

The following finite difference scheme arose in a closed network of queues (see Appendix):

$$(1) \quad y_i(k+1) = a_i \left[ \frac{1}{k} + y_i(k) \right] / \sum_{j=1}^m a_j \left[ \frac{1}{k} + y_j(k) \right] \quad 1 \leq i \leq m; \quad k=1,2,3,\dots$$

$$(2) \quad y_i(1) = a_i / \sum_{j=1}^m a_j \quad 1 \leq i \leq m;$$

where  $m \geq 2$  and every  $a_i > 0$ . For each  $k$ , the  $y$ 's are strictly positive and sum to unity. Show that  $\lim_{k \rightarrow \infty} y_i(k)$  exists for each  $i$  and find its value. Give an expression for  $y_i(k)$  if  $\{a_i\}_{i=1}^m$  are distinct.

## 2. SOLUTION

With the notation

$$b_k = \sum_{i=1}^m (a_i)^k$$

the first few iterates of (1) yield, after clearing fractions, the expressions

$$y_i(1) = a_i / b_1 \quad 1 \leq i \leq m$$

$$y_i(2) = [a_i b_1 + (a_i)^2] / [(b_1)^2 + b_2] \quad 1 \leq i \leq m$$

$$y_i(3) = \frac{\frac{(b_1)^2 + b_2}{2} a_i + b_1 (a_i)^2 + (a_i)^3}{\frac{(b_1)^2 + b_2}{2} b_1 + b_1 b_2 + b_3} \quad 1 \leq i \leq m.$$

This suggests the general form

$$(3) \quad y_i(k) = \frac{\sum_{j=1}^k c(k)_j (a_i)^j}{\sum_{j=1}^k c(k)_j b_j}$$

where the  $c$ 's are independent of  $i$ . Insertion of (3) into (1) yields

$$(4) \quad c(k+1)_1 = \frac{1}{k} \sum_{j=1}^k c(k)_j b_j \quad k \geq 1$$

$$c^{(k+1)}_j = c^{(k)}_{j-1} \quad 2 \leq j \leq k+1, \quad k \geq 1$$

which implies

$$(5) \quad c^{(k)}_j = d_{k-j} \quad 1 \leq j \leq k < \infty$$

and the recursion formula

$$(6) \quad d_k = \frac{1}{k} \sum_{j=1}^k d_{k-j} b_j \quad k = 1, 2, 3, \dots$$

with initial condition

$$(7) \quad d_0 = c^{(1)}_1 = 1.$$

We may solve for  $\{d_k\}_{k=0}^{\infty}$  via generating functions as follows. Let

$$(8) \quad B^*(z) = \sum_{k=1}^{\infty} z^k b_k = \sum_{i=1}^m (a_i z) / (1 - a_i z)$$

$$D^*(z) = \sum_{k=0}^{\infty} z^k d_k$$

with  $D^*(0) = d_0 = 1$ .

From (6) one obtains

$$z \frac{dD^*(z)}{dz} = \sum_{k=1}^{\infty} k z^k d_k = \sum_{k=1}^{\infty} z^k \sum_{j=1}^k d_{k-j} b_j$$

$$= \sum_{j=1}^{\infty} b_j z^j \left( \sum_{k=j}^{\infty} d_{k-j} z^{k-j} \right) = B^*(z) D^*(z),$$

$$\frac{d}{dz} \ln D^*(z) = \frac{B^*(z)}{z} = \sum_{i=1}^m a_i / (1 - a_i z).$$

Integrate upwards from  $z = 0$  to obtain

$$(9) \quad D^*(z) = 1 / \prod_{i=1}^m (1 - a_i z).$$

In a similar fashion, we write (3) as

$$(10) \quad y_i(k) = F_k(a_i) / G_k$$

and may solve for  $F_k$  and  $G_k$  via generating functions as follows. Let

$$\begin{aligned}
 F_k(a) &= \sum_{j=1}^k c(k)_j (a)^j = \sum_{j=1}^k d_{k-j} (a)^j \\
 (11) \quad F^*(z, a) &= \sum_{k=1}^{\infty} z^k F_k(a) = \sum_{k=1}^{\infty} \sum_{j=1}^k z^k d_{k-j} (a)^j \\
 &= \sum_{j=1}^{\infty} (az)^j \sum_{k=j}^{\infty} d_{k-j} z^{k-j} = \frac{azD^*(z)}{1-az}
 \end{aligned}$$

$$\begin{aligned}
 G_k &= \sum_{j=1}^k c(k)_j b_j = \sum_{j=1}^k d_{k-j} b_j \\
 (12) \quad G^*(z) &= \sum_{k=1}^{\infty} z^k G_k = B^*(z)D^*(z) \\
 &= \left[ \sum_{i=1}^m \frac{a_i z}{1-a_i z} \right] \frac{1}{\prod_{j=1}^m (1-a_j z)}.
 \end{aligned}$$

Let us first treat the special case where  $\{a_i\}_{i=1}^m$  are distinct. Notice  $G^*(z)$  has a double pole at each  $(a_j)^{-1}$ , and the following partial fraction expansion:

$$(13) \quad G^*(z) = \sum_{i=1}^m \left[ \frac{H_i}{(1-a_i z)^2} + \frac{Q_i}{(1-a_i z)} \right]$$

where

$$H_i = \lim_{z \rightarrow 1/a_i} (1-a_i z)^2 G^*(z) = 1 / \prod_{\substack{j=1 \\ j \neq i}}^m (1-a_j/a_i)$$

$$Q_i = H_i \left( -1 + \sum_{j \neq i} \frac{a_j/a_i}{1-a_j/a_i} \right) + \sum_{j \neq i} \frac{H_i}{1-a_j/a_i}.$$

This implies

$$(14) \quad G_k = \sum_{i=1}^m [(k+1)H_i + Q_i](a_i)^k \quad k = 1, 2, \dots$$

Similarly, from (11) follows

$$(15) \quad F^*(z, a_i) = \frac{a_i z}{(1-a_i z) \prod_{j=1}^m (1-a_j z)} \quad 1 \leq i \leq m$$

with a double pole at  $z = 1/a_i$  and a simple pole at all other  $1/a_j$ . A partial fraction expansion yields

$$(16) \quad F^*(z, a_i) = \frac{H_i}{(1-a_i z)^2} + \frac{R_i}{1-a_i z} + \sum_{\substack{j=1 \\ j \neq i}}^m \frac{B_{ij}}{(1-a_j z)}$$

where

$$B_{ij} = \lim_{z \rightarrow 1/a_j} (1-a_j z) F^*(z, a_i) = \frac{(a_i/a_j)H_j}{1-a_i/a_j}$$

$$R_i = F^*(0, a_i) - H_i - \sum_{\substack{j=1 \\ j \neq i}}^m B_{ij} = -(H_i + \sum_{\substack{j=1 \\ j \neq i}}^m B_{ij}).$$

This implies

$$(17) \quad F_k(a_i) = k H_i (a_i)^k - \left( \sum_{\substack{j=1 \\ j \neq i}}^m B_{ij} \right) (a_i)^k + \sum_{\substack{j=1 \\ j \neq i}}^m B_{ij} (a_j)^k \quad k = 1, 2, 3, \dots$$

Insertion of (14) and (17) into (10) produces an explicit expression for each  $y_i(k)$ . For the asymptotic behavior, let

$$(18) \quad a_{\max} = \max_{1 \leq i \leq m} a_i$$

and let  $i_0$  be the unique index with  $a_{i_0} = a_{\max}$ .

Put

$$\begin{aligned} F_k(a_i) &\sim k H_{i_0} (a_{\max})^k & a_i &= a_{\max} \\ &\sim B_{ii_0} (a_{\max})^k & a_i &\neq a_{\max} \\ G_k &\sim k H_{i_0} (a_{\max})^k \end{aligned}$$

into (10) to conclude

$$(19) \quad \begin{aligned} y_i(k) &\rightarrow 1 & a_i &= a_{\max} \\ &\rightarrow 0 \text{ as } k^{-1} & a_i &\neq a_{\max}. \end{aligned}$$

Returning now to the general case, retain the definition (18) and let  $p \geq 1$  denote the multiplicity of  $a_{\max}$  in  $\{a_i\}_{i=1}^m$ . Then (12) implies  $G^*(z)$  has a pole of order  $p+1$  at  $z = 1/a_{\max}$ , and

$$(20) \quad G_k \sim R k^p (a_{\max})^k$$

where

$$(21) \quad R = \lim_{z \rightarrow 1/a_{\max}} (1 - a_{\max} z)^{p+1} G^*(z) = p / \prod_{\substack{i=1 \\ a_i < a_{\max}}}^m (1 - a_i/a_{\max}).$$

Similarly, (15) implies  $F^*(z, a_i)$  has a pole at  $z = 1/a_{\max}$ , of order  $p+1$  if  $a_i = a_{\max}$  and of order  $p$  if  $a_i < a_{\max}$ . Consequently

$$(22) \quad F_k(a_{\max}) \sim S k^p (a_{\max})^k$$

where

$$(23) \quad S = \lim_{z \rightarrow 1/a_{\max}} (1 - a_{\max} z)^{p+1} F^*(z, a_{\max}) = R/p.$$

$$(24) \quad F_k(a_i) \sim T_i k^{p-1} (a_{\max})^k \quad \text{if } a_i < a_{\max}$$

where  $T_i$  is independent of  $k$ . Insertion of (20-24) into (10) now leads to

$$(25) \quad y_i(k) \rightarrow \frac{1}{p} \quad \text{if } a_i = a_{\max} \quad (\text{there are } p \text{ such } i\text{'s}) \\ \rightarrow 0 \quad \text{as } k^{-1} \quad \text{if } a_i < a_{\max}.$$

This completes the proof that  $\lim_{k \rightarrow \infty} y_i(k)$  exists.

### 3. APPENDIX: DERIVATION OF THE FINITE-DIFFERENCE EQUATION [1,2,3,]

Consider a closed network of queues with  $m$  servers and  $k \geq 1$  customers. Each customer has a mean service time  $S_i$  at server  $i$ , and when this service is completed he has probability  $p_{ij}$  of going next to server  $j$ . Here  $p_{ij} > 0$ ,  $\sum_{j=1}^m p_{ij} = 1$  and  $(p_{i,j})$  is assumed to be an irreducible Markov chain with a unique equilibrium distribution  $[\pi_i]$  satisfying  $\pi P = \pi$ ,  $\sum_{i=1}^m \pi_i = 1$  and  $\pi_i > 0$  for all  $i = 1, \dots, m$ .

The steady-state quantities of interest are  $\lambda_i(k)$  = throughput at server  $i$  (customers served per unit time),  $N_i(k)$  = mean number of customers in service or on queue at server  $i$ , and  $W_i(k)$  = mean waiting time (queueing plus service) experienced by customers at server  $i$ . These satisfy the three sets of equations



$$(A-1) \quad N_i(k) = \lambda_i(k)W_i(k) \quad 1 \leq i \leq m$$

$$(A-2) \quad \lambda_i(k) = k\pi_i / \left[ \sum_{j=1}^m \pi_j W_j(k) \right] \quad 1 \leq i \leq m$$

$$(A-3) \quad W_i(k) = S_i + S_i N_i(k-1) \quad 1 \leq i \leq m; \quad k \geq 1$$

with the understanding

$$(A-4) \quad N_i(0) = 0 \quad 1 \leq i \leq m.$$

Equation (A-1) is Little's formula applied to server  $i$ . Equation (A-2) says that each of the  $k$  customers has a mean time between visits to server  $i$  of  $(\pi_i W_i(k))/\pi_i$ . Together these two equations show

$$(A-5) \quad \sum_{i=1}^m N_i(k) = k.$$

Equation (A-3) expresses the waiting time  $W_i$  as the sum of the service time  $S_i$  and queueing time  $\bar{N}_i S_i$ , where  $\bar{N}_i$  is the mean number of customers on queue when a customer arrives. First-come first-served queue discipline and exponentially-distributed service times (with mean  $S_i$ ) are assumed; with the approximation  $\bar{N}_i = N_i(k-1)$ , equation (A-3) is obtained.[3]

If  $\lambda_i$  is eliminated via (A-2) and then  $W_i$  is eliminated via (A-3), the following equations for  $N_i$  are obtained:

$$(A-6) \quad N_i(k) = \frac{k\pi_i S_i [1+N_i(k-1)]}{\sum_{j=1}^m \pi_j S_j [1+N_j(k-1)]} \quad 1 \leq i \leq m.$$

Letting  $a_i = \pi_i S_i > 0$  and  $y_i(k) = N_i(k)/k$ , (A-6) reduces to (1) with boundary conditions (2).

The variable  $y_i(k)$  is the *fraction* of the  $K$  customers who are located at server  $i$ . These satisfy  $\sum_{i=1}^m y_i(k) = 1$  due to (A-5). The asymptotic result (25) says that as the number of customers becomes very large, almost all of them queue up at the set of most-heavily congested servers, who "pace" the rest of the system.

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