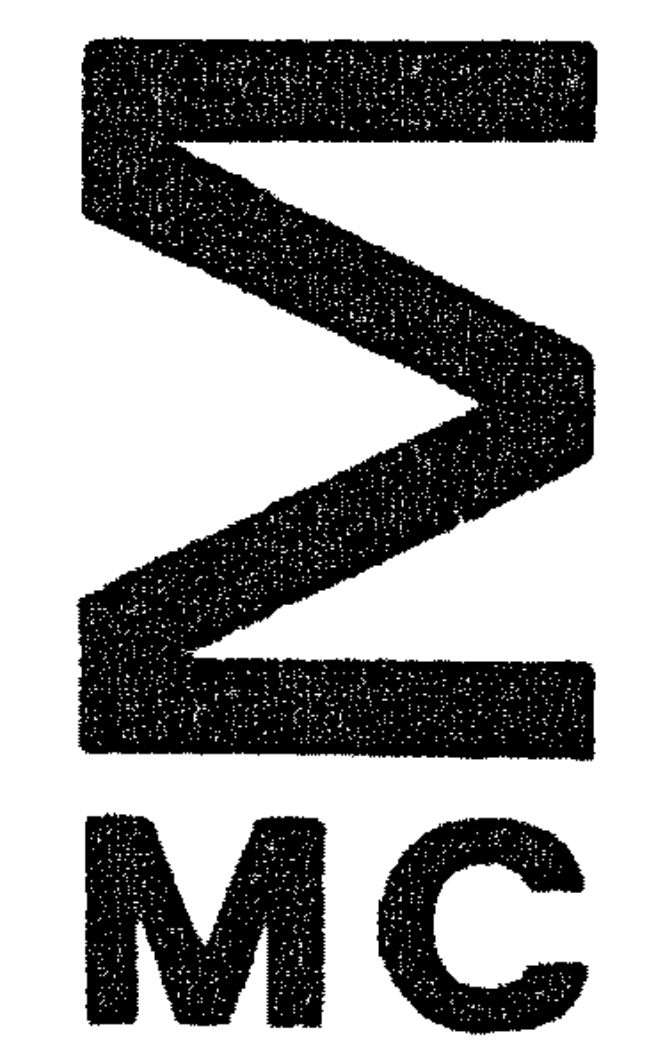


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AFDELING INFORMATICA

IW 3/73

MARCH

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MIXAL ASSEMBLY

2e boerhaavestraat 49 amsterdam

Printed at the Mathematical Centre, 49, 2e Boerhaavestraat, Amsterdam.

The Mathematical Centre, founded the 11-th of February 1946, is a non-profit institution aiming at the promotion of pure mathematics and its applications. It is sponsored by the Netherlands Government through the Netherlands Organization for the Advancement of Pure Research (Z.W.O), by the Municipality of Amsterdam, by the University of Amsterdam, by the Free University at Amsterdam, and by industries.

AMS (MOS) subject classification scheme: 68A10, 68A15

ACM-Computing Reviews-category 4.1, 4.2

Abstract

MIXAL is the assembly language for the (hypothetical) MIX computer, introduced by D.E. Knuth. In the present report, an assembly program for this language is described, based upon a slightly modified syntactic definition of MIXAL.

The assembler has been written first as an ALGOL 60 program. Care has been taken to provide it with a clear structure; in particular, the whole (12 pages) program contains only one goto statement. Next, the same algorithm has been implemented in MIXAL on the MIX computer. The latter version of the MIXAL assembler has been used as a MIXAL program to be assembled by the ALGOL 60 version. Both versions of the assembler have been reproduced in this report.

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1. Introduction.

In this report, an one-pass assembler for MIXAL is described, MIXAL being the assembly language for the (hypothetical) MIX computer, designed and described by Knuth [1].

The assembler has been written in ALGOL 60 as a boolean procedure, called "mixal assembler".

The assembler has been tested in the environment of a MIX simulator, written in ALGOL 60 by R.P. van de Riet [2].

The programs have been run on the EL X8 computer of the Mathematical Centre at Amsterdam, under supervision of the ALGOL 60 system, called MILLI [3].

The MIXAL system, consisting of assembler and simulator, accepts several MIXAL programs, textually being separated from each other by a job separator (a dollar symbol: \$), and which are translated and executed consecutively.

The interface of assembler and simulator consists of two parameters to procedure "mixal assembler", in addition to the boolean value, delivered by its procedure identifier. The first parameter "m" represents the memory of the MIX computer; actually it has to be an array, being declared with the bound pair 0:3999. All locations of this memory have to be set equal to zero, before calling procedure "mixal assembler"; those locations for which this is not the case, cannot be filled by the assembler (see section 5.3).

The second parameter "start address" is an output parameter of type integer, passing to the simulator the location at which the execution of the program has to be started.

At last, the boolean value, delivered by the procedure identifier of "mixal assembler", is the complement of the boolean variable "erroneous", which is local to procedure "mixal assembler"; "erroneous" indicates whether or not errors have been detected during assembly of the MIXAL program.

The assembler has to take into account that the simulator is a base ten machine (the bytesize being one hundred), with a special internal code for the characters, and a special operation code (a field specification (3:5) is coded as 35, being the value of $3 \cdot 10 + 5$, instead of $29 = 3 \cdot 8 + 5$).

Besides the ALGOL 60 version of the MIXAL assembler, we also have developed a version written in MIXAL, the algorithm being the same, except for some

very little details.

This MIXAL version has been reproduced in chapter 8, in the way it has been printed out by the MIXAL assembler in ALGOL 60, reproduced in chapter 7; thus we use this MIXAL version as a demonstration and a test case of the ALGOL 60 version. Finally, in chapter 9, we give a loading program, -appropriate to the MIX simulator, mentioned before- serving as a demonstration of the MIXAL version. As a thorough test for the MIXAL version, we applied the assembler to the MIXAL assembler in MIXAL itself, using however another MIX simulator, written in ELAN (being the assembly language of the EL X8 computer). This simulator satisfies the same specifications as the ALGOL 60 version (except for the memory size, which has been increased from 4000 to 10000), but it is much faster and can be operated by hand. The results of this test, however, have not been reproduced in this report.

2. Definition of MIXAL.

2.1 The syntax of MIXAL.

```
<MIX character> ::= <any available character on the MIX computer>
<number> ::= <digit> | <number><digit>
<identifier> ::= <letter> | <identifier> <letter> | <identifier> <digit>
<symbol> ::= <identifier> | <digit><symbol>
<comment> ::= <empty> | <comment><MIX character>
<separator> ::= ; | <nocr> | " <comment><nocr>
<local forward> ::= <digit> F
<local label> ::= <digit> H
<local backward> ::= <digit> B
<local symbol> ::= <local forward> | <local label> | <local backward>
<undefined symbol> ::= <symbol>
<defined symbol> ::= <symbol>
```


<unary operator> ::= + | -
<binary operator> ::= + | - | * | / | // | ;
<atomic expression> ::= <local backward> | <defined symbol> | <number> | *
<expression> ::= <atomic expression> | <unary operator><atomic expression> |
 <expression><binary operator><atomic expression>
<field specification> ::= (<expression>) | <empty>
<word value> ::= <expression><field specification> | <word value> ,
 <expression><field specification>

<literal constant> ::= = <word value> =
<future reference> ::= <local forward> | <undefined symbol> |
 <literal constant>

<address part> ::= <future reference> | <expression> | <empty>
<index part> ::= , <expression> | <empty>
<address value> ::= <address part><index part>

<MIX instruction symbol> ::= ADD | ... | SUB
<pseudo instruction symbol> ::= EQU | ORIG
<constant instruction symbol> ::= CON | ALF
<END instruction symbol> ::= END
<MIX machine instruction> ::= <MIX instruction symbol>
 <address value><field specification>
<pseudo instruction> ::= <pseudo instruction symbol><word value>
<char> ::= <MIX character>
<character code> ::= . <char><char><char><char><char>
<MIX constant instruction> ::= CON <word value> |
 ALF <character code>
<proper instruction> ::= <MIX machine instruction> |
 <pseudo instruction> |
 <MIX constant instruction>
<proper END instruction> ::= <END instruction symbol> . <word value>


```
<label> ::= <undefined symbol> | <local label>
<MIXAL instruction> ::= <proper instruction> | <label> : <MIXAL
    instruction> | <separator> <MIXAL instruction>
<END instruction> ::= <proper END instruction> | <label> : <END
    instruction> | <separator> <END instruction>
<MIXAL body> ::= <empty> | <MIXAL instruction><separator><MIXAL body>
<MIXAL program> ::= <MIXAL body><END instruction><separator>
```

2.2 Remarks.

The language MIXAL as defined in the preceding section, differs from the language defined by Knuth on a few points.

As one of the purposes has been to design a version of MIXAL being layout independent as much as possible (i.e. the position of any syntactical unit on the card or papertape should be of no significance), some additional punctuations have been introduced to prevent ambiguities; the main case being the colon following each label, the other cases being the quotation mark preceding comment, and the point in character code preceding the five MIX characters.

Layout independence means that spaces freely may be inserted into the MIXAL text; however, there are some exceptions:

- a) in character code, a space certainly is not layout, but stands for a space as MIX character,
- b) in numbers and symbols spaces may also be inserted as layout, but not two (or more) spaces consecutively.

The latter exception has been made as (in some situations) two symbols may succeed each other (or a symbol and a number), without being separated by another syntactical unit. Thus we use two (or more) spaces as separator between symbols and/or numbers.

A second difference is concerned with labels. The layout independence has made it possible to allow more than one label, preceding a proper instruction, each of these labels being equivalent to the same value.

As these changes do not essentially affect MIXAL as defined by Knuth, it suffices to refer to [1] for the meaning and use of the assembly language MIXAL.

3. Reading, printing and recognition of basic syntactical units.

We distinguish between two levels of basic reading procedures, the first one being concerned with the reading and printing of single MIX characters, and the second one combining these MIX characters to basic syntactical units, which are analyzed more closely.

3.1 The low level reading procedure "read char".

This procedure reads the next character of the MIXAL program from the input medium. Both the procedure identifier of "read char" and the global variable "char" get as value the internal code of this MIX character.

Furthermore, the listing of the MIXAL program -each line generally consisting of the address and code, corresponding to the MIXAL instruction, the line number, and the text of the MIXAL instruction- is arranged by "read char", storing the characters read into a buffer (the array "line"), and printing out this buffer at the right moment, i.e. generally after another part of the assembler has produced the code corresponding to the instruction of this line, and has placed it also in the buffer. The boolean variables "text" and "code" make it possible to indicate this moment outside procedure "read char". They indicate whether some text has been completed for printing, and whether address and code have been produced. Thus, "read char" builds up a line in the buffer, organized with two pointers "pos" and "last pos", running from zero upwards, the first one pointing at the first free position of the buffer, the latter one pointing at the lastly occupied position.

It will be clear that reading a <nocr> (new line carriage return, or new card) needs a special treatment, as, at the reading of this separator, the buffer generally may not be printed out (address and code may have to be determined yet). The <nocr> itself is not stored into the buffer, as it is not a MIX character actually, but, as an indication that a <nocr> just has been read, the pointer "pos" gets the value "-1", and the boolean variable "text" gets the value "true" (the latter being explained already).

When "read char" is called again, we may assume that address and code have been produced, if this were necessary. The value "true" for "text" (always) means that we may print this line before continuing, and the value "-1" for "pos" means that we have to start a new line at position zero, and that we also can store the new linenumber into the buffer.

During the initialisation of the assembler, when all MIX instruction symbols have to be read and stored into the symbol table, the high level reading procedures are used, which call "read char" in their turn. The text for the initialisation, however, is not read from an input medium, but is contained in the ALGOL 60 program by means of a string. Giving the boolean variable "from string" the value "true", the reading from this string is effectuated by means of the standard procedure "STRINGSYMBOL", called in "read from string", which delivers the representations of characters of a string [3].

remark:

When a dollar \$ (the job separator, mentioned in the introduction) is read, obviously the END instruction of the MIXAL program is missing. Having reported this error, we jump out of all nested procedure calls, back to the first statement of the assembler, in order to treat the next MIXAL program. We mention this jump explicitly here, as this is the only goto statement, appearing in the whole program.

3.2 Some auxiliary procedures, concerned with the output buffer.

The printing of the characters, contained in the buffer, is done by procedure "print line". This procedure prints the contents of the buffer (using the pointer "last pos"), and fills the buffer with spaces again.

In the first place, procedure "read char" is responsible for the filling of the buffer with text.

Furthermore, a procedure "buffer" is available, used for storing numbers into the buffer (used for address, code and linenumber). Procedure "buffer" has four formal parameters: "p", "n", "m" and "value". The first parameter "p" indicates the position in the buffer at which the units of the number, contained in "value", are placed. There is a possibility to insert spaces between the digits of the number, by means of the parameters "n" and "m".

The digits of the number are divided into "abs(n)" groups of "abs(m)" digits; between each two successive groups one space is inserted. The sign of "n" indicates whether or not (positive or negative, respectively) a sign has to be placed preceding the number; the sign of "m" indicates whether leading zero's are printed, or replaced by spaces (positive or negative, respectively).

Finally, a procedure "buffer symb" may be called, storing symbols from the symbol table into the buffer. Its first parameter indicates the place of the symbol in the symbol table (see section 3.4.), the second parameter indicates the position in the buffer at which the first character of the symbol is placed, and the third parameter is an output parameter, delivering the first free place in the buffer, following the last character of the symbol. Procedure "buffer symb" actually has been used for special purposes concerned with the completion of the assembly process (see section 6).

3.3 The high level reading procedure "read synt unit".

Procedure "read synt unit" combines a number of MIX characters to one basic syntactical unit. Most of the basic syntactical units consist of one MIX character only (*, :, ,, ,, (,), =, +, -, *, / and :); they need no further analysis. The basic syntactical units, consisting of more than one character are:
a) number, b) symbol, c) separator and d) special divisor (//).

a) Reading a number, we don't know whether we are actually reading a number or a symbol after all, before we have read it completely. That is why we treat a number as if it were a symbol; this means that the digits temporarily are stored into the symbol table (see b.). In the boolean variable "contains letter" we keep up whether at least one letter has been read; if this is the case, we are concerned with a symbol (discussed below), otherwise we have read a number, of which the value is calculated by means of the digits stored in the symbol table. This value is assigned to the global variable "value of number", and the procedure identifier of "read synt unit", as well as the global variable "synt unit", get the constant value of "number" (value 102).

b) Reading a symbol, the letters and digits are stored meanwhile into the symbol table, in a way to be considered later on (3.4.). When all letters and digits have been read, procedure "look for symbol" looks up the symbol just read among the symbols which have been stored already in the symbol table. If this symbol is not found, a new information cell in the symbol table is created, containing the letters and digits of the symbol already. Thus, the new symbol has been joined to the symbol table by procedure "look for symbol". By means of the function designator "look for symbol", the place in the symbol table is passed on to the variable "last symbol". Besides, the variable "type" gets more information about the kind of symbol just read. After all, the procedure identifier of "read synt unit" and "synt unit" get the constant value of "symbol" (value 103).

c) The separator is read entirely, which means that comment is skipped, up to the first <nocr>, by means of procedure "read char", and the procedure identifier of "read synt unit" (as well as "synt unit") gets the constant value of "separator" (value 101). Having read a semicolon as separator, the variable "text" is set on "true", as we expect that one MIXAL instruction has been completed, and the assembler will produce address and code, before the next character is read (at which moment the buffer will be printed out). The variable "pos" is left unchanged in this case; so succeeding characters will stay at the right position of the line, however, this will be a next line.

d) The reading of a special divisor (//) does not lead to any difficulty.

In most cases, the reading of one syntactical unit consists of the reading of one character, but, in some cases, we know that we have read one whole syntactical unit, just when we have read the first character of the following syntactical unit (this is the case when reading numbers, symbols and the division operator /).

Thus we arrive at the impossible situation that we sometimes have read already the first character of a following syntactical unit, and in other cases we have not. This can be resolved by consequently reading one

character too many in all cases, but this has an unpleasant effect on the synchronisation of possible error messages. We have chosen the following solution:

Instead of unnecessarily reading one character too many, we insert one extra character (a space as layout character), by simply giving the variable "char" the constant value of "space" (value 66) in those cases where we do not need to read the first character of the following syntactical unit. As, at the reading of the next syntactical unit, spaces are skipped beforehand by procedure "read synt unit", there will be no harmful effect of this insertion.

The main program of the assembler, and the expression reading procedures, mainly read the MIXAL program by means of this procedure "read synt unit", which means that one need not worry, at that level, about the listing of the MIXAL program, the insertion of, and search for symbols in the symbol table, and the reading of numbers.

3.4 The organisation of the symbol table.

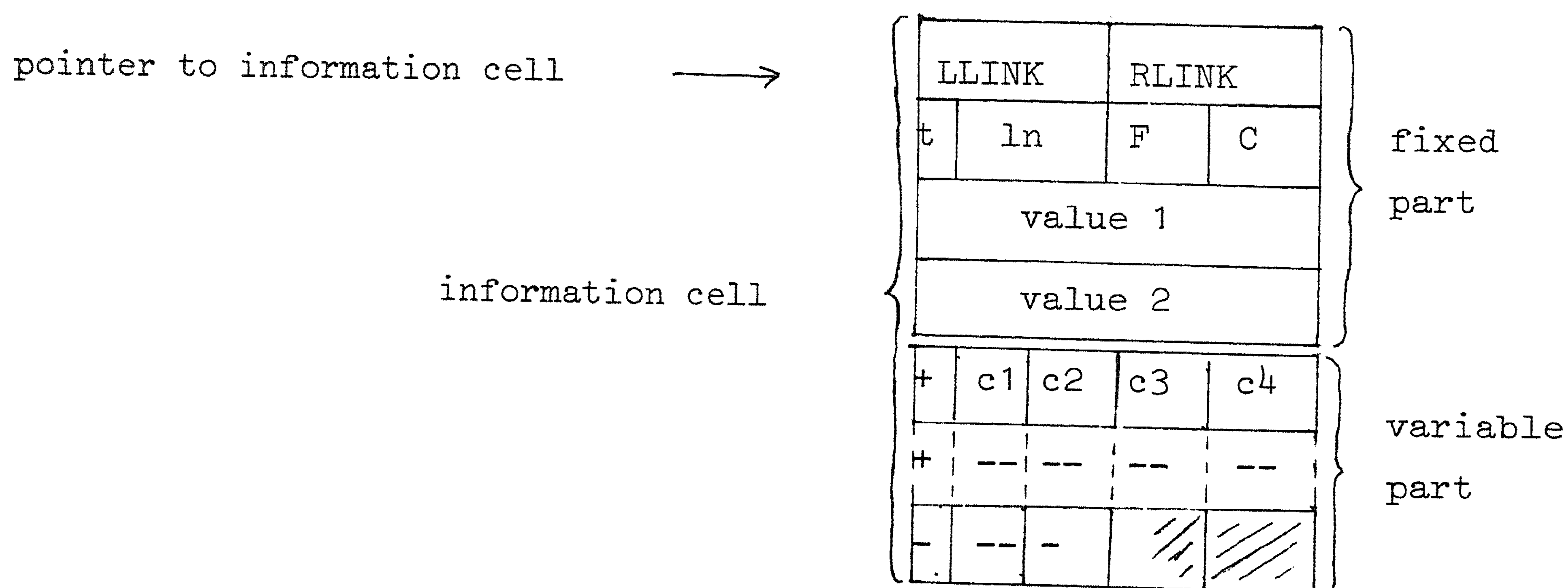
The symbol table is realised as an integer array "t", and consists of a set of information cells, each of which containing all information about one symbol. An information cell consists of a number of contiguous elements of "t", the address of the first element being used as the reference (pointer) of this information cell; the variable "last symbol" is used as a pointer to it. Local symbols, however, are not stored into this symbol table as other symbols are, but are detected explicitly in procedure "read synt unit", indicated by "type", while in these cases the variable "last symbol" gets as value the value of the digit, contained in the local symbol; this value can be seen as a pointer to an element of the arrays "h", "b", "f" and "defined local backward", into which the information concerning local symbols is stored. The set of information cells may be divided into three contiguous parts, the first two being filled during initialisation with the MIX instruction symbols, and with the constant instruction symbols, the END instruction symbol, and the pseudo instruction symbols. Into the third part, all other symbols, introduced by the MIXAL programmer, are stored. The boundary between the first

two parts is given by means of the variable "last inst" being a pointer to the information cell of the lastly stored MIX instruction symbol.

For the symbols contained in the second part, explicit pointers are used, viz.: "alf", "con", "end", "equ" and "orig".

Concerning the internal construction of information cells, there exists no difference between the distinct parts, as each instruction symbol (e.g. ADD, CON or END) may be used by the MIXAL programmer as its own symbol. For such a symbol there exists only one information cell (in the first or second part of the symbol table) into which both the information with respect to the MIX instruction, and the information concerning the symbol, used by the programmer as its own, are stored.

An information cell consists of two parts, a fixed part and a variable part, the latter being used to contain the internal codes of the letters and digits, the particular symbol consists of.



The last element of the variable part is made negative, as opposed to the other elements, indicating the end of the alpha numerical information. This element may contain fewer than four characters (the others always contain four of them), in which case it is filled up by dummy characters.

Next, the fixed part of the information cell is dealt with closely. The fields F and C are concerned with MIX instruction symbols only (the first part of the symbol table). F contains the standard field, and C the code of

the corresponding MIX instruction. These fields are filled during initialisation. The field *ln* is used to store the linenumber of the line, on which the symbol has occurred as label (has got a value). The field *t* (type) indicates whether the symbol in question is defined or undefined (negative or positive value, respectively), i.e. whether or not the symbol has become equivalent to some value by means of an occurrence as label.

For the fields *value 1* and *value 2*, we have to distinguish between undefined and defined symbols, the status being indicated by the value of the field *t* as described before. At information cells of undefined symbols, the field *value 1* may be used to hold a pointer to the information cell of the symbol, which directly precedes the symbol in question, both symbols being used as label preceding the same MIXAL instruction. If this preceding label happens to be a local label, the pointer to it has been inverted and decreased by one, forcing a negative value, indicating in this case a reference out of the symbol table.

In case the symbol in question is not preceded by another label, the initial value of zero (stored in procedure "look for symbol") is replaced by the value 9999, indicating an empty pointer (N.B. this value neither may be a reference into the symbol table, nor a reference to a local symbol). The field *value 2* may hold (if the symbol has occurred as future reference without being defined already) a pointer to that memory address of the MIX computer, at which the instruction has been placed, the address part of which contained the textually last future reference of the symbol in question. When the symbol has not (yet) occurred as future reference, *value 2* holds the empty pointer (N.B. in this case the value 9999 may not be a valid MIX address).

For local symbols, the fields *value 1* and *value 2* also exist, but they are recorded in the respective elements of the arrays "h" and "f".

Concerning defined symbols, both pointers in the fields *value 1* and *value 2* have become irrelevant, but merely the value, the symbol recently has become equivalent to, is important. This value is split up into two parts, called the head and the tail, as the numerical capacity of one element of "t" would not be sufficient to store a word value. The head, stored into the field *value 1*, has the value "abs (word value) ÷ 10¹⁶", and the tail, stored into the

field value 2, has the value "sign (word value)*(abs(word value)-head *10⁶)".

The remarks concerning defined symbols do not pertain to defined local symbols, as the fields value 1 and value 2 may be used again at a next occurrence of the local symbol as label. So, a distinct field is used for storing the value of the local symbol. This field is recorded in the array "b", and in an auxiliary array "defined local backward" we keep up in a boolean value whether the local symbol in question has occurred as local label at least once.

The search technique in the symbol table is a binary search technique in a tree-structured symbol table. To make a binary search applicable, we first have to define an order relation between symbols. The alphabetical ordering turned out to be an appropriate order relation for our purpose. As we also have to deal with digits, we have extended this ordering, defining the digits ordered naturally, preceding the letters (e.g. MIX, MIX 1, MIX 10, MIX 9, MIXAL, in this order). The tree structure of the symbol table is defined as follows:

1) Each node (actually an information cell) has at most two subtrees (called the left and right subtree, which are trees satisfying this definition), which are referred to by the pointers, contained in the fields LLINK and RLINK; if a subtree is not present, the corresponding field contains an empty reference (0).

2) All nodes in the left subtree of some node contain information cells of symbols, alphabetically preceding the symbol of this node, and all nodes of the right subtree contain information cells of symbols, succeeding the symbol of this node.

Due to this organisation, it is possible, after one comparison of an arbitrary symbol with the symbol contained in the root of the tree, to decide, in which of the two subtrees we have to search for this arbitrary symbol (if the symbols, just compared happen to be not the same). At each next comparison, we may exclude a whole class of symbols from the searching process. Trying to search for a symbol, which not yet has been inserted into the tree, we eventually come out at an empty pointer. At this point we know that an

information cell of this symbol does not occur in the symbol table, and we easily may join the new symbol to the symbol table, by replacing the empty pointer by a real pointer to the information cell (the variable part of which existing already) of the new symbol. Having filled in two empty pointers in the fields LLINK and RLINK of the newly created information cell, and initial values for the other fields (all within procedure "look for symbol"), we have extended the symbol table by one symbol, while the conditions, the tree structure satisfies, are still fulfilled. For, if we have to search for the recently inserted symbol, starting at the root of the tree, we do follow the same path through the tree as before, up to the node which recently has contained the empty pointer (nothing has been changed until so far). As this empty pointer has been replaced by the real pointer to the information cell of the symbol, we are searching for, we do not end in the blind alley now, but do find the appropriate information cell.

Finally, we direct our attention to the advantages gained by using the binary search. It may be clear that, offering new symbols to be inserted in alphabetical order, only branches of the tree to the right will originate, in which case we certainly do not gain any advantage over the linear search method. The order in which the MIXAL programmer introduces his symbols cannot be predicted, but, fortunately, we are able to introduce the MIX instruction symbols during initialisation, ordered in such a way, that the base of the tree will be filled homogeneously. The number of MIX instruction symbols amounts to 154 (including ALF, CON, etc.). Given a tree of seven levels, we can store into it $\sum_{i=0}^6 2^i = 127$ symbols. The remaining 27 symbols easily can be stored into the eighth level. This means that each of these instruction symbols can be found after eight comparisons at the most; the average number of comparisons amounting to:

$$1/154 * (\sum_{i=0}^6 2^i * (i+1) + 27 * 8) = 1001/154 = 6.5.$$

Thus we have made a considerable profit on the linear search method, which needs an average number of $154/2 = 77$ comparisons for an instruction symbol.

3.5 The initialisation of the symbol table and the global variables.

During initialisation, characters are read from the string, contained in procedure "read from string", as discussed already in section 3.1. This string satisfies the following syntactical structure.

```
<MIX characters> ::= l*;:.( ] = +-*/*: azbfh
<standard field and code> ::= <number (2.1)>
<MIX instruction definition> ::= <empty>|
    <MIX instruction symbol (2.1)><standard field and code>
<MIX instruction definitions> ::= <empty> | <MIX instruction definition >
    < separator > < MIX instruction definitions >
<constant instruction symbols> ::= ALF CON
<pseudo instruction symbols> ::= EQU ORIG

<initial string> ::= <MIX characters>
    <MIX instruction definitions>
    ( < constant instruction symbols >
    <END instruction symbol (2.1)>
    <pseudo instruction symbols>)
```

First, we read the MIX characters by means of procedure "read char", and assign their internal code to the variables: "underlining", "bar", "space", "asterisk", etc. These variables, strictly speaking, are constants, introduced to make the program more readable. Having assigned to the other constants and variables their (initial) values, and having filled the outputbuffer "line" with spaces, we continue with filling the first part of the symbol table with the MIX instruction symbols. Having constructed explicitly the root of the tree, the symbol table consists of (viz. an information cell with initialized fields LLINK, RLINK, t, value 1 and value 2), the MIX instruction definitions are read by means of procedure "read synt unit ", the insertion to the symbol table being automatically now. Following each symbol, we read the standard field and code, the value of this number being: (standard field * 100 + code), and record both quantities in the fields F and C respectively of the information

cell in question. Having read all MIX instruction definitions, terminated by a paren, the first part of the symbol table is completed by assigning to "last inst" the pointer to the last information cell.

In addition, the next five symbols are read, and pointers to their information cells are recorded in the variables "alf", "con", "end", "equ" and "orig", respectively.

Finally, having checked whether we have reached the closing parenthesis, we assign to the boolean variable "from string" the value "false" (the effect of which being discussed already in section 3.1), and finish the initialisation process.

3.6 Errormessages.

When during the assembly process an error is detected, procedure "error" is called. The actual parameter of "error" is a string, describing (in the English language) the kind of error detected. In procedure "error", first all characters contained in the outputbuffer are printed out by means of a call of "print line". Next, at a new line, the actual parameter of error is printed out. The pointer "pos" (3.2), holding the position in the buffer, remains unchanged; thus, following characters (when reading on) will be printed out at their correct position, some lines lower however.

We already have mentioned the synchronisation of errormessages (3.3).

Synchronisation means that the message is printed, textually immediately following those MIX characters, the errormessage is concerned with. In some cases however, we have read one character too many (see 3.3) when detecting an error, or even worse, a whole syntactical unit (including another following character). The latter phenomenon we meet at the reading of expressions (section 4). Reading too much means, in this context, that some text, belonging to following syntactical units, has been read already (while not being considered), merely to determine whether or not a whole unit (e.g. an expression) has been completed. In these cases, a possible errormessage will not be given in the right context.

Having detected and reported an error, the assembly process generally can be continued, and other errors may be detected and reported in the same way.

4. Expressions and word values.

4.1 Expressions.

As appears from the definitions in section 2.1, no priority rules exist for arithmetic operators in expressions; this means that operations are applied in order from left to right. Together with the fact that parenthesis are not allowed within expressions, this makes expressions very easy to evaluate; we do not need a stack for storing intermediate results, as these results immediately are used as left operand of an operator, the right operand being a single atomic expression, the value of which easily and directly is determined by means of procedure "read atomic expression". Of the atomic expressions, only symbols give rise to some further investigation. Symbols in expressions only may occur as defined symbol, or as defined local symbol, the values of which are denoted in the combined field value 1 and value 2, and in the array "b", respectively. The other possible occurrences of symbols are not admitted here, so they cause an error message, viz. "undefined symbol in expression", or "undefined local backward". How we can determine whether or not a symbol has been defined, has been explained already in section 3.4. The other possibilities for atomic expressions are very easy to deal with. When we are concerned with a number, we find its equivalent value in the global variable "value of number"; at this point it is important to observe that we have not read too much, as the following syntactical unit may be a number, the value of which would destroy the value of the global variable "value of number".

When we are concerned with an asterisk (called "self"), we only have to deliver the value of the location counter "lc" (see section 5).

For atomic expressions, we have to check, within procedure "read atomic expression", whether the equivalent value is not too large (exceeding the MIX word value). When executing arithmetical operations, by means of procedure "arithm", also capacity overflow may occur. In these cases an error message is given by the procedures concerned, viz. "overflow (atomic expression)", and "overflow (+)", "overflow (*)", etc., respectively, and, as resulting value, zero is delivered. The consequence of these tests,

within "read atomic expression" and "arithm", is that in procedure "read expression" only valid intermediate results occur, and no test on capacity overflow is necessary.

It is clear now that an expression can be evaluated in a straightforward manner by means of the procedures "read atomic expression" and "arithm". For reasons, later on to be explained, the first syntactical unit of an expression always has been read already when calling "read expression". Within this procedure we read the next syntactical unit and whenever this is an arithmetic operator, we go on reading an atomic expression, and executing the operation; otherwise, when no operator is read, we have done with the expression, and then, we did read one syntactical unit too much. This syntactical unit cannot be a syntactical unit belonging to the following MIXAL instruction, provided that two distinct MIXAL instructions are separated by a separator.

4.2 Word values.

A word value is used to denote the contents of one MIX word. According to the definition, we can describe a word value as:

$E_1 (F_1), E_2 (F_2), \dots, E_n (F_n), n \geq 1$ and E_i and F_i being expressions for $i=1, \dots, n$.

The value E_i is meant to be stored into the fields of a certain MIX word, specified by (F_i) , for $i=1, \dots, n$ consecutively. The field specification (F_i) may be omitted, in which case a standard field (0:5) applies. The desired result, delivered by procedure "read w value", is the eventual value which would remain at the memory address CON when executing the next fragment of a MIXAL program:

STZ CON; LDA C1; STA CON(F1); ...; LDA Cn; STA CON(Fn), where C_1, \dots, C_n denote the addresses at which the values of E_1, \dots, E_n are stored respectively. Using literal constants, the fragment also can read:

STZ CON; LDA = E1 =; STA CON (F1); ...; LDA = En =; STA CON (Fn).

The F_i ($i=1, \dots, n$), being the value of a field specification, must be of the form $L_i * 10 + R_i$ (denoted in MIXAL as $L_i : R_i$), where L_i and R_i have to satisfy: $0 \leq L_i \leq R_i \leq 5$.

Reading a word value, one syntactical unit too many is read (as it was the case in "read expression"), when testing whether or not we have done with the word value.

Also the word value (just as an expression) is followed by another syntactical unit which does not belong to the following MIXAL instruction. This means that we never incorrectly read (even when errors are detected) syntactical units belonging to following MIXAL instructions, at the level of the procedures "read expression" and "read w value".

5. The MIXAL instruction.

5.1 General outline.

An interesting feature of the MIXAL instruction is the fact that more than one label may precede the proper instruction. The first syntactical unit of an instruction has to be a symbol, being a label or an instruction symbol, the latter being the first syntactical unit of the proper instruction.

Finding out which of these two kinds of symbols we are concerned with, we have to consider the context in which the symbol occurs, as each instruction symbol also may be used as label. Checking whether the next syntactical unit is a colon, we are able to recognize a label as such.

However, if no label occurs, we thus have read the first syntactical unit belonging to the address value (if not empty), to the word value, or to succeeding character code.

When this last syntactical unit happens to be a symbol as well (this is very well possible), the global variables "last symbol" and "type", originally referring to the instruction symbol, have been overwritten; unless we have saved the original value of "last symbol" by means of a local variable of procedure "read mixal program" (actually called "inst"), we would have lost all information about the instruction symbol.

Another consequence of this looking for a colon is that the first syntactical unit of an expression, possibly succeeding the instruction symbol, has been read already, when we eventually come to the processing of the expression. (see section 4.1).

Now, the question arises, whether it is possible to give a symbol a value (make the symbol a defined symbol), immediately following the recognition of the symbol occurring as label. Unfortunately this value depends on the kind of proper instruction following. If this proper instruction happens to be the EQU pseudo instruction, the equivalent value for the label(s) is (are) denoted by the word value which follows EQU, otherwise the label(s) will become equivalent to the value which the location counter "lc" has at the moment

Consequently, we first have to read all labels preceding the proper instruction, including the instruction symbol (and the next syntactical unit to recognize the instruction symbol as such), before determining the equivalent value for these labels. Meanwhile, we retain all labels by linking them; each label is made to refer to its predecessor, and the first one refers to nothing, i.e. contains an empty pointer.

To hold this reference, the field value 1 of the concerning information cells in the symbol table is used. We already have seen that, for local labels, an element of the array "h" corresponds to the field value 1, and that a negative reference stands for a reference to "h" and a positive reference stands for a reference into the symbol table; this holds both for references contained in the symbol table and in the array "h" as well. Thus we have chained all labels, and a pointer to the last occurring label is kept in the variable "last label". However, we carefully have to avoid circular lists, originating when one symbol occurs more than once as label preceding the same instruction.

Symbols, which have been inserted in a chain as described above, can be recognized by the following two properties for their information cells:

- 1) The field t is positive (irrelevant in the case of local labels), which means that the symbol has not yet become a defined one;
- 2) The field value 1 has a value, unequal to the initial value zero, which means that the symbol has occurred as label, as value 1 contains a reference (possibly an empty reference).

Subsequently, we treat the proper instruction entirely, and finally we update the chained list of labels by following the references, while storing into the information cells the equivalent value of the symbols, and making

the fields t negative, to indicate the defined status of the symbols (by means of procedure "update").

If one of the chained symbols occurs in some expression, contained in the proper instruction, it has not yet the defined status, and an error message follows. In the case of the EQU pseudo instruction, this error message is appropriate, as the value of a symbol may not depend on itself. In the other cases, the value of such a symbol could have been known. For congruence reasons with the previous case, we decided however, to maintain the general scheme of treatment; the advantage being that one rule can be given, stating, at which moment symbols get a value, viz. after the entire instruction has been dealt with.

Some examples may serve as illustration:

- 1) F: EQU F+10; In this case, the value of F would depend on itself; we rightly get an error message
- 2) F: ADD F+10; An error message is given, as the value of F has not yet been assigned to F, when evaluating the expression F+10, although it is known.
- 3) F: ADD **10; The alternative construction of 2), using an asterisk which has the same value as F gets afterwards.

At the moment the equivalent value for the labels is known, we keep this value in the variable "equiv". When ultimately the procedure "update" is called, we pass on the value of "equiv" as actual parameter of "update".

5.2 Address value and field specification.

The following two facts, mentioned before, are recalled. Firstly, the pointer to the information cell, belonging to the instruction symbol, is hold in the variable "instr", local to procedure "read mixal program". This pointer is passed on to procedure "read mix machine inst" as actual parameter. The first action of procedure "read mix machine inst" will be then to fetch the standard field and code of the instruction from the fields F and C of the information cell concerned, and to assign these values to the variables "f part" and "c part", respectively.

Secondly, the first syntactical unit of address value and field specification has been read already at this moment. We start the actual processing by testing whether the address part contains a future reference, analyzing the present syntactical unit. When a future reference does occur, the variable "a part" will get the value of a reference, being discussed in the following section; otherwise, if the address part doesn't happen to be empty -i.e. if the present syntactical unit is not a comma (indicating an index part), nor a left parenthesis (indicating a field specification), nor a separator (indicating the absence of all quantities)-, we may expect an expression, treated by means of procedure "read expression", the value being assigned to the variable "a part". Of the quantities indexpart and field specification, easily can be determined now whether or not they are empty.

An empty index part gives rise to the assignment of the value zero to the variable "i part", and an empty field specification causes the standard field remaining in the variable "f part".

After all, the word value which has to be stored at the address indicated by "lc", can be assembled easily by means of the variables "a part", "i part", "f part" and "c part", and actually it is stored into the MIX memory by means of procedure "produce" (which also produces the output of address and code by means of the outputbuffer).

5.3 Future references.

Syntactically, a future reference may occur only as address part of a MIX machine instruction. The value for this future reference will become known in the future, when treating following instructions. The intention is, that, for the instruction in which the future reference occurs, code is produced as usual, except for the address part, which will be filled in later on: at the moment that the corresponding (local) label gets a value by means of procedure "update"; for literal constants it means at the end of the program (see section 6).

It is evident, that a symbol may occur as future reference more than once. This means that we have to update a number of MIX addresses, when the symbol in question gets its value. As for all these addresses it holds, that the

address part of the code has not yet been filled in, we can use these two bytes for linking the MIX addresses; thus we chain all addresses corresponding to a future reference of the same symbol. The beginning of this chain is kept in the symbol table, in the field value 2 of the information cell (for local forwards corresponding with an element of array "f"), and the chain is ended by means of an empty reference.

In the same way, a chain of literal constants is created, the access to it being kept in the variable "last constant". The difference between the chain of literal constants, and a chain of symbolic future references, is, that each MIX word in the first chain ultimately will obtain an unique address part (viz. the address at which the constant word value is stored by the assembler), while each MIX word of one other chain gets the same value in its address part. Chaining the MIX addresses, at which literal constants occur, is not sufficient; also the corresponding word values have to be retained somewhere until the end of the assembly process, when they are treated. We use the end of the array "t" for this purpose. For each word value to be stored, two consecutive elements of array "t" are used, identically to the way of coding a word value in the fields value 1 and value 2.

Describing the chaining of labels, we have seen that overwriting of references carefully has been avoided. Also in the case of future references, overwriting may take place, possibly leading to disastrous effects (such as an infinite loop of the assembler).

As we are able indeed, by means of the ORIG pseudo instruction, to start writing at each address we want, we necessarily have to protect ourselves against overwriting any address which contains a reference. As these addresses are hardly distinguishable from other ones - we would have to follow all references, starting at many places in the symbol table, and at the variable "last constant", to determine whether the address in question forms part of one of the chains - we decided to forbid each attempt to write on an address, which does not contain the value zero, the initial value of all addresses (see introduction).

We are aware of the fact that we unfortunately also forbid innocent

overwriting, but that is the price we have to pay, when we decide to desist from the very expensive algorithm, described above.

When the assembler detects an address, containing a value, unequal to the initial value zero, the error message "address already used" is given. Meeting an address, which does contain the value zero, necessarily one of the four following situations occurs:

- 1) the address never has been used,
- 2) the address has been filled by means of CON 0
- 3) the address has been filled by means of ALF .00000
- 4) the address has been filled by means of NOP 0

In the first three cases, the address does not contain a future reference, so there exist no danger (although it is a little inconsistent to do permit the overwriting in some cases). The address part of the instruction NOP may contain a future reference, previously occurred at address zero (in the address part of the NOP instruction, a reference to address zero is filled in). However, this may occur only once in a MIXAL program; thus, when we retain this address in a special variable "dangerous", we protect ourselves against overwriting of references, when testing whether a certain address contains the value zero, and this address is not the dangerous one.

6. The completion of the assembly process.

Having detected the END instruction, and having read the following word value, and assigned its value to the variable "start address", we have to deal with those symbolic future references, which have not yet been treated, and with the literal constants, by means of procedure "insert constants". At first, all information cells of local symbols and ordinary symbols are examined, the latter ones by means of procedure "traverse tree", which recursively walks through the tree, investigating the symbols in alphabetical order, and which calls at each node of the tree the procedure "define" as actual parameter for "investigate" (local symbols have been treated explicitly, but analogously to "define"). Only those (local) symbols, which have been used as future reference, without having occurred as (local)

label, need some special treatment. The information cells of these symbols, satisfy the following properties:

- 1) The field `t` is positive (irrelevant in the case of local labels), indicating that the symbol has not yet become a defined one;
- 2) The field value `2` contains a reference, unequal to the empty pointer, indicating that the symbol has been used as future reference;
- 3) The field value `1` still contains the initial value zero, indicating that the symbol has not occurred as label preceding the END instruction (this chain of symbols has not yet been updated).

For each of these symbols, consecutively one MIX word is reserved, by means of procedure "produce" (overwriting of addresses, which have been used already, still is impossible), and the addresses of these words are filled in at the locations at which the future references occurred, by means of procedure "update" (also procedure "produce" is called, for updating an address; now the overwriting test mechanism is eliminated, by testing whether the location counter does not point at this address).

To inform the MIXAL programmer, for which symbols a new MIX word is reserved, output is delivered, immediately following the END instruction, consisting of the symbol in question, followed by a colon and the pseudo instruction `CON 0` illustrating what happened.

Next, a number of words have to be reserved for the literal constants. Walking through the chain of occurrences of literal constants, we reserve for each one a consecutive MIX word and put the corresponding word value in it by means of procedure "produce", and we update the address part with the actual address of the corresponding word value. Having dealt with the insertions, the possible chain of labels, preceding the END instruction is updated.

Here we remark that property 3) for symbols to be inserted, is necessary, as we don't want that the labels preceding the END instruction, get their value by means of "insert constants", but by means of the call of "update", at the end of each instruction.

After all, for the convenience of the MIXAL programmer, a dump of the symbol table is given by means of procedure "dump symbol tabel". Thus, an alphabetical list of symbols used is printed out, each symbol preceded

by the value it has become equivalent to, and the linenumber of the line at which the symbol occurred as label. This information easily is fetched from the symbol table, which is walked through, in alphabetical order of symbols, by means of procedure "traverse tree" with the actual parameter "dump" for "investigate".

7. The ALGOL 60 procedure "mixal assembler".

Next follows the complete text of the MIXAL assembler in ALGOL 60 without further comment.


```
boolean procedure mixal assembler(m, start0); array m; integer start0;  
begin
```

```
boolean erroneous, from string, future, code, text;  
integer underlining, bar, space, asterisk, semicolon, colon, comma,  
point, paren, thesis, equal sign, plus, minus, times, div1, dec, quote,  
letter a, letter z, letter b, letter f, letter h, nlcr, tab,  
div2, symbol, number, sep,  
undefined, defined, local label, local backward, local forward,  
equ, orig, con, alf, end, last inst,  
last symbol, type, char, last constant, synt unit,  
read ptr, pos, last pos, line number, table ptr1,  
lit cons ptr, u1, b1, aux;  
real start,lc, dangerous, value of number, time 0;  
boolean array defined local backward[0 : 9];  
integer array line[-40 : 99], h,f[0 : 9], t[1 : 2000];  
real array b[0 : 9];
```

```
integer procedure read char;  
if from string  
then begin char:= read from string;  
read char:=char:= if char = 127 then 69  
else if char > 35 then char - 27  
else char  
end  
else begin integer s;  
if text  $\vee$  code then print line;  
if pos = -1  
then begin if line number < 6708  
then line number:= line number + 1;  
buffer(-4,-1,-4, line number); pos:=0  
end;  
s:= resym; for s:=s while s = 134 do s:= resym;  
if s = 135 then s:= 119;  
char:= if s = 127 then 69  
else if s > 35 then s - 27 else s;  
if char = underlining then u1:=1000  
else if char = bar then b1:=1000  
else begin if pos > 99  $\wedge$  char  $\neq$  nlcr  
then begin print line; pos:=0 end;  
if char = tab  $\vee$  char = nlcr  
then s:=ext(space);  
s:=2  $\times$  b1 + u1 + s; u1:=b1:=0;  
if s = 2028  $\vee$  s = 2055  
then begin errorrm( $\{$ end of input $\}$ );  
for aux:=0 step 1 until 3999 do  
m[aux]:=0; new page; goto begin  
end;  
line[pos]:=s; last pos:=pos;  
if char = tab then pos:=(pos+9):8x8  
else if char = nlcr  
then begin text:=true; pos:= -1 end  
else pos:=pos + 1  
end;  
end;
```



```
    if char > 99 then errorm({unknown character});  
    read char:=char  
end read char;
```

```
real procedure read 5 char;  
if synt unit = point  
then begin real value; integer i, c; value:=0;  
    for i:=1 step 1 until 5 do  
        begin c:=if char = nlcr then space else read char;  
            value:=value × 100 + (if c = nlcr then space else c)  
        end;  
        if char ≠ nlcr then char:= space;  
        read synt unit; read 5 char:= value  
    end  
else begin errorm({. required}); read 5 char:=0 end;
```

```
procedure print line;  
begin integer i, s; boolean u, b;  
    for i:= -40 step 1 until last pos do  
        begin s:=line[i]; u:=b:=false;  
            if s > 2000  
            then begin b:=true; s:=s - 2000 end;  
            if s > 1000  
            then begin u:=true; s:=s - 1000 end;  
            if u then prsym(ext(underlining));  
            if b then prsym(ext(bar));  
            prsym(s); line[i]:=ext(space)  
        end;  
        prsym(ext(nlcr)); code:=text:=false; last pos:=0  
end print line;
```

```
procedure errorm(s); string s;  
begin print line; printtext(s); prsym(ext(nlcr)); erroneous:=true end;
```

```
integer procedure read synt unit;  
begin for char:=char while char = space ∨ char = tab do read char;  
    if char < letter z  
    then begin integer p, n, q1, q, i, value;  
        boolean contains letter, compl;  
        contains letter:= letter a < char ∧ char < letter z;  
        p:=table ptr1 + 4; q1:=char; q:=read char; n:=1;  
        for value:=q1 + 1, value while 1 compl do  
            begin if char = space then read char;  
                if q = space then q:=char;  
                compl:= letter z < char;  
                if compl then value:= - value × 100⌊(4-n);  
                if compl ∨ n = 4  
                then begin t[p]:=value; value:=0;  
                    p:=p + 1; test on overflow(p)  
                end;  
                if 1 compl  
                then begin n:=if n < 4 then n + 1 else 1;  
                    value:=value × 100 + char + 1;
```



```
        if letter a < char & char < letter z
        then contains letter := true;
        read char
    end
end;
if p = table ptr1 + 5 & n = 2 & q1 < letter a &
(q = letter b ∨ q = letter f ∨ q = letter h)
then begin synt unit := symbol;
    type := if q = letter b then local backward
            else if q = letter f then local forward
            else local label;
    last symbol := q1
end
else if contains letter
then begin synt unit := symbol;
    last symbol := look for symbol(p);
    type := if t[last symbol+1] < 0
            then defined else undefined
end
else begin synt unit := number; value of number := 0;
    for p := table ptr1 + 4, p+1 while t[p-1] > 0 do
    begin value := abs(t[p]);
        for i := 3, i-1 while value ≠ 0 do
        begin aux := 100i; q := value : aux;
            value := value - q × aux;
            value of number := value of number × 10 + q - 1
        end
    end
end
end
else if char = quote ∨ char = nlcr
then begin synt unit := sep;
    for char := char while char ≠ nlcr do read char;
    char := space
end
else if char = semicolon
then begin synt unit := sep; char := space; text := true end
else if char ≠ div1
then begin synt unit := char; char := space end
else if read char = div1
then begin synt unit := div2; char := space end
else synt unit := div1;
    read synt unit := synt unit
end read synt unit;
```

```
integer procedure look for symbol(ptr); value ptr; integer ptr;
begin integer p, next, i, wp, w, comp, llink, rlink;
    for p := 1, next while p ≠ 0 do
    begin for i := 4, i+1 while comp = 0 & wp > 0 & w > 0 do
        begin wp := t[p + i]; w := t[table ptr1 + i];
            comp := abs(w) - abs(wp)
        end;
    if comp = 0 then comp := if wp > 0 then -1 else
        if w > 0 then 1 else 0;
    if comp = 0
```



```
then begin look for symbol:=p; next:=0;
      if table ptr1 = 1 then table ptr1:=ptr
      end
else begin aux:=t[p]; llink:=aux : 104; rlink:= aux - llink × 104;
      next:=if comp < 0 then llink else rlink;
      if next = 0
      then begin look for symbol:=table ptr1;
            if comp < 0
            then llink:=table ptr1
            else rlink:=table ptr1;
            t[p]:=llink × 104 + rlink;
            t[table ptr1 ]:=0;
            t[table ptr1 + 1]:=104;
            t[table ptr1 + 2]:=0;
            t[table ptr1 + 3]:=9999;
            table ptr1:=ptr
            end
      end
end
end look for symbol;
```

```
procedure test on overflow(ptr); value ptr; integer ptr;
if ptr > lit cons ptr
then begin errorm({overflow symbol table});
      for char:=char while true do read char
end test on overflow;
```

```
real procedure read atomic expression;
begin real at expr; read atomic expression:=at expr:=0;
  if synt unit = symbol
  then begin if type = defined
            then begin aux:=t[last symbol + 3];
                    at expr:=sgn(aux)×(t[last symbol+2]×106+abs(aux))
            end
            else if type = local backward
            then begin if defined local backward[last symbol]
                        then at expr:=b[last symbol]
                        else errorm({undefined local backward});
            end
            else if type = local forward
            then errorm({local forward in expression});
            else if type = local label
            then errorm({local label in expression});
            else errorm({undefined symbol in expression});
            read synt unit
        end
  else if synt unit = number
  then begin at expr:=value of number; read synt unit end
  else if synt unit = asterisk
  then begin at expr:=lc; read synt unit end
  else errorm({atomic expression required});
  if abs(at expr) < 1010
  then read atomic expression:=at expr
  else errorm({overflow(atomic expression)});
end read atomic expression;
```



```

real procedure read expression;
begin real expr; integer operator; read expression:=0;
  expr:=if 1 (synt unit = plus ∨ synt unit = minus)
    then read atomic expression else 0;
  for synt unit:=synt unit
  while synt unit = plus ∨ synt unit = minus ∨ synt unit = times ∨
    synt unit = div1 ∨ synt unit = div2 ∨ synt unit = dec
  do begin operator:=synt unit; read synt unit;
    expr:=arithm(expr, operator, read atomic expression)
  end;
  read expression:=expr
end read expression;

```

```

real procedure arithm(a, op, b); value a,op,b; real a,b; integer op;
if op = plus
then begin real value; value:=a + b; arithm:=0;
  if abs(value) < 1010 then arithm:=value
  else errorm(overflow (+))
end
else if op = minus
then begin real value; value:=a - b; arithm:=0;
  if abs(value) < 1010 then arithm:=value
  else errorm(overflow (-))
end
else if op = times
then begin integer s,h1,t1,h2,t2,c; real r0,r1; arithm:=0;
  s:=sgn(a × b); a:=abs(a); b:=abs(b);
  h1:=entier(a/106); t1:=a - h1 × 106;
  h2:=entier(b/106); t2:=b - h2 × 106;
  r0:=t1 × t2; c:=entier(r0/106); r0:=r0 - c × 106;
  r1:=t1×h2 + h1×t2 + c; c:=entier(r1/104); r1:=r1 - c×104;
  if h1 × h2 + c = 0 then arithm:=s × (r1 × 106 + r0)
  else errorm(overflow (×))
end
else if op = div1
then begin arithm:=0; if b ≠ 0
  then arithm:=sgn(a × b) × entier(abs(a/b))
  else errorm(overflow (/))
end
else if op = div2
then begin arithm:=0; if b ≠ 0 ∧ abs(a) < abs(b)
  then arithm:=sgn(a×b)×entier(abs(a/b)×1010)
  else errorm(overflow (//))
end
else begin real value; value:=a × 10 + b; arithm:=0;
  if abs(value) < 1010 then arithm:=value
  else errorm(overflow (:))
end arithm;

```

```

real procedure read w value;
begin real con, a, field, aux1, aux2, aux3;
  integer l,r,s; boolean comma read;
  con:=0;
  for synt unit:=synt unit, synt unit while comma read do

```



```
begin a:=read expression;
  if synt unit = paren
  then begin read synt unit; field:= read expression;
    if field = 5 then con:= a
    else if field = 0 then con:= sgn(a) × abs(con)
    else if field = 15 then con:= sgn(con) × abs(a)
    else begin l:=entier(field/10); r:=field - l × 10;
      if 0 < l ∧ l < r ∧ r < 5
      then begin if l ≠ 0 then s:=sgn(con)
        else begin s:=sgn(a); l:=1 end;
        a:=abs(a); con:=abs(con);
        aux1:=100l(6-l); aux2:=100l(5-r);
        aux3:=aux1/aux2;
        con:=s × (entier(con/aux1) × aux1 +
          (a - entier(a/aux3)×aux3) × aux2 +
          con - entier(con/aux2) × aux2)
      end
    else errorm(⟨illegal field⟩)
  end;
  if synt unit = thesis then read synt unit
  else errorm(⟨required⟩)
end
else con:=a;
  if synt unit = comma
  then begin comma read:=true; read synt unit end
  else comma read:=false
end;
read w value:=con
end read w value;
```

```
procedure read mixal program;
begin integer inst, last label; real equiv; last label:= 9999;
  for inst:= 0, inst while inst ≠ end do
  for synt unit:=colon, synt unit while synt unit = colon do
  if read synt unit ≠ symbol
  then begin if synt unit ≠ sep then errorm(⟨symbol required⟩);
    synt unit:= colon
  end
  else begin inst:=last symbol;
    if read synt unit = colon
    then begin if type = undefined
      then begin if t[last symbol+2] = 0
        then begin t[last symbol+2]:=last label;
          last label:=last symbol
        end
      else errorm(⟨label used twice⟩)
    end
    else if type = local label
      then begin if h[last symbol] = 0
        then begin h[last symbol]:=last label;
          last label:=last symbol-1
        end
      else errorm(⟨local label used twice⟩)
    end
    else errorm(⟨illegal label⟩)
  end
end
```



```

else begin equiv:=lc;
      if inst < last inst
      then produce(lc,read mix machine inst(inst))
      else if inst = equ
      then equiv:=read w value
      else if inst = orig
      then lc:=read w value
      else if inst = con
      then produce(lc,read w value)
      else if inst = alf
      then produce(lc,read 5 char)
      else if inst = end
      then begin start:=read w value;
            insert constants;
            equiv:=lc
            end
      else
      errorm(⟨unknown instruction⟩);
      if synt unit ≠ sep
      then errorm(⟨separator required⟩);
      for synt unit:=synt unit while synt unit ≠ sep do
      read synt unit;
      update(last label, equiv); last label:=9999
      end
end
end read mixal program;

```

```

boolean procedure possible(address); value address; real address;
possible:=if 0 < address ∧ address < 3999
          then m[address] = 0 ∧ address ≠ dangerous
          else false;

```

```

real procedure read mix machine inst(instruction); value instruction;
integer instruction;
begin real a part,i part,f part,c part;
  a part:=i part:=0; aux:=abs(t[instruction + 1]); f part:=aux : 102;
  c part:=aux - f part × 102; f part:=f part - f part : 102 × 102;
  if synt unit = symbol
  then begin if type = undefined
            then begin a part:=t[last symbol + 3]; future:=true;
                  if possible(lc) then t[last symbol + 3]:=lc;
                  read synt unit
            end
            else if type = local forward
            then begin a part:=f[last symbol]; future:=true;
                  if possible(lc) then f[last symbol]:=lc;
                  read synt unit
            end
            else a part:=read expression
            end
  else if synt unit = equal sign
  then begin real value;
        a part:=last constant; future:=true;
        if possible(lc) then last constant:=lc;
        read synt unit; value:=read w value;

```



```

        lit cons ptr:=lit cons ptr - 2;
        test on overflow(table ptr1);
        t[lit cons ptr + 1]:=aux:=entier(abs(value)/106);
        t[lit cons ptr + 2]:=sgn(value)×(abs(value) - aux×106);
        if synt unit = equal sign then read synt unit
        else errorm(⚡= required⚡)
    end
    else if 1 (synt unit=comma ∨ synt unit=paren ∨ synt unit=sep)
        then a part:=read expression;
    if synt unit = comma
    then begin read synt unit; i part:=read expression end;
    if synt unit = paren
    then begin read synt unit; f part:=read expression;
        if synt unit = thesis then read synt unit
        else errorm(⚡) required⚡)
    end;
    if abs(a part) > 9999
    then begin errorm(⚡abs(a part) > 9999⚡); a part:=0 end;
    if i part < 0 ∨ i part > 99
    then begin errorm(⚡i part < 0 ∨ i part > 99⚡); i part:=0 end;
    if f part < 0 ∨ f part > 99
    then begin errorm(⚡f part < 0 ∨ f part > 99⚡); f part:=0 end;
    read mix machine inst:=sgn(a part) ×
        (abs(a part) × 106 + i part × 104 + f part × 102 + c part)
end read mix machine inst;

```

```

procedure produce(address,value);value address,value;real address,value;
begin if code then print line;
    if 1 (0 < address ∧ address < 3999)
    then begin errorm(⚡wrong address⚡);buffer(-28,1,-10,address) end
    else begin if address = lc ∧ (address = dangerous ∨ m[address] ≠ 0)
        then errorm(⚡address already used⚡)
        else begin if future ∧ value = 0 then dangerous:=address;
            m[address]:=value;
            buffer(-14, if future then -3 else 5, 2, value)
        end;
        buffer(-35, -2, 2, address);
        if address = lc then lc:=lc+1 else line[-40]:=ext(asterisk)
    end;
    future:=false; code:=true
end produce;

```

```

procedure update(label,equiv);value label,equiv;integer label;real equiv;
for label:=label while label ≠ 9999 do
begin integer symb, link, address;
    if label > 0
    then begin symb:=label; label:=t[symb + 2];
        link:=t[symb + 3]; aux:=t[symb + 1];
        t[symb + 1]:= -(aux + (line number+1 - aux:104)×104);
        t[symb + 2]:=aux:=entier(abs(equiv)/106);
        t[symb + 3]:=sgn(equiv) × (abs(equiv) - aux × 106)
    end
    else begin symb:=-label-1; label:=h[symb]; h[symb]:=0;
        link:=f[symb]; f[symb]:=9999; b[symb]:=equiv;

```



```

        defined local backward[symb]:=true
    end;
    for link:=link while link ≠ 9999 do
    if abs(equiv) > 9999
    then begin link:=9999; errorm({abs(future ref) > 9999}) end
    else begin address:=link; link:=entier(m[address]/106);
        produce(address,sgn(equiv)×(m[address]+(abs(equiv)-link)×106))
    end
    end update;

```

```

procedure buffer(p,n,m,value); value p,n,m,value; integer p,n,m; real value;
begin integer i, j, q; real aux1;
    if n > 0
    then line[p - (abs(m)+1)×n - 1]:=ext(if value<0 then minus else plus);
    aux1:=104abs(n×m); value:=abs(value) - entier(abs(value)/aux1) × aux1;
    for i:=abs(n)-1 step -1 until 0 do
    for j:=abs(m)-1 step -1 until 0 do
    begin aux1:=104(m × i + j); q:=entier(value/aux1);
        value:=value - q × aux1;
        if q > 0 ∨ (i = 0 ∧ j = 0) then m:=abs(m);
        if m > 0 then line[p - ((m+1) × i + j)]:=ext(q)
    end
    end buffer;

```

```

procedure traverse tree(investigate, symbol); value symbol;
integer symbol; procedure investigate;
if symbol ≠ 0
then begin integer llink, rlink; aux:=t[symbol];
    llink:=aux : 104; rlink:=aux - llink × 104;
    traverse tree(investigate, llink);
    investigate(symbol);
    traverse tree(investigate, rlink)
    end traverse tree;

```

```

procedure buffer symb(s, lb, ub); value s, lb; integer s, lb, ub;
if s < 0
then begin line[lb]:= -s-1; line[lb+1]:=ext(letter h); ub:=lb+2 end
else begin integer p, q, value, i; ub:=lb;
    for p:=s + 4, p + 1 while t[p - 1] > 0 do
    begin value:=abs(t[p]);
        for i:=3, i - 1 while value ≠ 0 do
        begin aux:=100i; q:=value : aux;
            value:=value - q × aux;
            if ub = 99
            then begin line[99]:=2000 + ext(space);
                last pos:=99; print line; ub:=0
            end;
            line[ub]:=q - 1; ub:=ub + 1
        end
    end
    end buffer symb;

```



```
procedure buffer constant(s); value s; integer s;  
begin integer ptr;  
  if code  $\vee$  text then print line;  
  buffer symb(s,0,ptr); line[ptr]:=ext(colon);  
  if ptr > 15 then begin last pos:=ptr; print line end;  
  buffer symb(con,16,ptr); last pos:=ptr + 5; line[last pos]:=ext(0)  
end buffer constant;
```

```
procedure insert constants;  
begin procedure define(symbol); value symbol; integer symbol;  
  if t[symbol+1] > 0  $\wedge$  t[symbol+2] = 0  $\wedge$  t[symbol+3]  $\neq$  9999  
  then begin integer equiv; equiv:=lc;  
    buffer constant(symbol); produce(lc, 0);  
    t[symbol+2]:=9999; update(symbol, equiv)  
  end define;  
  integer symb, address, ln;  
  code:=true; ln:=line number; line number:=6709;  
  for symb:= -1 step -1 until -10 do  
  if f[-symb-1]  $\neq$  9999  
  then begin integer equiv; equiv:=lc;  
    buffer constant(symb); produce(lc, 0);  
    h[-symb-1]:=9999; update(symb, equiv)  
  end;  
  traverse tree(define, 1);  
  line number:=ln;  
  for address:=last constant while address  $\neq$  9999 do  
  begin real value; last constant:=entier(m[address]/106);  
    value:=m[address] + (lc - last constant)  $\times$  106;  
    aux:=t[lit cons ptr + 2];  
    produce(lc, sgn(aux) $\times$ (t[lit cons ptr + 1]  $\times$  106 + abs(aux)));  
    lit cons ptr:=lit cons ptr + 2;  
    produce(address, value)  
  end  
end insert constants;
```

```
procedure dump(symbol); value symbol; integer symbol;  
if t[symbol+1] < 0  
then begin integer ptr; real value; aux:=t[symbol+3];  
  value:=sgn(aux) $\times$ (t[symbol+2] $\times$ 106+abs(aux));  
  line number:=abs(t[symbol+1])  $\div$  104 - 1;  
  buffer(-14, 5, 2, value);  
  if line number > 6708 then line[-3]:= ext(plus)  
  else buffer(-4, -1, -4, line number);  
  buffer symb(symbol, 0, ptr); last pos:= ptr - 1;  
  print line  
end dump;
```

```
procedure dump symbol table;  
begin new page;  
  printtext({value and line number of defined symbols:});  
  prsym(ext(nlcr)); prsym(ext(nlcr));  
  traverse tree(dump, 1); new page  
end dump symbol table;
```



```

procedure initialize;
begin integer i;
  procedure rc(s); integer s; s:=read char;
  erroneous:=false; from string:=true; read ptr:=0;
  rc(underlining); rc(bar); rc(space); rc(asterisk); rc(semicolon);
  rc(colon); rc(comma); rc(point); rc(paren); rc(thesis); rc(equal sign);
  rc(plus); rc(minus); rc(times); rc(div1); rc(dec); rc(letter a);
  rc(letter z); rc(letter b); rc(letter f); rc(letter h); read char;
  quote:=94; nlcr:=92; tab:=91;
  sep:=101; number:=102; symbol:=103; div2:=104;
  undefined:=0; defined:=1; local label:=2; local backward:=3;
  local forward:=4;
  table ptr1:=1; lit cons ptr:=2000; t[1]:=t[3]:=0; t[4]:=9999;
  for i:= -40 step 1 until 99 do line[i]:=ext(space);
  for i:=0 step 1 until 9 do
  begin h[i]:=0; f[i]:=9999; defined local backward[i]:=false end;
  last constant:=9999; lc:=0; dangerous:= -1;
  future:=code:=text:=false;
  pos:=last pos:=line number:=u1:=b1:=0;
  for synt unit:=sep, synt unit while synt unit = sep do
  if read synt unit = symbol
  then begin read synt unit;
    t[last symbol + 1]:=104 + read atomic expression
  end
  else if synt unit ≠ sep ∧ synt unit ≠ paren then error(er 01);
  if synt unit = paren
  then begin last inst:=last symbol;
    read synt unit; alf:=last symbol;
    read synt unit; con:=last symbol;
    read synt unit; end:=last symbol;
    read synt unit; equ:=last symbol;
    read synt unit; orig:=last symbol;
    if read synt unit ≠ thesis then error(er 03)
  end
  else error(er 02);
  if erroneous then
  begin error(error detected during initialisation of assembler);
    exit
  end;
  from string:=false
end initialize;

```

```

integer procedure read from string;
begin read from string:=stringsymbol(read ptr, |x;:,.( )=+-x/:azbfh

```

j3np	553;	ent5	265;	ld4	514;	dec6	166;	inc6	066
j1	447;	slax	206;	cmp6	576;	enn5	365;	fsub	602
j1p	251;	j6n	056;	jxp	257;	ldx	517;	st4	534
cmp2	572;	dec2	162;	enn1	361;	ent1	261;	fadd	601
inc2	062;	j1n	051;	j2np	552;	j4p	254;	janp	550

jxn	057;	ld2	512;	ld6	516;	nop	000;	src	506
stj	240;	char	105;	cmp4	574;	cmpx	577;	dec4	164
decx	167;	enn3	363;	enna	360;	ent3	263;	enta	260
fdiv	604;	in	044;	inc4	064;	incx	067;	j1np	551
j2n	052;	j2p	252;	j4n	054;	j5np	555;	j6p	256

jbus	042;	jnov	347;	jxnp	557;	ld1	511;	ld3	513
ld5	515;	lda	510;	move	107;	out	045;	sra	106
st2	532;	st6	536;	stz	541;	add	501;	cmp1	571
cmp3	573;	cmp5	575;	cmpa	570;	dec1	161;	dec3	163
dec5	165;	deca	160;	div	504;	enn2	362;	enn4	364

enn6	366;	ennx	367;	ent2	262;	ent4	264;	ent6	266
entx	267;	fcmp	670;	fmul	603;	hlt	205;	inc1	061
inc3	063;	inc5	065;	inca	060;	ioc	043;	j1n	351
j1nz	451;	j1z	151;	j2nn	352;	j2nz	452;	j3n	053
j3p	253;	j4np	554;	j5n	055;	j5p	255;	j6np	556

jan	050;	jap	250;	jg	647;	jmp	047;	jred	046
jxn	357;	jxnz	457;	jxz	157;	ld1n	521;	ld2n	522
ld3n	523;	ld4n	524;	ld5n	525;	ld6n	526;	ldan	520
ldxn	527;	mul	503;	num	005;	sla	006;	slc	406
srax	306;	st1	531;	st3	533;	st5	535;	sta	530

stx	537;	sub	502;	j2z	152;	j3nn	353;	j3nz	453
j3z	153;	j4nn	354;	j4nz	454;	j4z	154;	j5nn	355
j5nz	455;	j5z	155;	j6nn	356;	j6nz	456;	j6z	156
jann	350;	janz	450;	jaz	150;	je	547;	jge	747
jle	947;	jne	847;	jov	247;	jsj	147		

```

      (      alf   con   end   equ   orig )      †);
  read ptr:=read ptr + 1
end read from string;

```

```

integer procedure sgn(a); value a; real a;
sgn:=if a < 0 then -1 else +1;

```

```

integer procedure ext(c); value c; integer c;
ext:=if c = 69 then 127
      else if c > 35 then c + 27
      else c;

```

```

begin: initialize; time 0:= time;
      read mixal program; print line;
prsym(ext(nlcr)); printtext(†start address: †); fixt(10,0,start);
prsym(ext(nlcr));
if 1 (0 < start ^ start < 3999)
then error(†wrong start address†)

```



```
else start0:=start;  
if erroneous then errorm(⟨mixal program incorrect⟩);  
prsym(ext(nlcr)); printtext(⟨translation time:⟩);  
absfixt(4,2,time - time 0); printtext(⟨sec.⟩);  
dump symbol table;  
mixal assembler:= 7 erroneous
```

```
end mixal assembler;
```



```

1  "CHAPTER 8      MIXAL ASSEMBLER DESCRIBED IN MIXAL.
2
3
4
5  " UNIT NUMBERS FOR APPARATUS
6  READER:        EQU    16          " CARD READER
7  CPUNCH:        EQU    17          " CARD PUNCH
8  PRINTER:       EQU    18          " LINE PRINTER
9
10
11 " INTERNAL CODE OF CHARACTERS
12 LETTER B:      EQU    11          " B
13 LETTER F:      EQU    15          " F
14 LETTER H:      EQU    17          " H
15 LETTER Z:      EQU    35          " Z
16 PLUS:          EQU    37          " +
17 MINUS:         EQU    38          " -
18 TIMES:         EQU    38          " *
19   ASTERISK:    EQU    39          " *
20 DIV:           EQU    40          " /
21 EQUAL SIGN:    EQU    43          " =
22 COMMA:         EQU    60          " ,
23 POINT:         EQU    61          " .
24 DEC:           EQU    61          " :
25   COLON:       EQU    63          " :
26 SEMI COLON:   EQU    64          " ;
27 SPACE:        EQU    66          "
28 PAREN:         EQU    71          " (
29 THES IS:      EQU    72          " )
30 NLCR:         EQU    92          "
31 QUOTE:        EQU    94          " "
32
33
34 " CONSTANTS USED BY ASSEMBLER
35 SEP:           EQU    101
36 NUMBER:        EQU    102
37 SYMBOL:        EQU    103
38 DIV2:          EQU    104          " //
39 UNDEFINED SYMB: EQU    0
40 DEFINED SYMB:  EQU    1
41 LOCAL LABEL:   EQU    2
42 LOCAL BACKWARD: EQU    3
43 LOCAL FORWARD: EQU    4
44 VALUE1:        EQU    0 : 2       " SPECIFICATION OF FIELD VALUE1
45 VALUE2:        EQU    4 : 5       " SPECIFICATION OF FIELD VALUE2
46 MOD:           EQU    0           " CONTENTS TO BE MODIFIED
47 N:             EQU    - 1         " USED TO NEGATE EXPRESSIONS
48 LENGTH OF TABLE: EQU    1500
49
50
51 " WORKING SPACE OF ASSEMBLER
52 OR'G           EQU    1200
53 BEG'N OF ASS:
54
55
56 CHAR:          CON    0
57 S CHAR:        CON    0
58 LINE NUMBER:   CON    0

```

12 00 + 00 00 00 00 00
12 01 + 00 00 00 00 00
12 02 + 00 00 00 00 00

12 03	+	00 00 00 00 00	59	CODE:	CON	0
12 04	+	00 00 00 00 00	60	TEXT:	CON	0
12 05	+	00 00 00 00 00	61	SYNT UNIT:	CON	0
12 06	+	00 00 00 00 00	62	LAST SYMBOL:	CON	0
12 07	+	00 00 00 00 00	63	TYPE:	CON	0
12 08	+	00 00 00 00 00	64	VALUE OF NUMBER:	CON	0
12 09	+	00 00 00 00 00	65	EPRCNECUS:	CON	0
12 10	+	00 00 00 00 00	66	AT EXPR:	CON	0
12 11	+	00 00 00 00 00	67	EXPR:	CON	0
12 12	+	00 00 00 00 00	68	VALUE:	CON	0
12 13	+	00 00 00 00 00	69	CONSTANT:	CON	0
12 14	+	00 00 00 00 00	70	ADDRESS:	CON	0
12 15	+	00 00 00 00 00	71	A PART:	CON	0
12 16	+	00 00 00 00 00	72	I PART:	CON	0
12 17	+	00 00 00 00 00	73	F PART:	CON	0
12 18	+	00 00 00 00 00	74	INSTRUCTION:	CON	0
12 19	+	00 00 00 00 00	75	LC:	CON	0
12 20	+	00 00 00 00 00	76	FUTURE:	CON	0
12 21	+	00 00 00 00 00	77	DANGEROUS:	CON	0
12 22	+	00 00 00 00 00	78	LAST LABEL:	CON	0
12 23	+	00 00 00 00 00	79	EQUIV:	CON	0
12 24	+	00 00 00 00 00	80	INST:	CON	0
12 25	+	00 00 00 00 00	81	LAST CONSTANT:	CON	0
12 26	+	00 00 00 00 00	82	LIT CONS PTR:	CON	0
12 27	+	00 00 00 00 00	83	START ADDRESS:	CON	0
			84			
			85			
			86		ORIG	* + 10
			87	BACKWARD:		" SPACE FOR ARRAY "B"
			88		ORIG	* + 10
			89	HERE:		" SPACE FOR ARRAYS 'M' AND 'F'
			90	INBUF X:	ORIG	* + 16
			91	INBUF:	ORIG	* + 16
			92	OUTBUF X:	ORIG	* + 24
			93	OUTBUF:	ORIG	* + 24
			94	TABLE:	ORIG	* + LENGTH OF TABLE
			95			" SPACE FOR SYMBOL TABLE
			96			
			97	" BEGIN OF PERMANENT PART OF ASSEMBLER		
			98			
			99			
28 28	+	00 00 00 14 99	100	TABLE PTR1:	CON	LENGTH OF TABLE - 1
28 29	+	00 00 00 00 00	101	LAST INST:	CON	0
28 30	+	00 00 00 00 00	102	ALF:	CON	0
28 31	+	00 00 00 00 00	103	CON:	CON	0
28 32	+	00 00 00 00 00	104	END:	CON	0
28 33	+	00 00 00 00 00	105	EQU:	CON	0
28 34	+	00 00 00 00 00	106	ORIG:	CON	0
28 35	+	00 00 00 00 10	107	BASE:	CON	10
			108	BYTESIZE: S:	EQU	1 (4 : 4)
28 36	+	00 00 00 00 99	109	MAX BYTE:	CON	S-1
			110	END MARKER:		
28 37	+	00 00 00 99 99	111	MAX ADDRESS:	CON	S-1 (4:4), S-1 (5:5)
28 38	+	99 99 99 99 99	112	MAX WORD:	CON	S-1 (1:1), S-1 (2:2), S-1 (3:3), S-1 (4:4), S-1 (5:5)
			113	LENGTH OF MEMORY:		
28 39	+	00 00 00 99 99	114		CON	3999
28 40	+	00 00 10 00 00	115	HUNDREDTHOUSAND:	CON	100000
28 41	+	66 66 66 66 66	116	SPACES:	ALF	.
28 42	+	66 66 66 94 94	117	JOB SEPARATOR:	ALF	. "
28 43	+	14 27 27 24 27	118	ERROR:	ALF	.ERROR


```

28 44 + 01 01 01 01 01 119 CORRECTION: CON 1(1:1), 1(2:2), 1(3:3), 1(4:4), 1(5:5)
120
121
122 " SUBROUTINES OF ASSEMBLER
123
124
125 READ CHAR: " RA IS AN OUTPUT PARAMETER FOR THE CODE OF THE CHARACTER,
126 " R1 AND R2 ARE PERMANENT READING POINTERS,
28 45 00 02 40 127 STJ 9F " SAVE RETURN ADDRESS
28 46 + 12 04 00 05 10 128 LDA TEXT " ) IF TEXT
28 47 + 12 03 00 05 01 129 ADD CODE " ) OR CODE
28 48 00 02 50 130 JAP PRINT LINE " ) THEN PRINT THIS LINE
28 49 00 05 52 131 J2NP READC " JUMP IF BUFFER NOT EXHAUSTED
28 50 + 12 79 00 05 10 132 LDA INBUF + 15 " POSSIBLE JOB SEPARATOR
28 51 + 28 42 00 45 70 133 CMPA JOB SEPARATOR (4 ; 5) " IS IT A JOB SEPARATOR ?
28 52 00 08 47 134 JNE 1F " ELSE CONTINUE THE READING
28 53 + 00 98 00 02 66 135 ENT6 98 " REPORT ERROR:
28 54 00 00 47 136 JMP ERRORM " END OF INPUT
28 55 00 00 47 137 JMP RESET TABLE " RESET THE INITIAL SYMBOL TABLE
28 56 00 00 47 138 JMP NEXT PROGRAM " TREAT THE NEXT PROGRAM
28 57 + 12 02 00 05 10 139 1H: LDA LINE NUMBER " ) INCREASE
28 58 + 28 57 00 08 47
28 59 + 00 01 00 00 60 140 INCA 1 " ) LINE NUMBER
28 60 + 12 02 00 05 30 141 STA LINE NUMBER " ) BY ONE
28 61 00 00 47 142 JMP BUFFER LN " ) PRINT THE LINE NUMBER
28 62 + 00 16 00 02 62 143 ENT2 16 " ONE AFTER LAST WORD OF BUFFER
28 63 + 28 62 00 16 42 144 JBUS * (READER) " WAIT FOR CARD READER
28 64 + 12 64 00 02 61 145 ENT1 INBUF " ) MOVE NEXT CARD FROM APPARATUS
28 65 + 12 48 00 16 07 146 MOVE INBUF X (16) " ) BUFFER TO INPUT BUFFER
28 66 + 12 79 00 05 10 147 LDA INBUF + 15 " POSSIBLE JOB SEPARATOR
28 67 + 28 42 00 45 70 148 CMPA JOB SEPARATOR (4 ; 5) " IS IT A JOB SEPARATOR ?
28 68 + 28 69 00 05 47 149 JE * + 2 " THEN SKIP NEXT INSTRUCTION
28 69 + 12 48 00 16 44 150 IN INBUF X (READER) " READ NEXT CARD
28 70 + 00 01 00 01 62 151 2H: DEC2 1 " PREVIOUS WORD OF BUFFER
28 71 00 00 52 152 J2N 3F " JUMP IF NOT AVAILABLE
28 72 + 12 64 02 05 20 153 LDAN INBUF, 2 " CONTENTS OF THE WORD (INVERTED)
28 73 + 28 69 00 05 47 154 CMPA SPACES (1 : 5) " ) IF SPACES
28 74 + 28 69 00 05 47 155 JE 2B " ) THEN LOOP
28 75 + 12 65 02 00 30 156 3H: STA INBUF + 1, 2 (0 : 0) " END OF BUFFER INDICATOR
28 76 + 28 74 00 00 52
28 77 - 00 16 00 02 62 157 ENT2 -16 " INITIALIZE READING PTR (TO WORDS)
28 78 + 00 55 03 03 63 158 READC: ENN3 5 : 5 /N ,3 " PREVIOUS CHARACTER IN BYTE 5 ?
28 79 + 28 76 00 05 52
28 80 00 02 53 159 J3P 3F " ELSE JUMP
28 81 + 00 55 00 02 63 160 ENT3 5 : 5 " INITIALIZE READING PTR (TO BYTES)
28 82 + 00 01 00 00 62 161 INC2 1 " NEXT WORD
28 83 00 05 52 162 J2NP 2F " JUMP IF WORD IS AVAILABLE
28 84 + 00 01 00 02 62 163 1H: ENT2 1 " POSITION = - 1
28 85 + 00 92 00 02 60 164 ENTA NLCR " END OF LINE READ
28 86 + 12 04 00 05 30 165 STA TEXT " TEXT OF ONE LINE COMPLETED
28 87 00 00 47 166 JMP 8F " READY
28 88 + 12 79 02 05 10 167 2H: LDA INBUF + 15, 2 " END OF BUFFER INDICATOR ?
28 89 + 28 85 00 05 52
28 90 + 28 81 00 00 50 168 JAN 1B " THEN SKIP REMAINING SPACES
28 91 - 00 66 03 03 63 169 3H: ENN3 6 : 6 /N ,3 " POINTER TO NEXT BYTE
28 92 + 28 87 00 02 53
28 93 + 28 89 00 44 33 170 ST3 * + 1 (4 : 4) " MODIFY NEXT INSTRUCTION
28 94 + 12 79 02 00 10 171 LDA INBUF + 15 ,2 (MOD) " ) READ THE CHARACTER
28 95 + 28 91 00 44 33 172 ST3 * + 1 (4 : 4) " MODIFY NEXT INSTRUCTION
28 96 + 13 27 02 00 30 173 STA OUTBUF + 8 + 15 ,2 (MOD) " ) PRINT THE CHARACTER

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28 92 + 12 00 00 05 30 174 8H: STA CHAR " INTERNAL CODE OF CHARACTER
*28 84 + 28 92 00 00 47
28 93 + 00 00 00 00 47 175 9H: JMP MOD " RETURN
*28 95 + 28 93 00 02 40
176
177
178 READ 5 CHAR: " RA IS AN OUTPUT PARAMETER FOR THE CODE OF THE CHARACTERS
28 94 00 02 40 179 STJ 9F " SAVE RETURN ADDRESS
28 95 + 12 05 00 05 10 180 LDA SYNT UNIT " ) IS SYNTACTICAL UNIT
28 96 + 00 61 00 01 60 181 DECA POINT " ) A POINT ?
28 97 00 04 50 182 JANZ 8F " ) ELSE ERROR
28 98 + 00 55 00 02 65 183 ENT5 5 : 5 " INITIALIZE BYTE POINTER
28 99 + 28 45 00 00 47 184 JMP READ CHAR " READ THE FIRST CHARACTER
29 00 - 00 66 05 03 65 185 2H: ENNS 6 : 6 /N ,5 " POINTER TO NEXT BYTE
29 01 + 00 92 00 02 67 186 ENTX NLCR " ) REPLACE POSSIBLE
29 02 + 12 00 00 05 77 187 CMPX CHAR " ) END OF CARD (NLCK)
29 03 + 29 05 00 08 47 188 JNE * + 2
29 04 + 00 66 00 02 60 189 ENTA SPACE " ) BY SPACES
29 05 + 29 06 00 44 35 190 ST5 * + 1 (4 : 4) " MODIFY NEXT INSTRUCTION
29 06 + 12 01 00 00 30 191 STA 5 CHAR (MOD) " STORE THE CHARACTER
29 07 - 00 55 05 03 65 192 ENNS 5 : 5 /N ,5 " IS IT THE FIFTH CHARACTER ?
29 08 + 28 45 00 08 47 193 JNE READ CHAR " READ NEXT CHARACTER
29 09 + 29 00 00 02 55 194 J5P 2B " ELSE CONTINUE THE READING
29 10 00 00 47 195 JMP READ SYNT UNIT " NEXT SYNTACTICAL UNIT
29 11 + 12 01 00 05 10 196 LDA 5 CHAR " RESULTING VALUE OF CHARACTERS
29 12 + 00 00 00 00 47 197 9H: JMP MOD " RETURN
29 13 + 00 10 00 02 66 198 8H: ENT6 10 " REPORT ERROR:
*29 14 + 29 13 00 04 50
29 14 00 00 47 199 JMP ERRORM " , REQUIRED
29 15 + 29 12 00 00 47 200 JMP 9B " RETURN
201
202
203 PRINT LINE: " NO PARAMETERS
29 15 00 02 40 204 STJ 9F " SAVE RETURN ADDRESS
*28 48 + 29 16 00 02 50
29 17 + 29 17 00 18 42 205 JBUS * (PRINTER) " WAIT FOR LINE PRINTER
29 18 + 12 80 00 02 61 206 ENT1 OUTBUF X " ) MOVE NEXT LINE FROM OUTPUT
29 19 + 13 04 00 24 07 207 MOVE OUTBUF (24) " ) BUFFER TO APPARATUS BUFFER
29 20 + 12 80 00 18 45 208 OUT OUTBUF X (PRINTER) " PRINT THE LINE
29 21 + 28 41 00 01 07 209 MOVE SPACES (1) " ) REFILL BUFFER
29 22 + 13 04 00 23 07 210 MOVE OUTBUF (23) " ) WITH SPACES
29 23 + 12 03 00 05 41 211 STZ CODE " NO CODE IN BUFFER
29 24 + 12 04 00 05 41 212 STZ TEXT " NO TEXT IN BUFFER
29 25 + 00 00 00 00 47 213 9H: JMP MOD " RETURN
*29 16 + 29 25 00 02 40
214
215
216 ERRORM: " R16 CONTAINS THE ERROR NUMBER
29 26 00 02 40 217 STJ 9F " SAVE RETURN ADDRESS
*29 14 + 29 26 00 00 47
*28 54 + 29 26 00 00 47
29 27 + 29 16 00 00 47 218 JMP PRINT LINE " PRINT THE (INCOMPLETE) LINE
29 28 + 28 43 00 05 10 219 LDA ERROR " ) 'PRINT' THE
29 29 + 13 05 00 05 30 220 STA OUTBUF + 1 " ) CHARACTERS: E R R O R
29 30 + 00 00 05 02 60 221 ENTA 0 ,6 " ERROR NUMBER
29 31 + 00 00 00 01 05 222 CHAR " ) 'PRINT' IT
29 32 + 13 06 00 45 37 223 STX OUTBUF + 2 (4 : 5)
29 33 + 12 05 00 05 10 224 LDA SYNT UNIT " SYNTACTICAL UNIT
29 34 + 00 00 00 01 05 225 CHAR " ) 'PRINT' IT

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29 35	+	13 08 00 05 37	226	STX	OUTBUF + 4 (3 : 5)	
29 36	+	12 11 00 05 10	227	LDA	EXPR	" VALUE OF EXPRESSION
29 37	+	00 00 00 01 05	228	CHAR		
29 38	+	13 10 00 05 30	229	STA	OUTBUF + 6	") 'PRINT'
29 39	+	13 11 00 05 37	230	STX	OUTBUF + 7	") THE VALUE
29 40	+	29 16 00 00 47	231	JMP	PRINT LINE	" PRINT THIS LINE
29 41	+	00 01 00 02 66	232	ENT6	1	") NOTE
29 42	+	12 09 00 05 36	233	ST6	ERRONEOUS	") THE ERROR
29 43	+	00 00 00 00 47	234	JMP	MOD	" RETURN
*29 26	+	29 43 00 02 40				
			235			
			236			
			237	READ SYNT UNIT:	" RA IS AN OUTPUT PARAMETER FOR THE CODE OF THE SYNTACTICAL UNIT	
			238		" R11 DELIVERS A POINTER TO A SYMBOL READ	
			239	STJ	9F	" SAVE RETURN ADDRESS
29 44		00 02 40				
*29 10	+	29 44 00 00 47				
29 45	+	12 00 00 05 10	240	LDA	CHAR	" LAST CHARACTER
29 46	+	00 66 00 01 60	241	DECA	SPACE	" IS IT A SPACE ?
29 47		00 04 50	242	JANZ	LETGIT	" ELSE JUMP
29 48	+	28 45 00 00 47	243	JMP	READ CHAR	") SKIP
29 49	+	29 46 00 00 47	244	JMP	1B	") SPACES
29 50	-	00 31 00 01 60	245	DECA	LETGIT:	LETTER Z = SPACE
29 47	+	29 50 00 04 50				
29 51		00 02 50	246	JAP	SEPARATOR	" JUMP IF NOT A LETTER OR DIGIT
29 52	-	00 26 00 01 60	247	DECA	9 - LETTER Z	" IS IT A LETTER ?
29 53		00 02 30	248	STA	COMPL (0 : 2)	" NOTE THIS (MODIFY INSTRUCTION)
29 54	+	28 28 00 05 14	249	LD4	TABLE PTR1	") FREE SPACE
29 55	+	00 03 00 01 64	250	DECA	3	") IN SYMBOL TABLE
29 56	+	12 26 00 05 16	251	LD6	LIT CONS PTR	" END OF SPACE FOR SYMBOLS
29 57	+	00 00 04 01 66	252	DEC6	0 , 4	" SPACE AVAILABLE ?
29 58		00 02 56	253	J6P	EXIT	" ELSE EXIT
29 59	+	13 28 04 05 41	254	STZ	TABLE , 4	" CLEAN SPACE FOR CHARACTERS
29 60	+	00 10 00 00 60	255	INCA	9 + 1	" MODIFY THE CHARACTER (ADD 1)
29 61	+	13 28 04 11 30	256	STA	TABLE , 4 (1 : 1)	" STORE THE FIRST CHARACTER
29 62	+	00 11 00 02 65	257	ENT5	1 : 1	" INITIALIZE BYTE POINTER
29 63	+	28 45 00 00 47	258	JMP	READ CHAR	" READ NEXT CHARACTER
29 64	+	00 66 00 01 60	259	DECA	SPACE	" IS IT A SPACE ?
29 65	+	28 45 00 01 50	260	JAZ	READ CHAR	" THEN SKIP ONE
29 66	+	12 00 00 05 10	261	LDA	CHAR	" CHARACTER
29 67	+	00 35 00 01 60	262	DECA	LETTER Z	" IS IT A LETTER OR A DIGIT ?
29 68		00 05 50	263	JANP	2F	" THEN JUMP
29 69	+	00 01 00 03 61	264	ENN1	1	") INDICATOR FOR END OF
29 70	+	13 28 04 00 31	265	ST1	TABLE , 4 (0 : 0)	") ALPHA NUMERICAL INFORMATION
29 71		00 00 47	266	JMP	COMPL	" SYNTACTICAL UNIT COMPLETED
29 72	-	00 26 00 01 60	267	DECA	9 - LETTER Z	" IS THE CHARACTER A LETTER ?
*29 68	+	29 72 00 05 50				
29 73	+	29 75 00 05 50	268	JANP	* + 2	" ELSE JUMP
29 74		00 02 34	269	ST4	COMPL (0 : 2)	" NOTE OCCURRENCE OF LETTER
29 75	-	00 55 05 03 65	270	ENN5	5 : 5 /N , 5	" PREVIOUS CHARACTER IN BYTE 5 ?
29 76		00 02 55	271	J5P	3F	" ELSE JUMP
29 77	+	00 55 00 02 65	272	ENT5	5 : 5	" INITIALIZE BYTE POINTER
29 78	+	00 01 00 01 64	273	DEC4	1	" NEXT WORD WITH ALPHA NUMERICALS
29 79	+	12 26 00 05 16	274	LD6	LIT CONS PTR	" END OF SPACE FOR SYMBOLS
29 80	+	00 00 04 01 66	275	DEC6	0 , 4	" SPACE AVAILABLE ?
29 81		00 02 56	276	J6P	EXIT	" ELSE EXIT
29 82	+	13 28 04 05 41	277	STZ	TABLE , 4	" CLEAN SPACE FOR CHARACTERS
29 83	-	00 66 05 03 65	278	ENN5	6 : 6 /N , 5	" POINTER TO NEXT BYTE
*29 76	+	29 83 00 02 55				
29 84	+	00 10 00 00 60	279	INCA	9 + 1	" MODIFY THE CHARACTER (ADD 1)
29 85	+	29 86 00 44 35	280	ST5	* + 1 (4 : 4)	" MODIFY NEXT INSTRUCTION


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29 86 + 13 28 04 00 30 281
29 87 + 29 63 00 00 47 282
29 88 + 00 00 00 02 60 283
*29 74 + 29 88 00 02 34
*29 71 + 29 88 00 00 47
*29 53 + 29 88 00 02 30
29 89 + 28 28 00 05 11 284
29 90 + 12 08 00 05 41 285
29 91 + 00 00 04 02 66 286
29 92 - 00 04 01 01 66 287
29 93 00 02 56 288
29 94 00 01 56 289
29 95 + 13 24 01 05 10 290
29 96 + 28 44 00 05 02 291
29 97 + 00 05 00 03 06 292
29 98 + 13 25 01 05 10 293
30 00 + 28 44 00 05 02 294
30 01 + 00 00 00 00 05 295
30 02 + 12 08 00 05 17 296
30 03 00 02 57 297
30 04 + 12 08 00 05 30 298
30 05 + 00 02 00 01 61 300
30 06 + 29 92 00 00 47 301
30 07 + 12 08 00 05 10 302
*29 95 + 30 07 00 01 56
30 08 + 28 40 00 05 03 303
30 09 00 04 50 304
30 10 + 12 08 00 05 37 305
30 11 + 13 29 04 05 10 306
30 12 + 28 44 00 05 02 307
30 13 + 00 05 00 03 06 308
30 14 + 00 00 00 00 05 309
30 15 + 12 08 00 05 01 310
30 16 + 12 08 00 05 30 311
30 17 + 00 00 00 02 60 312
*29 94 + 30 17 00 02 56
30 18 - 00 55 05 03 67 313
30 19 00 05 04 314
30 20 + 30 21 00 02 30 315
30 21 + 00 00 00 02 66 316
30 22 + 28 40 00 05 10 317
30 23 + 00 00 00 01 05 318
30 24 + 00 00 06 03 06 319
30 25 + 00 00 00 00 05 320
30 26 00 05 30 321
30 27 + 12 08 00 05 10 322
30 28 00 05 03 323
30 29 00 04 50 324
30 30 + 12 08 00 05 37 325
30 31 + 13 28 04 05 17 326
30 32 + 28 44 00 05 10 327
30 33 - 00 05 06 03 66 328
30 34 + 00 00 06 02 06 329
30 35 + 28 44 00 05 02 330
30 36 + 00 05 00 03 06 331
30 37 + 00 00 00 00 05 332
30 38 + 12 08 00 05 01 333
30 39 + 12 08 00 05 30 334
30 40 + 01 02 00 02 60 335

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COMPL:

1H:

2H:

3H:

1H:

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STA TABLE ,4 (MOD) " STORE THE LETTER OR DIGIT
JMP READ " PROCEED WITH NEXT CHARACTERS
ENTA MOD " AT LEAST ONE LETTER READ ?

JAP CONTAINS LETTER " THEN JUMP, AND TREAT SYMBOL
LD1 TABLE PTR1 " POINTER TO NEW INFO CELL
STZ VALUE OF NUMBER " INITIALIZE VALUE OF NUMBER
ENT6 0 ,4 " END OF INFO JUST STORED
DEC6 - 4 ,1 " LENGTH OF INFO
J6P 3F " JUMP IF ONLY ONE WORD WITH DIGITS
J6Z 2F " JUMP IF ONLY TWO WORDS WITH DIGITS
LDA TABLE - 4 ,1 " FIFTH TILL TENTH DIGIT
SUB CORRECTION " CANCEL MODIFICATION OF CHARACTERS
SRAX 5 " SHIFT DIGITS INTO RX
LDA TABLE - 3 ,1 " FIRST TILL FIFTH DIGIT
SUB CORRECTION " CANCEL MODIFICATION OF CHARACTERS
NUM " NUMERICAL VALUE OF TEN DIGITS
LOX VALUE OF NUMBER " MORE DIGITS TREATED ALREADY ?
JXP OVERFLOW " THEN OVERFLOW NOW
STA VALUE OF NUMBER " STORE VALUE OF THE TEN DIGITS
DEC1 2 " POINTER TO NEXT DIGITS
JMP 1B " CONTINUE THE PROCESS
LDA VALUE OF NUMBER " ) SHIFT VALUE OBTAINED

MUL HUNDREDTHOUSAND " ) OVER FIVE DECIMAL POSITIONS
JANZ OVERFLOW " JUMP IF OVERFLOW OCCURS
STX VALUE OF NUMBER " STORE INTERMEDIATE VALUE
LDA TABLE + 1 ,4 " FIVE LAST BUT FIVE DIGITS
SUB CORRECTION " CANCEL MODIFICATION OF CHARACTERS
SRAX 5 " SHIFT DIGITS INTO RX, RA = 0
NUM " NUMERICAL VALUE OF FIVE DIGITS
ADD VALUE OF NUMBER " ADD TO INTERMEDIATE VALUE
STA VALUE OF NUMBER " RESULTING VALUE (NO OVERFLOW)
ENTA 0 " RA = 0

ENNX 5 : 5 /N ,5 " ) NUMBER OF BYTES IN LAST WORD
DIV = 1 : 1 = " ) WHICH HAVE NOT BEEN FILLED IN
STA * + 1 (0 : 2) " MODIFY NEXT INSTRUCTION
ENT6 MOD " NUMBER OF BYTES NOT FILLED IN
LDA HUNDREDTHOUSAND " ) ONE, FOLLOWED BY
CHAR " ) FIVE DECIMAL ZERUS
SRAX 0 ,6 " SHIFT OUT THE BYTES NOT FILLED IN
NUM " ) NUMERICAL VALUE INDICATING
STA AUX " ) THE DECIMAL SHIFT OF VALUE OF NUMBER
LDA VALUE OF NUMBER " VALUE OF NUMBER
MUL AUX " SHIFT OVER DECIMAL POSITIONS
JANZ OVERFLOW " JUMP IF OVERFLOW OCCURS
STX VALUE OF NUMBER " STORE INTERMEDIATE VALUE
LDX TABLE ,4 " LAST DIGITS IN RX
LDA CORRECTION " (MODIFIED) ZEROS IN RA
ENNA - 5 ,6 " NUMBER OF DIGITS IN LAST WORD
SLAX 0 ,6 " SHIFT DIGITS INTO RA
SUB CORRECTION " CANCEL MODIFICATION OF CHARACTERS
SRAX 5 " SHIFT DIGITS INTO RX, RA = 0
NUM " NUMERICAL VALUE OF LAST DIGITS
ADD VALUE OF NUMBER " ADD TO INTERMEDIATE VALUE
STA VALUE OF NUMBER " RESULTING VALUE (NO OVERFLOW)
ENTA NUMBER " ) THIS SYNTACTICAL UNIT

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30 41	+	12 05 00 05 30	336		STA	SYNT UNIT	") IS A NUMBER
30 42		00 00 00 47	337		JMP	9F	" RETURN
30 43	+	12 08 00 05 41	338	OVERFLOW:	STZ	VALUE OF NUMBER	" RESULTING VALUE
*30 29	+	30 43 00 04 50					
*30 09	+	30 43 00 04 50					
*30 03	+	30 43 00 02 57					
30 44	+	00 11 00 02 66	339		ENT6	11	" REPORT ERROR:
30 45	+	29 26 00 00 47	340		JMP	ERRORM	" VALUE OF NUMBER TOO LARGE
30 46	+	30 40 00 00 47	341		JMP	1B	" READY
			342	CONTAINS LETTER:			
30 47	+	28 28 00 05 11	343		LD1	TABLE PTR1	" POINTER TO NEW INFO CELL
*29 89	+	30 47 00 02 50					
30 48	+	00 03 04 01 61	344		DEC1	3 , 4	" ONLY ONE WORD WITH CHARACTERS ?
30 49		00 04 51	345		J1NZ	LOOK FOR SYMBOL	" ELSE ORDINARY SYMBOL
30 50	+	00 22 00 01 65	346		DEC5	2 : 2	" TWO CHARACTERS ?
30 51		00 04 55	347		J5NZ	LOOK FOR SYMBOL	" ELSE ORDINARY SYMBOL
30 52	+	13 28 04 11 11	348		LD1	TABLE , 4 (1 : 1)	" FIRST CHARACTER
30 53	+	00 10 00 01 61	349		DEC1	9 + 1	" IS IT A DIGIT ?
30 54		00 02 51	350		J1P	LOOK FOR SYMBOL	" ELSE ORDINARY SYMBOL
30 55	+	13 28 04 22 17	351		LDX	TABLE , 4 (2 : 2)	" SECOND CHARACTER
30 56	+	00 12 00 01 67	352		DECX	LETTER B + 1	" IS IT LETTER B ?
30 57		00 01 57	353		JXZ	B	" THEN JUMP
30 58	+	00 04 00 01 67	354		DECX	LETTER F - LETTER B	
30 59		00 01 57	355		JXZ	F	" JUMP IF IT IS LETTER F
30 60	+	00 02 00 01 67	356		DECX	LETTER H - LETTER F	
30 61		00 04 57	357		JXNZ	LOOK FOR SYMBOL	" JUMP IF IT IS NOT LETTER H
30 62	+	00 02 00 02 67	358		ENTX	LOCAL LABEL	" TYPE LOCAL LABEL
30 63		00 00 47	359		JMP	1F	" READY
30 64	+	00 04 00 02 67	360	F:	ENTX	LOCAL FORWARD	" TYPE LOCAL FORWARD
*30 59	+	30 64 00 01 57					
30 65		00 00 47	361		JMP	1F	" READY
30 66	+	00 03 00 02 67	362	B:	ENTX	LOCAL BACKWARD	" TYPE LOCAL BACKWARD
30 67	+	30 66 00 01 57					
30 67	+	12 07 00 05 37	363	1H:	STX	TYPE	" STORE THE TYPE
*30 65	+	30 67 00 00 47					
*30 65	+	30 67 00 00 47					
30 68	+	00 10 01 03 61	364		ENN1	9 + 1 , 1	" POINTER (NEGATIVE-1) TO LOCAL SYMBOL
30 69		00 00 47	365		JMP	7F	" LOCAL SYMBOL FOUND
			366	LOOK FOR SYMBOL:			
30 70	+	14 99 00 02 61	367		ENT1	LENGTH OF TABLE = 1	" ROOT OF BINARY TREE
*30 61	+	30 70 00 04 57					
*30 54	+	30 70 00 02 51					
*30 51	+	30 70 00 04 55					
*30 49	+	30 70 00 04 51					
30 71	-	00 03 01 02 65	368	NEXT:	ENT5	- 3 , 1	" POINTER TO LETTERS IN TREE
30 72	+	28 28 00 05 16	369		LD6	TABLE PTR1	") POINTER TO LETTERS OF SYMBOL
30 73	+	00 03 00 01 66	370		DEC6	3	") TO BE LOOKED FOR
30 74	+	13 28 06 05 10	371	1H:	LDA	TABLE , 6	" COMPARE THE NEW SYMBOL
30 75	+	13 28 05 25 70	372		CMPA	TABLE , 5 (1 : 5)	" WITH THE SYMBOL IN TREE
30 76		00 08 47	373		JNE	3F	" SELECT NEXT SYMBOL , IF NOT EQUAL
30 77	+	13 28 05 05 17	374		LOX	TABLE , 5	") JUMP IF THESE ARE THE FINAL
30 78		00 05 57	375		JXNP	2F	") CHARACTERS OF THE SYMBOL IN TREE
30 79		00 05 50	376		JANP	LEFT	" JUMP IF NEW SYMBOL EXHAUSTED
30 80	+	00 01 00 01 65	377		DEC5	1	") NEXT
30 81	+	00 01 00 01 66	378		DEC6	1	") FIVE LETTERS
30 82	+	30 74 00 00 47	379		JMP	1B	") TO BE COMPARED
30 83		00 05 50	380	2H:	JANP	FOUND	" JUMP IF NEW SYMBOL ALSO EXHAUSTED
*30 76	+	30 83 00 05 57					
30 84		00 00 47	381		JMP	RIGHT	" NEW SYMBOL IS LONGER
30 85		00 06 47	382	3H:	JG	RIGHT	" JUMP IF NEW SYMBOL SUCCEEDS OLD ONE

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*30 76	+	30 85 00 08 47					
30 85	+	13 28 01 23 10	383	LEFT:	LDA	TABLE ,1 (2 : 3)	" LEFT BRANCH
*30 79	+	30 86 00 05 50					
30 87		00 04 50	384		JANZ	1F	" JUMP IF NOT EMPTY
30 88	+	00 00 01 03 67	385		ENNX	0 ,1	" NEW THREAD
30 89	+	28 28 00 05 10	386		LDA	TABLE PTR1	" NEW LEFT BRANCH
30 90	+	13 28 01 23 30	387		STA	TABLE ,1 (2 : 3)	" STORE NEW LEFT BRANCH
30 91		00 00 47	388		JMP	NEW	" CREATE NEW INFO CELL
30 92	+	30 93 00 02 30	389	1H:	STA	* + 1 (0 : 2)	" MODIFY NEXT INSTRUCTION
*30 87	+	30 92 00 04 50					
30 93	+	00 00 00 02 61	390		ENT1	MOD	" POINTER TO NEXT INFO CELL
30 94	+	30 71 00 00 47	391		JMP	NEXT	" CONTINUE THE SEARCH
30 95	+	13 28 01 05 10	392	RIGHT:	LDA	TABLE ,1	" RIGHT BRANCH
*30 85	+	30 95 00 06 47					
*30 84	+	30 95 00 00 47					
30 95	+	30 92 00 02 50	393		JAP	1B	" JUMP IF IT IS NOT A THREAD
30 97	+	13 28 01 45 27	394		LDXN	TABLE ,1 (4 : 5)	" TAKE THE THREAD (NEGATIVE)
30 98	+	28 28 00 05 10	395		LDA	TABLE PTR1	" NEW RIGHT BRANCH
30 99	+	13 28 01 00 30	396		STA	TABLE ,1 (0 : 0)	" NOT A THREAD
31 00	+	13 28 01 45 30	397		STA	TABLE ,1 (4 : 5)	" STORE RIGHT BRANCH
31 01	+	28 28 00 05 11	398	NEW:	LD1	TABLE PTR1	" POINTER TO NEW INFO CELL
*30 91	+	31 01 00 00 47					
31 02	+	13 28 01 05 37	399		STX	TABLE ,1	" STORE A THREAD
31 03	+	13 27 01 05 41	400		STZ	TABLE - 1 ,1	" INITIALIZE
31 04	+	28 37 00 05 10	401		LDA	END MARKER	" NEW
31 05	+	13 26 01 05 30	402		STA	TABLE - 2 ,1	" INFO CELL
31 06	+	00 01 00 01 64	403		DEC4	1	
31 07	+	28 28 00 05 34	404		STA	TABLE PTR1	" INCREASE TABLE POINTER
31 08		00 00 47	405		JMP	6F	" ORDINARY SYMBOL STORED
31 09	+	28 28 00 05 10	406	FOUND:	LDA	TABLE PTR1	" POINTER TO NEW SYMBOL
*30 86	+	31 09 00 05 50					
31 10	+	14 99 00 01 60	407		DECA	LENGTH OF TABLE + 1	
31 11	+	28 37 00 05 27	408		LDXN	END MARKER	" FINAL THREAD
31 12	+	31 01 00 01 50	409		JAZ	NEW	" JUMP IF ONLY ONE SYMBOL IN TREE
31 13	+	13 27 01 01 10	410	6H:	LDA	TABLE - 1 ,1 (0 : 1)	
*31 08	+	31 13 00 00 47					
31 14	+	12 07 00 05 30	411		STA	TYPE	" SAVE TYPE OF SYMBOL
31 15	+	12 06 00 05 31	412	7H:	ST1	LAST SYMBOL	" POINTER TO THIS SYMBOL
*30 89	+	31 15 00 00 47					
31 16	+	01 03 00 02 60	413		ENTA	SYMBOL	" THIS SYNTACTICAL UNIT
31 17	+	12 05 00 05 30	414		STA	SYNT UNIT	" IS A SYMBOL
31 18		00 00 47	415		JMP	9F	" RETURN
31 19	+	00 59 00 01 60	416	SEPARATOR:	DECA	QUOTE - LETTER Z	" IS THE CHARACTER A QUOTE ?
*29 01	+	31 19 00 02 50					
31 20		00 04 50	417		JANZ	2F	" ELSE JUMP
31 21	+	28 45 00 00 47	418	1H:	JMP	READ CHAR	" READ NEXT CHARACTER
31 22	+	00 92 00 01 60	419		DECA	NLCR	" IS IT A NLCR ?
31 23	+	31 21 00 04 50	420		JANZ	1B	" ELSE SKIP UNTIL NLCR
31 24		00 00 47	421		JMP	3F	" NLCR READ
31 25	-	00 02 00 01 60	422	2H:	DECA	NLCR - QUOTE	" IS IT A NLCR ?
*31 20	+	31 25 00 04 50					
31 26		00 04 50	423		JANZ	4F	" ELSE JUMP
31 27	+	01 01 00 02 60	424	3H:	ENTA	SEP	" THIS SYNTACTICAL UNIT
*31 24	+	31 27 00 00 47					
31 28		00 00 47	425		JMP	8F	" IS A SEPARATOR
31 29	-	00 28 00 01 60	426	4H:	DECA	SEMICOLON - NLCR	" IS THE CHARACTER A SEMICOLON ?
*31 25	+	31 29 00 04 50					
31 30		00 04 50	427		JANZ	01V	" ELSE JUMP
31 31	+	00 01 00 02 61	428		ENT1	1	" TEXT OF ONE INSTRUCTION
31 32	+	12 04 00 05 31	429		ST1	TEXT	" COMPLETED NOW

31 33	+	31 27 00 00 47	430		JMP	3B	" SEPARATOR READ
31 34	-	00 24 00 01 60	431	DIV:	DECA	DIV1 - SEMICOLON	" IS THE CHARACTER A DIVISION SIGN ?
*31 35	+	31 34 00 04 50					
31 35		00 04 50	432		JANZ	OTHER	" ELSE JUMP
31 36	+	28 45 00 00 47	433		JMP	READ CHAR	" READ NEXT CHARACTER
31 37	+	00 40 00 01 60	434		DECA	DIV1	" IS IT ALSO A DIVISION SIGN ?
31 38		00 04 50	435		JANZ	5F	" ELSE JUMP
31 38	+	01 04 00 02 60	436		ENTA	DIV2	") THIS SYNTACTICAL UNIT
31 40		00 00 47	437		JMP	8F	") IS A SPECIAL DIVISION OPERATOR
31 41	+	00 40 00 02 60	438	5H:	ENTA	DIV1	" COMMON DIVISION OPERATOR
*31 38	+	31 41 00 04 50					
31 42	+	12 05 00 05 30	439		STA	SYNT UNIT	" THIS SYNTACTICAL UNIT
31 43		00 00 47	440		JMP	9F	" RETURN
31 44	+	00 40 00 00 60	441	OTHER:	INCA	DIV1	" OTHER SINGLE CHARACTER
*31 35	+	31 44 00 04 50					
31 45	+	12 05 00 05 30	442	8H:	STA	SYNT UNIT	" THIS SYNTACTICAL UNIT
*31 40	+	31 45 00 00 47					
*31 28	+	31 45 00 00 47					
31 46	+	00 66 00 02 67	443		ENTX	SPACE	" SIMULATE A SPACE
31 47	+	12 00 00 05 37	444		STX	CHAR	" INSERT IT
31 48	+	00 00 00 00 47	445	9H:	JMP	MOD	" RETURN
*31 43	+	31 48 00 00 47					
*31 18	+	31 48 00 00 47					
*31 42	+	31 48 00 00 47					
*29 44	+	31 48 00 02 40					
			446				
			447				
			448	READ AT EXPRESSION:		" RA IS AN OUTPUT PARAMETER FOR THE VALUE OF THE AT EXPR	
31 49		00 02 40	449		STJ	9F	" SAVE RETURN ADDRESS
31 50	+	29 44 00 00 47	450		JMP	READ SYNT UNIT	" READ ONE SYNTACTICAL UNIT
31 51	+	31 54 00 00 47	451		JMP	* + 3	
31 52		00 02 40	452	AT EXPRESSION:	STJ	9F	" SAVE RETURN ADDRESS
31 53	+	12 05 00 05 10	453		LDA	SYNT UNIT	" LAST SYNTACTICAL UNIT
31 54	+	01 03 00 01 60	454		DECA	SYMBOL	" IS IT A SYMBOL ?
31 55		00 01 50	455		JAZ	SYMB	" THEN JUMP
31 56	-	00 01 00 01 60	456		DECA	NUMBER - SYMBOL	" OR IS IT A NUMBER ?
31 57		00 01 50	457		JAZ	NUMB	" THEN JUMP
31 58	-	00 63 00 01 60	458		DECA	ASTERISK - NUMBER	
31 59		00 01 50	459		JAZ	ASTSK	" JUMP IF IT IS AN ASTERISK
31 60	+	00 21 00 02 66	460		ENT6	21	" REPORT ERROR:
31 61	+	29 26 00 00 47	461	1H:	JMP	ERRORM	" ATOMIC EXPRESSION REQUIRED
31 62	+	00 00 00 02 60	462		ENTA	0	" RESULTING VALUE
31 63		00 00 47	463		JMP	8F	" READY
31 64	+	12 06 00 05 11	464	SYMB:	LD1	LAST SYMBOL	" POINTER TO INFO CELL
*31 55	+	31 64 00 01 50					
31 65	+	12 07 00 05 10	465		LDA	TYPE	" TYPE OF SYMBOL
31 66	+	00 01 00 01 60	466		DECA	DEFINED SYMB	" IS IT A DEFINED SYMBOL ?
31 67		00 01 50	467		JAZ	DEFINED	" THEN JUMP
31 68	+	00 02 00 01 60	468		DECA	LOCAL BACKWARD - DEFINED SYMB	
31 69		00 01 50	469		JAZ	LOC DEFINED	" JUMP IF IT IS A LOCAL BACKWARD
31 70	+	00 23 00 02 66	470		ENT6	23	" REPORT ERROR:
31 71	+	31 61 00 00 47	471		JMP	1B	" UNDEFINED SYMBOL IN EXPRESSION
31 72	+	13 26 01 05 10	472	DEFINED:	LDA	TABLE - 2 ,1	" VALUE OF SYMBOL
*31 67	+	31 72 00 01 50					
31 73		00 00 47	473		JMP	8F	" READY
31 74	+	12 48 01 03 10	474	LOC DEFINED:	LDA	HERE ,1 (3 : 3)	" CORRESPONDING LABEL OCCURRED ?
*31 69	+	31 74 00 01 50					
31 75		00 02 50	475		JAP	2F	" THEN JUMP
31 76	+	00 22 00 02 66	476		ENT6	22	" REPORT ERROR:
31 77	+	31 61 00 00 47	477		JMP	1B	" UNDEF. LOCAL BACKWARD IN EXPRESSION

31 75	+	12 38 01 05 10	478	2H:	LDA	BACKWARD ,1	" VALUE OF LOCAL BACKWARD
*31 75	+	31 78 00 02 50					
31 77	+	00 00 47	479		JMP	8F	" READY
31 83	+	12 08 00 05 10	480	NUMB:	LDA	VALUE OF NUMBER	" VALUE OF NUMBER
*31 87	+	31 80 00 01 50					
31 81	+	00 00 47	481		JMP	8F	" READY
31 82	+	12 19 00 05 10	482	ASTSK:	LDA	LC	" VALUE OF LOCATION COUNTER
*31 89	+	31 82 00 01 50					
31 83	+	12 10 00 05 30	483	8H:	STA	AT EXPR	" VALUE OF ATOMIC EXPRESSION
*31 81	+	31 83 00 00 47					
*31 79	+	31 83 00 00 47					
*31 76	+	31 83 00 00 47					
*31 63	+	31 83 00 00 47					
31 84	+	00 00 00 00 47	484	9H:	JMP	MOD	" RETURN
*31 82	+	31 84 00 02 40					
*31 49	+	31 84 00 02 40					
			485				
			486				
			487			READ EXPRESSION:" RA IS AN OUTPUT PARAMETER FOR THE VALUE OF THE EXPRESSION	
31 85		00 02 40	488		STJ	9F	" SAVE RETURN ADDRESS
31 86	+	12 11 00 05 41	489		STZ	EXPR	" INITIALIZE EXPR
31 87	+	12 05 00 05 10	490		LDA	SYNT UNIT	" LAST SYNTACTICAL UNIT
31 88	+	00 37 00 01 60	491		DECA	PLUS	" IS IT AN ADDING OPERATOR ?
31 89		00 01 50	492		JAZ	ADD	" THEN JUMP
31 90	+	00 01 00 01 60	493		DECA	MINUS - PLUS	" IS IT A SUBTRACTING OPERATOR ?
31 91		00 01 50	494		JAZ	SUB	" THEN JUMP
31 92	+	31 52 00 00 47	495		JMP	AT EXPRESSION	" EXPECT AN ATOMIC EXPRESSION
31 93	+	12 11 00 05 30	496	NEXT OPERAND:	STA	EXPR	" SAVE THE VALUE OF IT
31 94		00 02 47	497		JOV	ARITHM OV	" IF OVERFLOW THEN REPORT IT
31 95	+	29 44 00 00 47	498		JMP	READ SYNT UNIT	" READ NEXT SYNTACTICAL UNIT
31 96	+	00 37 00 01 60	499		DECA	PLUS	" IS IT AN ADDING OPERATOR ?
31 97		00 01 50	500		JAZ	ADD	" THEN JUMP
31 98	+	00 01 00 01 60	501		DECA	MINUS - PLUS	" IS IT A SUBTRACTING OPERATOR ?
31 99		00 01 50	502		JAZ	SUB	" THEN JUMP
32 00	+	00 01 00 01 60	503		DECA	TIMES - MINUS	" IS IT A MULTIPLICATION OPERATOR ?
32 01		00 01 50	504		JAZ	MUL	" THEN JUMP
32 02	+	00 01 00 01 60	505		DECA	DIV1 - TIMES	" IS IT A COMMON DIVISION OPERATOR ?
32 03		00 01 50	506		JAZ	DIV X	" THEN JUMP
32 04	+	00 64 00 01 60	507		DECA	DIV2 - DIV1	" IS IT A SPECIAL DIVISION OPERATOR ?
32 05		00 01 50	508		JAZ	DIV A	" THEN JUMP
32 06	-	00 41 00 01 60	509		DECA	DEC - DIV2	" IS IT A DECIMAL OPERATOR ?
32 07		00 01 50	510		JAZ	TIMES BASE PLUS	" THEN JUMP
32 08	+	12 11 00 05 10	511		LDA	EXPR	" RESULTING VALUE OF EXPRESSION
32 09	+	00 00 00 00 47	512	9H:	JMP	MOD	" RETURN
*31 82	+	32 09 00 02 40					
			513	TIMES BASE PLUS:			
32 10	+	12 11 00 15 10	514		LDA	EXPR (1 : 5)	" ABSOLUTE VALUE OF EXPRESSION
*32 07	+	32 10 00 01 50					
32 11	+	28 35 00 05 03	515		MUL	BASE	" MULTIPLY BY TEN (BASE)
32 12	+	12 11 00 15 37	516		STX	EXPR (1 : 5)	" RESTORE THE VALUE
32 13	+	28 38 00 05 01	517		ADD	MAX WORD	" FORCE OVERFLOW IF RA ≠ 0
32 14	+	31 49 00 00 47	518	ADD:	JMP	READ AT EXPRESSION	
*31 97	+	32 14 00 01 50					
*31 89	+	32 14 00 01 50					
32 15	+	12 11 00 05 01	519		ADD	EXPR	" EXECUTE ADDING OPERATION
32 16	+	31 93 00 00 47	520		JMP	NEXT OPERAND	" PROCEED WITH NEXT OPERANDS
32 17	+	31 49 00 00 47	521	SUB:	JMP	READ AT EXPRESSION	
*31 99	+	32 17 00 01 50					
*31 91	+	32 17 00 01 50					
32 18	+	12 11 00 05 10	522		LDA	EXPR	" EXECUTE

32 19	+	12 10 00 05 02	523		SUB	AT EXPR	") SUBTRACTING OPERATION
32 20	+	31 93 00 00 47	524		JMP	NEXT OPERAND	" PROCEED WITH NEXT OPERANDS
32 21	+	31 49 00 00 47	525	MUL:	JMP	READ AT EXPRESSION	
*32 01	+	32 21 00 01 50					
32 22	+	12 11 00 05 03	526		MUL	EXPR	") EXECUTE
32 23	+	12 11 00 05 37	527		STX	EXPR	") MULTIPLICATION OPERATION
32 24	+	31 95 00 01 50	528		JAZ	NEXT OPERAND + 2	" JUMP IF NO OVERFLOW
32 25	+	00 31 00 02 66	529	ARITHM OV:	ENT6	31	" REPORT ERROR:
*32 94	+	32 25 00 02 47					
32 26	+	29 26 00 00 47	530		JMP	ERRORM	" OVERFLOW BY ARITHMETIC OPERATION
32 27	+	00 00 00 02 60	531		ENTA	0	" RESULTING VALUE
32 28	+	31 93 00 00 47	532		JMP	NEXT OPERAND	" PROCEED WITH NEXT OPERANDS
32 29	+	31 49 00 00 47	533	DIV X:	JMP	READ AT EXPRESSION	
*32 03	+	32 29 00 01 50					
32 30	+	00 00 00 02 60	534		ENTA	0	" RA = 0
32 31	+	12 11 00 05 17	535		LDX	EXPR	" RX = EXPR
32 32	+	12 10 00 05 04	536		DIV	AT EXPR	" EXECUTE DIVISION OPERATION
32 33	+	31 93 00 00 47	537		JMP	NEXT OPERAND	" PROCEED WITH NEXT OPERANDS
32 34	+	31 49 00 00 47	538	DIV A:	JMP	READ AT EXPRESSION	
*32 05	+	32 34 00 01 50					
32 35	+	12 11 00 05 10	539		LDA	EXPR	" RA = EXPR
32 36	+	00 00 00 02 67	540		ENTX	0	" RX = 0
32 37	+	12 10 00 05 04	541		DIV	AT EXPR	" EXECUTE DIVISION OPERATION
32 38	+	31 93 00 00 47	542		JMP	NEXT OPERAND	" PROCEED WITH NEXT OPERANDS
			543				
			544				
			545	READ W VALUE:		" RA IS AN OUTPUT PARAMETER FOR THE VALUE OF THE WORD VALUE	
32 39		00 02 40	546		STJ	9F	" SAVE RETURN ADDRESS
32 40	+	12 13 00 05 41	547		STZ	CONSTANT	" INITIALIZE CONSTANT
32 41	+	32 43 00 00 47	548		JMP	* + 2	
32 42	+	29 44 00 00 47	549	NEXT SPEC:	JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT
32 43	+	31 85 00 00 47	550		JMP	READ EXPRESSION	" READ AN EXPRESSION
32 44	+	12 12 00 05 30	551		STA	VALUE	" VALUE OF EXPRESSION
32 45	+	12 05 00 05 10	552		LDA	SYNT UNIT	" NOT BELONGING TO EXPRESSION
32 46	+	00 71 00 01 60	553		DECA	PAREN	" IS IT A PAREN ?
32 47		00 04 50	554		JANZ	STANDARD	" ELSE STANDARD FIELD SPECIFICATION
32 48	+	29 44 00 00 47	555		JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT
32 49	+	31 85 00 00 47	556		JMP	READ EXPRESSION	" READ FIELD SPECIFICATION
32 50	+	00 05 00 03 06	557		SRAX	5	" RX = FIELD SPECIFICATION, RA = 0
32 51	+	28 35 00 05 04	558		DIV	BASE	" RA = LEFT BYTE, RX = RIGHT BYTE
32 52		00 00 50	559		JAN	WRONG FIELD	" ERROR IF LEFT BYTE < 0
32 53		00 05 37	560		STX	AUX	" SAVE THE RIGHT BYTE
32 54		00 05 70	561		CMPA	AUX (1:5)	" LEFT BYTE > RIGHT BYTE ?
32 55		00 06 47	562		JG	WRONG FIELD	" THEN ERROR
32 56	+	00 05 00 01 67	563		DECX	5	" RIGHT BYTE > 5 ?
32 57		00 02 57	564		JXP	WRONG FIELD	" THEN ERROR
32 58	+	12 12 00 05 10	565		LDA	VALUE	" VALUE TO BE STORED
32 59	+	12 11 00 05 11	566		LD1	EXPR	" FIELD SPECIFICATION
32 60	+	32 61 00 04 31	567		ST1	* + 1 (4 : 4)	" MODIFY NEXT INSTRUCTION
32 61	+	12 13 00 00 30	568		STA	CONSTANT (MOD)	" STORE VALUE INTO BYTES SPECIFIED
32 62	+	12 05 00 05 10	569	1H:	LDA	SYNT UNIT	" NOT BELONGING TO EXPRESSION
32 63	+	00 72 00 01 60	570		DECA	THESIS	" IS SYNTACTICAL UNIT A THESIS ?
32 64		00 01 50	571		JAZ	2F	" THEN O.K.
32 65	+	00 42 00 02 66	572		ENT6	42	" REPORT ERROR:
32 66	+	29 26 00 00 47	573		JMP	ERRORM	") REQUIRED
32 67		00 00 47	574		JMP	3F	" PROCEED WITH NEXT BYTES
32 68	+	00 41 00 02 66	575	WRONG FIELD:	ENT6	41	" REPORT ERROR:
*32 57	+	32 68 00 02 57					
*32 58	+	32 68 00 06 47					
*32 59	+	32 68 00 00 50					

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32 69	+	29 26 00 00 47	576	JMP	ERRORM	" ILLEGAL FIELD	
32 70	+	32 62 00 00 47	577	JMP	1B	" CONTINUE	
32 71	+	12 12 00 05 17	578	STANDARD:	LDX	VALUE	" VALUE TO BE STORED
*32 72	+	32 71 00 04 50					
32 72	+	12 13 00 05 37	579	STX	CONSTANT	" STORE THE VALUE IN ALL BYTES	
32 73		00 00 47	580	JMP	3F	" PROCEED WITH NEXT BYTES	
32 74	+	29 44 00 00 47	581	2H:	JMP	READ SYNT UNIT	" FOLLOWING THE THESIS
*32 64	+	32 74 00 01 50					
32 75	+	12 05 00 05 10	582	3H:	LDA	SYNT UNIT	" SYNTACTICAL UNIT
*32 73	+	32 75 00 00 47					
*32 67	+	32 75 00 00 47					
32 76	+	00 60 00 01 60	583	DECA	COMMA	" IS IT A COMMA ?	
32 77	+	32 42 00 01 50	584	JAZ	NEXT SPEC	" THEN TREAT NEXT SPECIFICATION	
32 78	+	12 13 00 05 10	585	LDA	CONSTANT	" RESULTING VALUE	
32 79	+	00 00 00 00 47	586	9H:	JMP	MOD	" RETURN
*32 39	+	32 79 00 02 40					
			587				
			588				
			589	READ PROGRAM:	" NO PARAMETERS		
32 80		00 02 40	590	STJ	9F	" SAVE RETURN ADDRESS	
32 81	+	28 37 00 05 10	591	NEXT INST:	LDA	END MARKER	
32 82	+	12 22 00 05 30	592	STA	LAST LABEL	" INITIALIZE CHAIN OF LABELS	
32 83	+	29 44 00 00 47	593	NEXT SYMBOL:	JMP	READ SYNT UNIT	" READ POSSIBLE LABEL
32 84	+	01 03 00 01 60	594	DECA	SYMBOL	" IS IT A SYMBOL ?	
32 85		00 01 50	595	JAZ	2F	" THEN JUMP	
32 86	-	00 02 00 01 60	596	DECA	SEP - SYMBOL	" OR IS IT A SEPARATOR ?	
32 87	+	32 83 00 01 50	597	JAZ	NEXT SYMBOL	" THEN SKIP IT	
32 88	+	00 51 00 02 66	598	ENT6	51	" REPORT ERROR:	
32 89	+	29 26 00 00 47	599	1H:	JMP	ERRORM	" SYMBOL REQUIRED
32 90	+	32 83 00 00 47	600	JMP	NEXT SYMBOL	" SKIP SYNTACTICAL UNIT	
32 91	+	12 24 00 05 31	601	2H:	ST1	INST	" SAVE POINTER TO THE SYMBOL
*32 87	+	32 91 00 01 50					
32 92	+	29 44 00 00 47	602	JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT	
32 93	+	00 63 00 01 60	603	DECA	COLON	" IS IT A COLON ?	
32 94		00 04 50	604	JANZ	TREAT INSTRUCTION	" ELSE JUMP	
32 95	+	12 07 00 05 10	605	LDA	TYPE	" TYPE OF LABEL	
32 96		00 01 50	606	JAZ	UNDEF SYMB	" JUMP IF IT IS AN UNDEF. SYMBOL	
32 97	+	00 02 00 01 60	607	DECA	LOCAL LABEL	" IS IT A LOCAL LABEL ?	
32 98		00 01 50	608	JAZ	LOC LAB	" THEN JUMP	
32 99	+	00 52 00 02 66	609	ENT6	52	" REPORT ERROR:	
33 00	+	32 89 00 00 47	610	JMP	1B	" ILLEGAL LABEL	
33 01	+	13 26 00 02 66	611	UNDEF SYMB:	ENT6	TABLE - 2 , 1	" POINTER TO FIELD VALUE1
32 96	+	33 01 00 01 50					
33 02	+	12 02 00 05 10	612	LDA	LINE NUMBER	") NUMBER OF THIS LINE	
33 03	+	13 27 01 23 30	613	STA	TABLE - 1 , 1 (2 : 3)	") IN INFO CELL	
33 04	+	33 06 00 00 47	614	JMP	* + 2		
33 05	+	12 48 01 02 66	615	LOC LAB:	ENT6	HERE, 1	" POINTER TO FIELD VALUE1
*32 98	+	33 05 00 01 50					
33 06	+	00 00 05 02 10	616	LDA	0 , 6 (VALUE1)	" CONTENTS OF FIELD VALUE1	
33 07		00 01 50	617	JAZ	3F	" JUMP IF INITIAL CONTENTS	
33 08	+	00 53 00 02 66	618	ENT6	53	" REPORT ERROR:	
33 09	+	32 89 00 00 47	619	JMP	1B	" (LOCAL) LABEL USED TWICE	
33 10	+	12 22 00 05 10	620	3H:	LDA	LAST LABEL	" CHAIN OF LABELS
*33 07	+	33 10 00 01 50					
33 11	+	00 00 06 02 30	621	STA	0 , 6 (VALUE1)	") ADD NEW LABEL TO	
33 12	+	12 22 00 05 31	622	ST1	LAST LABEL	") CHAIN OF LABELS	
33 13	+	32 83 00 00 47	623	JMP	NEXT SYMBOL	" TRY NEXT LABELS	
			624	TREAT INSTRUCTION:			
33 14	+	12 19 00 05 10	625	LDA	LC	" LOCATION COUNTER	
*32 94	+	33 14 00 04 50					

33 15	+	12 23 00 05 30	626	STA	EQUIV	" INITIALIZE EQUIV
33 16	+	12 24 00 05 10	627	LDA	INST	" INSTRUCTION SYMBOL
33 17	+	28 29 00 05 02	628	SUB	LAST INST	" IS IT A MIX MACHINE INSTRUCTION ?
33 18		00 00 50	629	JAN	6F	" ELSE JUMP
33 19	+	12 24 00 05 14	630	LD4	INST	" POINTER TO INSTRUCTION SYMBOL
33 20		00 00 47	631	JMP	READ MIX MACHINE	INST
33 21	+	12 19 00 05 17	632	LDX	LC	" LOCATION COUNTER
33 22		00 00 47	633	JMP	PRODUCE	" PRODUCE CODE
33 23	+	12 05 00 05 10	634	LDA	SYNT UNIT	" SYNTACTICAL UNIT
33 24	+	01 01 00 01 60	635	DECA	SEP	" IS IT A SEPARATOR ?
33 25		00 01 50	636	JAZ	2F	" THEN O.K.
33 26	+	00 55 00 02 66	637	ENT6	55	" REPORT ERROR!
33 27	+	29 26 00 00 47	638	JMP	ERRORM	" SEPARATOR REQUIRED
33 28	+	29 44 00 00 47	639	JMP	READ SYNT UNIT	") SKIP
33 29	+	01 01 00 01 60	640	DECA	SEP	") UNTIL
33 30	+	33 28 00 04 50	641	JANZ	1B	") SEPARATOR
33 31	+	12 22 00 05 15	642	LD5	LAST LABEL	" CHAIN OF LABELS
33 32	+	33 31 00 01 50				
33 32		00 00 47	643	JMP	UPDATE	" UPDATE THIS CHAIN
33 33	+	32 81 00 00 47	644	JMP	NEXT INST	" PROCEED WITH NEXT INSTRUCTION
33 34	+	12 24 00 05 10	645	LDA	INST	" INSTRUCTION SYMBOL
*33 18	+	33 34 00 00 50				
33 35	+	28 34 00 05 02	646	SUB	ORIG	" IS IT ORIG ?
33 36		00 04 50	647	JANZ	6F	" ELSE TRY OTHER
33 37	+	32 39 00 00 47	648	JMP	READ W VALUE	" READ THE WORD VALUE
33 38	+	12 19 00 05 30	649	STA	LC	" NEW VALUE FOR LOCATION COUNTER
33 39	+	33 23 00 00 47	650	JMP	5B	" GENERAL TREATMENT
33 40	+	12 24 00 05 10	651	LDA	INST	" INSTRUCTION SYMBOL
*33 36	+	33 40 00 04 50				
33 41	+	28 33 00 05 02	652	SUB	EQU	" IS IT EQU ?
33 42		00 04 50	653	JANZ	6F	" ELSE TRY OTHER
33 43	+	32 39 00 00 47	654	JMP	READ W VALUE	" READ THE WORD VALUE
33 44	+	12 23 00 05 30	655	STA	EQUIV	" NEW VALUE FOR EQUIV
33 45	+	33 23 00 00 47	656	JMP	5B	" GENERAL TREATMENT
33 46	+	12 24 00 05 10	657	LDA	INST	" INSTRUCTION SYMBOL
*33 42	+	33 46 00 04 50				
33 47	+	28 31 00 05 02	658	SUB	CON	" IS IT CON ?
33 48		00 04 50	659	JANZ	6F	" ELSE TRY OTHER
33 49	+	32 39 00 00 47	660	JMP	READ W VALUE	" READ THE WORD VALUE
33 50	+	33 21 00 00 47	661	JMP	4B	" GENERAL TREATMENT
33 51	+	12 24 00 05 10	662	LDA	INST	" INSTRUCTION SYMBOL
*33 48	+	33 51 00 04 50				
33 52	+	28 30 00 05 02	663	SUB	ALF	" IS IT ALF ?
33 53		00 04 50	664	JANZ	6F	" ELSE TRY OTHER
33 54	+	28 94 00 00 47	665	JMP	READ 5 CHAR	" READ THE FIVE CHARACTERS
33 55	+	33 21 00 00 47	666	JMP	4B	" GENERAL TREATMENT
33 56	+	12 24 00 05 10	667	LDA	INST	" INSTRUCTION SYMBOL
*33 53	+	33 56 00 04 50				
33 57	+	28 32 00 05 02	668	SUB	END	" IS IT END ?
33 58		00 04 50	669	JANZ	6F	" ELSE TRY OTHER
33 59	+	32 39 00 00 47	670	JMP	READ W VALUE	" READ THE WORD VALUE
33 60	+	12 27 00 05 30	671	STA	START ADDRESS	" VALUE FOR START ADDRESS
33 61		00 00 47	672	JMP	INSERT CONSTANTS	" INSERTION OF CONSTANTS
33 62	+	12 19 00 05 10	673	LDA	LC	" LOCATION COUNTER (AFTER INSERT)
33 63	+	12 23 00 05 30	674	STA	EQUIV	" NEW VALUE FOR EQUIV
33 64	+	12 22 00 05 15	675	LD5	LAST LABEL	" CHAIN OF LABELS PRECEDING END
33 65		00 00 47	676	JMP	UPDATE	" UPDATE THIS CHAIN
33 66	+	00 00 00 00 47	677	JMP	MOD	" RETURN
*32 80	+	33 66 00 02 40				
33 67	+	00 54 00 02 66	678	ENT6	54	" REPORT ERROR!


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*33 58 + 33 67 00 04 50
33 68 + 29 26 00 00 47
33 69 + 33 23 00 00 47
679
680
681
682
683
684
33 70 00 02 40 685
33 71 + 12 14 00 05 37 686
33 72 + 28 39 00 05 11 687
33 73 + 00 00 01 01 67 688
33 74 00 02 57 689
33 75 + 00 00 01 00 67 690
33 76 00 01 67 691
33 77 00 03 57 692
33 78 00 00 67 693
33 79 + 12 00 00 01 67 694
33 80 00 03 57 695
33 81 + 12 00 00 00 67 696
33 82 00 00 57 697
33 83 + 12 14 00 05 17 698
*33 77 + 33 83 00 03 57 699
33 84 + 00 00 00 00 47 700
*33 70 + 33 84 00 02 40 701
33 85 + 00 00 00 02 61 702
*33 82 + 33 85 00 00 57 703
*33 80 + 33 85 00 03 57 704
*33 74 + 33 85 00 02 57 705
33 86 + 33 83 00 00 47 706
707
33 87 00 02 40 707
*33 20 + 33 87 00 00 47 708
33 88 + 12 15 00 05 41 709
33 89 + 12 16 00 05 41 710
33 90 + 13 27 04 44 10 711
33 91 + 12 17 00 05 30 712
33 92 + 13 27 04 55 10 713
33 93 + 12 18 00 05 30 714
33 94 + 12 05 00 05 10 715
33 95 + 00 43 00 01 60 716
33 96 00 01 50 717
33 97 + 00 60 00 01 60 718
33 98 00 01 50 719
34 00 - 00 43 00 01 60 720
34 01 + 00 11 00 01 60 721
34 02 00 01 50 722
34 03 + 00 30 00 01 60 723
34 04 00 01 50 724
34 05 + 31 85 00 00 47 725
34 06 + 12 15 00 05 30 726
34 07 00 00 47 727
34 08 + 29 44 00 00 47 728
*33 96 + 34 08 00 01 50 729
34 09 + 32 39 00 00 47 730
34 10 + 12 26 00 05 16

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JMP ERRORM " UNKNOWN INSTRUCTION
JMP 5B " GENERAL TREATMENT

POSSIBLE: " RX CONTAINS THE ADDRESS,
" R11 IS A BOOLEAN OUTPUT PARAMETER: R11 # U # POSSIBLE
STJ 9F " SAVE RETURN ADDRESS
STX ADDRESS " SAVE THE ADDRESS
LDJ LENGTH OF MEMORY " ) LAST ADDRESS OF MEMORY
DECX 0, 1 " ) < ADDRESS ?
JXP 1F " THEN NOT POSSIBLE
INCX 0, 1 " ) ADDRESS
DECX END OF ASS " ) ≥ END OF ASSEMBLER ?
JXNN 8F " THEN POSSIBLE ADDRESS
INCX END OF ASS " ) ADDRESS
DECX BEGIN OF ASS " ) ≥ BEGIN OF ASSEMBLER ?
JXNN 1F " THEN NOT POSSIBLE
INCX BEGIN OF ASS " ADDRESS
JXN 1F " < 0 ? THEN NOT POSSIBLE
LDX ADDRESS " RESTORE ADDRESS INTO RX

8H:
9H: JMP MOD " RETURN
1H: ENT1 0 " NOT POSSIBLE ADDRESS INDICATION

701 JMP 8B " READY

READ MIX MACHINE INST: " R14 CONTAINS A POINTER TO THE INSTRUCTION SYMBOL,
" RA IS AN OUTPUT PARAMETER CONTAINING THE
" VALUE WHICH HAS BEEN ASSEMBLED
STJ 9F " SAVE RETURN ADDRESS
STZ A PART " INITIALIZE A PART
STZ I PART " INITIALIZE I PART
LDA TABLE - 1, 4 (4 : 4)
STA F PART " STANDARD FIELD INTO F PART
LDA TABLE - 1, 4 (5 : 5)
STA INSTRUCTION " SAVE INSTRUCTION CODE
LDA SYNT UNIT " LAST SYNTACTICAL UNIT
DECA EQUAL SIGN " IS IT AN EQUAL SIGN ?
JAZ LIT CONS " THEN TREAT LITERAL CONSTANT
DECA SYMBOL - EQUAL SIGN " IS IT A SYMBOL ?
JAZ FUT SYMB " THEN TRY FUTURE REFERENCE
DECA COMMA - SYMBOL " IS IT A COMMA ?
JAZ COMMA CHECKED " THEN TREAT I PART
DECA PAREN - COMMA " IS IT A PAREN ?
JAZ PAREN CHECKED " THEN TREAT F PART
DECA SEP - PAREN " IS IT A SEPARATOR ?
JAZ 8F " THEN READY
COMMON EXPR: JMP READ EXPRESSION " READ A PART
STA A PART " SAVE THE VALUE
JMP CHECK COMMA " PROCEED WITH OTHER PARTS
LIT CONS: JMP READ SYNT UNIT " NEXT SYNTACTICAL UNIT
JMP READ W VALUE " READ THE WORD VALUE
LD6 LIT CONS PTR " STACK POINTER FOR LITERALS

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34 11	+	13 28 06 05 30	731		STA	TABLE ,6	") STACK
34 12	+	00 01 00 00 66	732		INCS	1	") THE
34 13	+	12 26 00 05 36	733		ST6	LIT CONS PTR	") WORD VALUE
34 14	+	28 28 00 05 14	734		LD4	TABLE PTR1	" TABLE POINTER
34 15	+	00 00 05 01 64	735		DECA	0 ,6	" SPACE AVAILABLE ?
34 16		00 00 00 54	736		JAN	EXIT	" ELSE EXIT
34 17	+	12 05 00 05 10	737		LDA	SYNT UNIT	" SYNTACTICAL UNIT
34 18	+	00 43 00 01 60	738		DECA	EQUAL SIGN	" IS IT AN EQUAL SIGN ?
34 19	+	12 25 00 02 66	739		ENT6	LAST CONSTANT	" CHAIN OF LITERAL CONSTANTS
34 20		00 02 36	740		ST6	FUT REF (0 : 2)	" MODIFY INSTRUCTION
34 21		00 01 50	741		JAZ	2F	" READ, AND EXTEND CHAIN
34 22	+	00 61 00 02 66	742		ENT6	61	" REPORT ERROR:
34 23	+	29 26 00 00 47	743		JMP	ERRORM	" = REQUIRED
34 24		00 00 47	744		JMP	FUT REF	" EXTEND CHAIN OF FUTURE REFERENCES
34 25	+	12 07 00 05 10	745	FUT SYMB:	LDA	TYPE	" UNDEFINED SYMBOL ?
34 26	+	34 25 00 01 50					
34 27	+	13 26 01 02 66	746		JANZ	LOC	" ELSE TRY LOCAL FORWARD
34 28		00 04 50	747		ENT6	TABLE - 2 ,1	" POINTER TO FIELD VALUE2
34 29	+	00 04 00 01 60	748		JMP	1F	" TREAT FUTURE REFERENCE
34 30	+	34 29 00 04 50	749	LOC:	DECA	LOCAL FORWARD	" LOCAL FORWARD OCCURRING ?
34 31	+	34 05 00 04 50	750		JANZ	COMMON EXPR	" ELSE TRY COMMON EXPRESSION
34 32	+	12 48 01 02 66	751		ENT6	HERE ,1	" POINTER TO FIELD VALUE2
34 33		00 02 36	752	1H:	ST6	FUT REF (0 : 2)	" MODIFY INSTRUCTION
34 34	+	34 32 00 00 47					
34 35	+	29 44 00 00 47	753	2H:	JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT
34 36	+	34 33 00 01 50					
34 37	+	00 00 00 02 66	754	FUT REF:	ENT6	MOD	" RESTORE R16
34 38	+	34 34 00 02 36					
34 39	+	34 34 00 00 47					
34 40	+	34 34 00 02 36					
34 41	+	00 00 06 45 10	755		LDA	0 ,6 (VALUE2)	") CHAIN OF FUTURE REFERENCES
34 42	+	12 15 00 05 30	756		STA	APART	") INTO A PART
34 43	+	00 01 00 02 64	757		ENT4	1	
34 44	+	12 20 00 05 34	758		ST4	FUTURE	" NOTE OCCURRENCE OF FUTURE REFERENCE
34 45	+	12 19 00 05 17	759		LDX	LC	" LOCATION COUNTER
34 46	+	33 70 00 00 47	760		JMP	POSSIBLE	" IS IT A POSSIBLE ADDRESS ?
34 47		00 05 51	761		JMP	CHECK COMMA	" ELSE SKIP ADMINISTRATION
34 48	+	12 19 00 05 14	762		LD4	LC	" LOCATION COUNTER
34 49	+	00 00 04 05 10	763		LDA	0 ,4	" CONTENTS OF THIS ADDRESS
34 50		00 04 50	764		JANZ	CHECK COMMA	" JUMP IF USED ALREADY
34 51	+	12 21 00 05 10	765		LDA	DANGEROUS	" DANGEROUS ADDRESS
34 52	+	00 00 04 01 60	766		DECA	0 ,4	" EQUAL TO THIS ADDRESS ?
34 53		00 01 50	767		JAZ	CHECK COMMA	" THEN SKIP ADMINISTRATION
34 54	+	00 00 06 45 34	768		ST4	0 ,6 (VALUE2)	" JOIN THE ADDRESS TO THE CHAIN
34 55	+	12 05 00 05 10	769	CHECK COMMA:	LDA	SYNT UNIT	" SYNTACTICAL UNIT
34 56	+	34 49 00 01 50					
34 57	+	34 49 00 04 50					
34 58	+	34 49 00 05 51					
34 59	+	34 49 00 00 47					
34 60	+	00 60 00 01 60	770		DECA	COMMA	" IS IT A COMMA ?
34 61		00 04 50	771		JANZ	CHECK PAREN	" ELSE TRY FIELD SPECIFICATION
34 62	+	29 44 00 00 47	772	COMMA CHECKED:	JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT
34 63	+	34 52 00 01 50					
34 64	+	31 85 00 00 47	773		JMP	READ EXPRESSION	" READ 1 PART
34 65	+	12 16 00 05 30	774		STA	1 PART	" SAVE THE VALUE
34 66	+	12 05 00 05 10	775	CHECK PAREN:	LDA	SYNT UNIT	" SYNTACTICAL UNIT
34 67	+	34 55 00 04 50					
34 68	+	00 71 00 01 60	776		DECA	PAREN	" IS IT A PAREN ?
34 69		00 04 50	777		JANZ	CHECK VALUES	" ELSE JUMP

34 58	+	29 44 00 00 47	778	PAREN CHECKED:	JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT
*34 02	+	34 58 00 01 50					
34 59	+	31 85 00 00 47	779		JMP	READ EXPRESSION	" READ F PART
34 60	+	12 17 00 05 30	780		STA	F PART	" SAVE THE VALUE
34 61	+	12 05 00 05 10	781		LDA	SYNT UNIT	" SYNTACTICAL UNIT
34 62	+	00 72 00 01 60	782		DECA	THESIS	" IS IT A THESIS ?
34 63		00 04 50	783		JANZ	7F	" ELSE ERROR
34 64	+	29 44 00 00 47	784		JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT
34 65	+	28 37 00 05 10	785	CHECK VALUES:	LDA	MAX ADDRESS	" MAX, VALUE OF A PART
*34 57	+	34 65 00 04 50					
34 66	+	12 15 00 15 70	786		CMPA	A PART (1 : 5)	" ABS(A PART) O.K. ?
34 67		00 07 47	787		JGE	1F	" THEN CHECK I PART
34 68	+	00 63 00 02 66	788		ENT6	63	" REPORT ERROR:
34 69	+	29 26 00 00 47	789		JMP	ERRORM	" VALUE OF A PART NOT O.K.
34 70	+	12 15 00 05 41	790		STZ	A PART	" REPLACE ERRONEOUS A PART
34 71	+	28 36 00 05 10	791	1H:	LDA	MAX BYTE	" MAX, VALUE OF I PART
*34 67	+	34 71 00 07 47					
34 72	+	12 16 00 05 70	792		CMPA	I PART	" I PART TOO LARGE ?
34 73		00 04 47	793		JL	3F	" THEN ERROR
34 74	+	12 16 00 05 17	794		LDX	I PART	" I PART ≥ 0 ?
34 75		00 03 57	795		JXNN	1F	" THEN CHECK F PART
34 76	+	00 64 00 02 66	796	3H:	ENT6	64	" REPORT ERROR:
*34 73	+	34 76 00 04 47					
34 77	+	29 26 00 00 47	797		JMP	ERRORM	" VALUE OF I PART NOT O.K.
34 78	+	12 16 00 05 41	798		STZ	I PART	" REPLACE ERRONEOUS I PART
34 79	+	28 36 00 05 10	799		LDA	MAX BYTE	" MAX VALUE OF F PART
34 80	+	12 17 00 05 70	800	1H:	CMPA	F PART	" F PART TOO LARGE ?
*34 77	+	34 80 00 03 57					
34 81		00 04 47	801		JL	3F	" THEN ERROR
34 82	+	12 17 00 05 17	802		LDX	F PART	" F PART ≥ 0 ?
34 83		00 03 57	803		JXNN	8F	" THEN CHECKS O.K.
34 84	+	00 65 00 02 66	804	3H:	ENT6	65	" REPORT ERROR:
*34 81	+	34 84 00 04 47					
34 85	+	29 26 00 00 47	805		JMP	ERRORM	" VALUE OF F PART NOT O.K.
34 86	+	12 17 00 05 41	806		STZ	F PART	" REPLACE ERRONEOUS F PART
34 87	-	12 15 00 05 10	807	8H:	LDA	A PART	")
*34 83	+	34 87 00 03 57					
*34 04	+	34 87 00 01 50					
34 88	+	12 18 00 02 30	808		STA	INSTRUCTION (0 : 2)	") ASSEMBLE
34 89	+	12 16 00 05 10	809		LDA	I PART	") THE
34 90	+	12 18 00 03 30	810		STA	INSTRUCTION (3 : 3)	") WHOLE
34 91	+	12 17 00 05 10	811		LDA	F PART	") INSTRUCTION
34 92	+	12 18 00 04 30	812		STA	INSTRUCTION (4 : 4)	")
34 93	+	12 18 00 05 10	813		LDA	INSTRUCTION	" RESULTING VALUE
34 94	+	00 00 00 00 47	814	9H:	JMP	MOD	" RETURN
*33 87	+	34 94 00 02 40					
34 95	+	00 62 00 02 66	815	7H:	ENT6	62	" REPORT ERROR:
*34 83	+	34 95 00 04 50					
34 96	+	29 26 00 00 47	816		JMP	ERRORM	") REQUIRED
34 97	+	34 65 00 00 47	817		JMP	CHECK VALUES	" CONTINUE
			818				
			819				
			820	PRODUCE:		" RA CONTAINS THE VALUE TO BE PRODUCED,	
			821			" RX CONTAINS THE ADDRESS	
34 98		00 02 40	822		STJ	9F	" SAVE RETURN ADDRESS
*33 22	+	34 98 00 00 47					
34 99	+	33 70 00 00 47	823		JMP	POSSIBLE	" POSSIBLE ADDRESS TO BE FILLED IN ?
35 00		00 02 51	824		J1P	3F	" JUMP IF POSSIBLE ADDRESS
35 01	+	00 71 00 02 66	825		ENT6	71	" REPORT ERROR:
35 02	+	29 26 00 00 47	826		JMP	ERRORM	" WRONG ADDRESS

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3D 03	+	12 14 00 05 17	827		LDX	ADDRESS	" IMPOSSIBLE ADDRESS
3D 04	+	35 08 00 03 57	828		JXNM	* + 4	" JUMP IF NOT NEGATIVE
3D 05	+	00 38 00 02 64	829		ENTA	MINUS	") 'PRINT'
3D 06	+	13 04 00 44 34	830		ST4	OUTBUF (4 : 4)	") MINUS SIGN
3D 07	+	35 10 00 00 47	831		JMP	* + 3	
3D 08	+	00 37 00 02 64	832		ENTA	PLUS	") 'PRINT'
3D 09	+	13 04 00 44 34	833		ST4	OUTBUF (4 : 4)	") PLUS SIGN
3D 10	+	00 05 00 02 06	834		SLAX	5	" RA = IMPOSSIBLE ADDRESS, RX = 0
3D 11	+	00 00 00 01 05	835		CHAR		" CONVERT TO CHARACTERS
3D 12	+	13 05 00 05 30	836		ST4	OUTBUF + 1	") 'PRINT'
3D 13	+	13 06 00 05 37	837		STX	OUTBUF + 2	") THEN
3D 14		00 00 00 47	838		JMP	6F	" GIVE UP
3D 15	+	12 03 00 05 14	839	3H:	LDA	CODE	") IF CODE PRODUCED ALREADY
3D 16	+	35 15 00 02 51					
3D 17	+	29 16 00 02 54	840		J4P	PRINT LINE	") THEN PRINT THE LINE
3D 18	+	12 14 00 05 11	841		LD1	ADDRESS	" POSSIBLE ADDRESS
3D 19	+	12 19 00 05 17	842		LDX	LC	" LOCATION COUNTER
3D 20	+	00 00 01 01 67	843		DECX	0 ,1	" ARE THEY UNEQUAL ?
3D 21		00 04 57	844		JXNZ	4F	" THEN SKIP TEST
3D 22	+	00 00 01 05 17	845		LDX	0 ,1	" CONTENTS OF ADDRESS, ZERO ?
3D 23		00 04 57	846		JXNZ	1F	" ELSE ERROR
3D 24	+	12 21 00 05 17	847		LDX	DANGEROUS	") ADDRESS UNEQUAL
3D 25	+	00 00 01 01 67	848		DECX	0 ,1	") TO DANGEROUS ONE ?
3D 26		00 04 57	849		JXNZ	4F	" THEN O.K.
3D 27	+	00 72 00 02 66	850	1H:	ENTA	72	" REPORT ERROR;
3D 28	+	35 26 00 04 57					
3D 29	+	29 26 00 00 47	851		JMP	EPRORM	" ADDRESS ALREADY USED
3D 30	+	12 14 00 05 11	852		LD1	ADDRESS	
3D 31		00 00 47	853		JMP	6F	" PRODUCE OUTPUT ONLY
3D 32	+	12 12 00 05 30	854	4H:	ST4	VALUE	" SAVE THE VALUE TO BE PRODUCED
3D 33	+	35 30 00 04 57					
3D 34	+	35 30 00 04 57					
3D 35		00 04 50	855		JANZ	2F	" JUMP IF IT IS NOT ZERO
3D 36	+	12 20 00 05 17	856		LDX	FUTURE	" FUTURE REFERENCE TO BE PRODUCED ?
3D 37		00 05 57	857		JXNP	2F	" ELSE JUMP
3D 38	+	12 21 00 05 31	858		ST1	DANGEROUS	" DANGEROUS FUTURE REFERENCE PRODUCED
3D 39	+	00 00 01 01 05	859	2H:	CHAR		" CONVERT THE VALUE TO CHARACTERS
3D 40	+	35 35 00 05 57					
3D 41	+	35 35 00 04 50					
3D 42	+	13 07 00 05 30	860		ST4	OUTBUF + 3 (5 : 5)	
3D 43	+	13 08 00 05 37	861		STX	OUTBUF + 4	") 'PRINT' SIX DECIMAL POSITIONS
3D 44	+	12 20 00 05 14	862		LDA	FUTURE	" FUTURE REFERENCE ?
3D 45		00 02 54	863		J4P	5F	" THEN OMIT A PART
3D 46	+	13 07 00 05 30	864		ST4	OUTBUF + 3	") 'PRINT' A PART
3D 47	+	35 45 00 03 50	865		JANN	* + 4	" JUMP IF VALUE ≥ 0
3D 48	+	00 38 00 02 64	866		ENTA	MINUS	
3D 49	+	13 06 00 03 34	867		ST4	OUTBUF + 2 (3 : 3)	") 'PRINT' MINUS SIGN
3D 50		00 00 47	868		JMP	5F	
3D 51	+	00 37 00 02 64	869		ENTA	PLUS	
3D 52	+	13 06 00 03 34	870		ST4	OUTBUF + 2 (3 : 3)	") 'PRINT' PLUS SIGN
3D 53	+	12 12 00 05 10	871	5H:	LDA	VALUE	" VALUE TO BE PRODUCED
3D 54	+	35 47 00 00 47					
3D 55	+	35 47 00 02 54					
3D 56	+	00 00 01 05 30	872		ST4	0 ,1	" STORE IT INTO MEMORY
3D 57	+	00 00 01 02 60	873	6H:	ENTA	0 ,1	" ADDRESS
3D 58	+	35 49 00 00 47					
3D 59	+	00 00 01 01 05	874		CHAR		" CONVERT TO CHARACTERS
3D 60	+	13 04 00 05 37	875		STX	OUTBUF (2 : 5)	") 'PRINT' IT
3D 61	+	12 19 00 05 17	876		LDX	LC	" LOCATION COUNTER
3D 62	+	00 00 01 01 67	877		DECX	0 ,1	" UNEQUAL TO ADDRESS ?


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3D 54 + 35 58 00 04 57      878
3D 55 + 00 01 00 00 61      879
3D 56 + 12 29 00 05 31      880
3D 57 + 00 00 00 00 47      881
3D 58 + 00 39 00 02 64      882
3D 59 + 13 04 00 01 34      883
3D 50 + 12 20 00 05 41      884
*3D 51 + 35 60 00 00 47
*3D 14 + 35 60 00 00 47
3D 61 + 00 01 00 02 61      885
3D 62 + 12 03 00 05 34      886
3D 63 + 00 00 00 00 47      887
*3D 98 + 35 63 00 02 40

      888
      889
      890
      891
3D 64          00 02 40
*3D 65 + 35 64 00 00 47
*3D 32 + 35 64 00 00 47
3D 67 + 28 37 00 05 75      892
3D 68 + 00 00 00 05 47      893
*3D 69 + 35 66 00 02 40
3D 67 + 00 00 05 02 64      894
3D 68          00 05 55      895
3D 69 + 13 26 04 02 15      896
3D 70 + 13 26 04 05 16      897
3D 71 + 00 01 00 02 60      898
3D 72 + 13 27 04 01 30      899
3D 73 + 12 23 00 05 10      900
3D 74 + 13 26 04 05 30      901
3D 75          00 00 47      902
3D 76 + 12 48 04 02 15      903
*3D 68 + 35 76 00 05 55
3D 77 + 12 48 04 02 41      904
3D 78 + 12 48 04 05 16      905
3D 79 + 28 37 00 05 10      906
3D 80 + 12 48 04 05 30      907
3D 81 + 12 23 00 05 10      908
3D 82 + 12 38 04 05 30      909
3D 83 + 00 01 00 02 61      910
3D 84 + 12 48 04 03 31      911
3D 85 + 28 37 00 05 76      912
*3D 75 + 35 85 00 00 47
3D 86 + 35 65 00 05 47      913
3D 87 + 12 23 00 05 10      914
3D 88 + 28 37 00 05 70      915
3D 89          00 09 47      916
3D 90 + 00 81 00 02 66      917
3D 91 + 29 26 00 00 47      918
3D 92 + 35 65 00 00 47      919
3D 93 + 00 00 00 02 61      920
*3D 89 + 35 93 00 09 47
3D 94 + 00 00 01 02 16      921
3D 95 + 12 23 00 05 10      922
3D 96 + 00 00 01 02 30      923
3D 97 + 00 00 01 02 67      924
3D 98 + 00 00 01 05 10      925
3D 99 + 34 98 00 00 47      926
3D 00 + 35 85 00 00 47      927
      928

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      JXNZ * + 4          " THEN JUMP
      INC1 1            " INCREASE LOCATION COUNTER
      ST1 LC           " BY ONE
      JMP 8F           " READY
      ENT4 ASTERISK     " CHARACTER CODE OF ASTERISK
      ST4 OUTBUF (1 : 1) " PRINT IT
      STZ FUTURE      " RESET FUTURE ON FALSE

      ENT1 1           " NOW, CODE
      ST4 CODE         " HAS BEEN PRODUCED
      JMP MOD          " RETURN

      888
      889
      890
      891 UPDATE:      " R15 CONTAINS A POINTER TO THE CHAIN TO BE UPDATED
      STJ 9F           " SAVE RETURN ADDRESS

      892 LOOP1:      CMP5 END MARKER          " REMAINING CHAIN EMPTY ?
      893 9H:         JE MOD                   " THEN RETURN

      894          ENT4 0 , 5                  " REMAINING CHAIN
      895          J5NF 1F                     " JUMP IF REFERENCE TO LOCAL LABEL
      896          LD5 TABLE - 2 , 4 (VALUE1) " DELETE LAST LABEL
      897          LD6 TABLE - 2 , 4 (VALUE2) " CHAIN OF FUTURE REFERENCES
      898          ENTA DEFINED SYMB          " SYMBOL
      899          ST1 TABLE - 1 , 4 (0 : 1) " DEFINED NOW
      900          LDA EQUIV                   " VALUE OF
      901          ST4 TABLE - 2 , 4         " DEFINED SYMBOL
      902          JMP LOOP2                  " UPDATE CHAIN OF FUTURE REFERENCES
      903 1H:         LD5 HERE , 4 (VALUE1)   " DELETE LAST LABEL FROM CHAIN

      904          STZ HERE , 4 (VALUE1)      " INITIALIZE CHAIN AGAIN
      905          LD6 HERE , 4 (VALUE2)      " CHAIN OF FUTURE REFERENCES
      906          LDA END MARKER
      907          ST4 HERE , 4 (VALUE2)     " EMPTY AGAIN
      908          LDA EQUIV                   " VALUE OF
      909          ST4 BACKWARD , 4          " LOCAL BACKWARD
      910          ENT1 DEFINED SYMB         " NOTE OCCURRENCE OF
      911          ST1 HERE , 4 (3 : 3)     " LOCAL BACKWARD
      912 LOOP2:      CMP6 END MARKER          " REMAINING CHAIN OF FUTURE REFERENCES

      913          JE LOOP1                   " IF EMPTY, PROCEED WITH LABEL CHAIN
      914          LDA EQUIV (1 : 5)         " ABS(EQUIV)
      915          CMPA MAX ADDRESS          " NOT TOO LARGE ?
      916          JLE 1F                     " THEN O.K.
      917          ENT6 81                    " REPORT ERROR!
      918          JMP ERRORM                " VALUE FOR FUTURE REFERENCE TOO LARGE
      919          JMP LOOP1                 " CONTINUE WITH LABEL CHAIN
      920 1H:         ENT1 0 , 6             " ADDRESS TO BE UPDATED

      921          LD6 0 , 1 (0 : 2)         " NEXT ADDRESS IN CHAIN
      922          LDA EQUIV                   " UPDATE A PART
      923          ST4 0 , 1 (0 : 2)         " OF THE ADDRESS
      924          ENTX 0 , 1                 " ADDRESS IN RX
      925          LDA 0 , 1                  " NEW CONTENTS IN RA
      926          JMP PRODUCE                " PRODUCE OUTPUT
      927          JMP LOOP2                 " CONTINUE WITH FUTURE REFERENCES

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929
930 TRAVERSE TREE: " R14 CONTAINS A SUBROUTINE FOR INVESTIGATE
931 STJ 9F " SAVE RETURN ADDRESS
932 ST4 INVESTIGATE (0 : 2) " ACTUAL SUBROUTINE INVESTIGATE
933 ENT2 LENGTH OF TABLE - 1 " POINTER TO ROOT OF TREE
934 1H: LD4 TABLE ,2 (2 : 3) " LEFT BRANCH, EMPTY ?
935 J4Z INVESTIGATE " THEN JUMP
936 ENT2 0 ,4 " TRAVERSE LEFT BRANCH
937 JMP 1B " REPEAT
938 INVESTIGATE: JMP MOD " INVESTIGATE THE NODE

939 LDA TABLE ,2 " THREAD ?
940 LD2 TABLE ,2 (4 : 5) " RIGHT BRANCH
941 JAP 1B " TRAVERSE TO LEFT IF NOT A THREAD
942 CMP2 END MARKER " WHOLE TREE TRAVERSED ?
943 JNE INVESTIGATE " ELSE CONTINUE
944 9H: JMP MOD " RETURN

945
946
947 BUFFER CHAR: " RA CONTAINS THE CHARACTER TO BE BUFFERED,
948 " R14 AND R15 ARE PERMANENT FILL POINTERS
949 STJ 9F " SAVE RETURN ADDRESS
950 ST5 * + 1 (4 : 4) " MODIFY NEXT INSTRUCTION
951 ST4 OUTBUF ,4 (MOD) " 'PRINT' THE CHARACTER
952 ENN5 5 : 5 /N ,5 " BYTE FIVE OF A WORD FILLED ?
953 J5P 8F " ELSE JUMP
954 ENT5 5 : 5 " ) BYTE 0 OF
955 INC4 1 " ) NEXT WORD
956 8H: ENN5 6 : 6 /N ,5 " POINTER TO NEXT BYTE

957 9H: JMP MOD " RETURN

958
959
960 BUFFER SYMB: " R12 CONTAINS A POINTER TO THE SYMBOL TO BE BUFFERED,
961 " R14 AND R15 CONTAIN THE POSITION FOR THE FIRST LETTER
962 STJ 9F " SAVE RETURN ADDRESS
963 J2P 4F " JUMP IF ORDINARY SYMBOL
964 ENNA 1 ,2 " DIGIT OF LOCAL SYMBOL
965 JMP BUFFER CHAR " 'PRINT' IT
966 ENTA LETTER H " LETTER H
967 JMP BUFFER CHAR " 'PRINT' IT
968 9H: JMP MOD " RETURN

969 4H: ENT1 - 3 ,2 " POINTER TO CHARACTERS

970 1H: LDX TABLE ,1 " NEXT FIVE CHARACTERS
971 ENT6 5 : 5 " INITIALIZE BYTE POINTER
972 2H: ENN6 6 : 6 /N ,6 " POINTER TO NEXT BYTE
973 ST6 * + 1 (4 : 4) " MODIFY NEXT INSTRUCTION
974 LD4 TABLE ,1 (MOD) " 'READ' NEXT CHARACTER
975 JAZ 9B " IF EMPTY THEN READY
976 DECA 1 " CANCEL MODIFICATION OF CHARACTER
977 JMP BUFFER CHAR " 'PRINT' THE CHARACTER
978 ENN6 5 : 5 /N ,6 " BYTE FIVE 'READ' ?
979 J6P 2B " ELSE LOOP
980 JXN 0B " READY IF LAST WORD
981 DECA 1 " NEXT WORD WITH CHARACTERS

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-57-


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36 44 + 36 32 00 00 47 982          JMP      1B          " CONTINUE
983
984
985  BUFFER CONSTANT: " R11 CONTAINS A POINTER TO THE SYMBOL
36 45          00 02 40 986          STJ      9F          " SAVE RETURN ADDRESS
36 46 + 12 03 00 05 10 987          LDA      CODE        " CODE PRODUCED ?
36 47 + 12 04 00 05 01 988          ADD      TEXT        " OR TEXT AVAILABLE ?
36 48 + 29 16 00 02 50 989          JAP      PRINT LINE  " THEN PRINT THE LINE
36 49 + 00 08 00 02 64 990          ENT4    8           " ) POSITION IN BUFFER
36 50 + 00 11 00 02 65 991          ENT5    1 : 1       " ) FOR SYMBOL
36 51 + 36 24 00 00 47 992          JMP      BUFFER SYMB " 'PRINT' THE SYMBOL
36 52 + 00 63 00 02 60 993          ENT4    COLON      "
36 53 + 36 15 00 00 47 994          JMP      BUFFER CHAR " 'PRINT' A COLON
36 54 + 13 16 00 11 10 995          LDA      OUTBUF+8 + 4 (1 : 1) " POSITION FOR FIRST LETTER
36 55 + 00 66 00 01 60 996          DECA   SPACE      " USED ALREADY ?
36 56 + 29 16 00 04 50 997          JANZ   PRINT LINE " THEN PRINT THE LINE
36 57 + 00 12 00 02 64 998          ENT4    8 + 4      " ) POSITION IN BUFFER
36 58 + 00 11 00 02 65 999          ENT5    1 : 1       " ) FOR THIS SYMBOL
36 59          00 02 32 1000         ST2     1F (0 : 2)   " SAVE POINTER TO SYMBOL
36 60 + 28 31 00 05 12 1001         LD2     CON         " POINTER TO CONSTANT INSTRUCTION SYMBOL
36 61 + 36 24 00 00 47 1002         JMP      BUFFER SYMB " 'PRINT' CONSTANT INSTRUCTION SYMBOL
36 62 + 00 00 00 02 62 1003 1H:      ENT2    MOD         " RESTORE POINTER TO SYMBOL
36 63 + 00 01 00 00 64 1004         INC4   1           "
36 64 + 13 04 04 22 41 1005         STZ    OUTBUF ,4 (2 : 2) " 'PRINT' WORD VALUE ZERO
36 65 + 00 00 00 00 47 1006 9H:      JMP      MOD         " RETURN
36 66          00 02 40 1007
36 67          00 02 40 1008
36 68          00 02 40 1009  DEFINE:      " R12 CONTAINS A POINTER TO THE SYMBOL TO BE DEFINED
36 69 + 13 27 02 01 10 1010         STJ      9F          " SAVE RETURN ADDRESS
36 70          00 02 50 1011         LDA      TABLE - 1 ,2 (0 : 1) " DEFINED SYMBOL ?
36 71 + 13 26 02 02 10 1012         JAP      9F          " THEN READY
36 72          00 04 50 1013         LDA      TABLE - 2 ,2 (VALUE1) " CONTENTS OF FIELD VALUE1
36 73 + 13 26 02 02 63 1014         JANZ   9F          " READY IF NOT INITIAL VALUE
36 74          00 02 63 1015         ENT3    TABLE - 2 ,2 " POINTER TO FIELD VALUE2
36 75          00 02 40 1016  DEFINE1:   " R 3 CONTAINS A POINTER TO FIELD VALUE2
36 76          00 02 40 1017         STJ      9F          " SAVE RETURN ADDRESS (AGAIN)
36 77 + 00 00 03 05 10 1018         LDA      0 ,3 (VALUE2) " CHAIN OF FUTURE REFERENCES
36 78 + 28 37 00 05 70 1019         CMPA   END MARKER " EMPTY CHAIN ?
36 79          00 05 47 1020         JE      9F          " THEN READY
36 80 + 36 45 00 00 47 1021         JMP      BUFFER CONSTANT " 'PRINT' A CONSTANT INSTRUCTION
36 81 + 12 19 00 05 17 1022         LDX     LC          " LOCATION COUNTER
36 82 + 12 23 00 05 37 1023         STX     EQUIV      " NEW VALUE FOR EQUIV
36 83 + 00 00 00 02 60 1024         ENT4    0          " WORD VALUE OF CONSTANT INSTRUCTION
36 84 + 34 98 00 00 47 1025         JMP      PRODUCE   " PRODUCE IT
36 85 + 28 37 00 05 10 1026         LDA      END MARKER " ) SIMULATE OCCURRENCE
36 86 + 00 00 03 02 30 1027         STA     0 ,3 (VALUE1) " ) AS LABEL (FOR UPDATE)
36 87 + 00 00 02 02 65 1028         ENT5    0 ,2      " PARAMETER FOR UPDATE
36 88 + 35 64 00 00 47 1029         JMP      UPDATE    " UPDATE THE CHAIN
36 89 + 00 00 00 00 47 1030 9H:      JMP      MOD         " RETURN
36 90 + 36 85 00 05 47 1031
36 91 + 36 85 00 02 40 1032
36 92 + 36 85 00 04 50 1033
36 93 + 36 85 00 02 50 1034  INSERT CONSTANTS: " NO PARAMETERS
36 94 + 36 85 00 02 40 1034         STJ      9F          " SAVE RETURN ADDRESS

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*33 61 + 36 86 00 00 47
36 87 + 00 01 00 02 60 1035
36 88 + 12 03 00 05 30 1036
36 89 - 00 10 00 02 62 1037
36 90 + 12 48 02 02 63 1038
36 91 + 36 72 00 00 47 1039
36 92 + 00 01 00 00 62 1040
36 93 + 36 90 00 00 52 1041
36 94 + 36 66 00 02 64 1042
36 95 + 36 01 00 00 47 1043
36 96 + 12 25 00 05 15 1044
36 97 + 28 37 00 05 75 1045
36 98 + 00 00 00 05 47 1046
36 99 + 36 98 00 02 40
37 00 + 12 25 00 05 30 1048
37 01 + 12 19 00 05 17 1049
37 02 + 00 00 05 02 37 1050
37 03 + 12 26 00 05 16 1051
37 04 + 00 01 00 01 66 1052
37 05 + 12 26 00 05 36 1053
37 06 + 13 28 06 05 10 1054
37 07 + 34 98 00 00 47 1055
37 08 + 00 00 05 02 67 1056
37 09 + 00 00 05 05 10 1057
37 10 + 34 98 00 00 47 1058
37 11 + 36 96 00 00 47 1059
1060
1061
1062
37 12 + 00 02 40 1063
37 13 + 37 12 00 00 47 1064
37 14 + 00 00 00 01 05 1065
37 15 + 13 10 00 25 37 1066
37 16 + 00 66 00 02 60 1067
37 17 + 13 10 00 22 17 1068
37 18 + 00 04 57 1069
37 19 + 13 10 00 22 30 1070
37 20 + 13 10 00 03 17 1071
37 21 + 00 04 57 1072
37 22 + 13 10 00 33 30 1073
37 23 + 13 10 00 44 17 1074
37 24 + 00 04 57 1075
37 25 + 13 10 00 44 30 1076
*37 26 + 00 00 00 00 47 9H:
*37 27 + 37 25 00 04 57
*37 28 + 37 25 00 04 57
*37 29 + 37 25 00 04 57
*37 30 + 37 25 00 04 57
*37 31 + 37 25 00 04 57
*37 32 + 37 25 00 04 57
*37 33 + 37 25 00 02 40
1077
1078
1079
37 34 + 00 02 40 1080
37 35 + 13 27 02 01 10 1081
37 36 + 00 01 50 1082
37 37 + 13 26 02 05 10 1083
37 38 + 00 00 00 01 05 1084
37 39 + 13 07 00 05 30 1085
37 40 + 13 08 00 05 37 1086
37 41 + 00 00 50 1087

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1H:
ENTA 1 " CAUSE THE PRINTING OF
STA CODE " A NEW LINE IF NECESSARY
ENT2 - 10 " INITIALIZE LOCAL SYMBOL POINTER
ENT3 HERE , 2 " POINTER TO FIELD VALUE2
JMP DEFINE1 " DEFINE THE SYMBOL IF NECESSARY
INC2 1 " INCREASE POINTER, READY ?
J2N 1B " ELSE CONTINUE
ENT4 DEFINE " PARAMETER FOR TRAVERSE TREE
JMP TRAVERSE TREE " DEFINE SYMBOLS IF NECESSARY
2H:
LD5 LAST CONSTANT " CHAIN OF LITERAL CONSTANTS
CMP5 END MARKER " IS IT EMPTY ?
9H:
JE MOD " THEN RETURN

LDA 0 , 5 ( 0 : 2) " DELETE LAST CONSTANT
STA LAST CONSTANT " FROM CHAIN
LDX LC " LOCATION COUNTER
STX 0 , 5 ( 0 : 2) " UPDATE A PART OF ADDRESS
LD6 LIT CONS PTR " UNSTACK
DEC6 1 " VALUE FOR
ST6 LIT CONS PTR " LITERAL CONSTANT
LDA TABLE , 6 " VALUE OF IT
JMP PRODUCE " PRODUCE THE INSERTED WORD VALUE
ENTX 0 , 5 " UPDATED ADDRESS
LDA 0 , 5 " NEW CONTENTS OF IT
JMP PRODUCE " PRODUCE OUTPUT
JMP 2B " CONTINUE

BUFFER LN:
" R4 CONTAINS THE LINE NUMBER
STJ 9F " SAVE RETURN ADDRESS

CHAR " CONVERT LINE NUMBER TO CHARACTERS
STX OUTBUF+8 - 2 ( 2 : 5) " STORE THEM INTO BUFFER
ENTA SPACE " INTERNAL CODE OF A SPACE
LDX OUTBUF+8 - 2 ( 2 : 2) " FIRST DECIMAL
JXNZ 9F " READY IF NOT A ZERO
STA OUTBUF+8 - 2 ( 2 : 2) " REPLACE THE ZERO BY A SPACE
LDX OUTBUF+8 - 2 ( 3 : 3) " SECOND DECIMAL
JXNZ 9F " READY IF NOT A ZERO
STA OUTBUF+8 - 2 ( 3 : 3) " REPLACE THE ZERO BY A SPACE
LDX OUTBUF+8 - 2 ( 4 : 4) " THIRD DECIMAL
JXNZ 9F " READY IF NOT A ZERO
STA OUTBUF+8 - 2 ( 4 : 4) " REPLACE THE ZERO BY A SPACE
9H:
JMP MOD " RETURN

DUMP:
" R12 CONTAINS A POINTER TO THE SYMBOL
STJ 9F " SAVE RETURN ADDRESS
LDA TABLE - 1 , 2 ( 0 : 1)
JAZ 9F " READY IF SYMBOL UNDEFINED YET
LDA TABLE - 2 , 2 " VALUE OF SYMBOL
CHAR " PRINT
STA OUTBUF+8 - 5 " THIS
STX OUTBUF+8 - 4 " VALUE
JAN 1F " JUMP IF THE VALUE IS NEGATIVE

```


37 34	+	00 37 00 02 60	1088		ENTA	PLUS	") 'PRINT' PLUS SIGN
37 35	+	13 06 00 44 30	1089		STA	OUTBUF+8 - 6 (4 : 4)	
37 36		00 00 00 47	1090		JMP	2F	
37 37	+	00 38 00 02 60	1091	1H:	ENTA	MINUS	") 'PRINT' MINUS SIGN
*37 38	+	37 37 00 00 50					
37 38	+	13 06 00 44 30	1092		STA	OUTBUF+8 - 6 (4 : 4)	
37 39	+	13 27 02 23 10	1093	2H:	LDA	TABLE - 1 , 2 (2 : 3)	" LINE OF DEFINING OCCURRENCE
*37 36	+	37 39 00 00 47					
37 40		00 04 50	1094		JANZ	3F	" JUMP IF IT HAS BEEN OCCURRED
37 41	+	00 37 00 02 60	1095		ENTA	PLUS	") 'PRINT' A PLUS SIGN
37 42	+	13 11 00 11 30	1096		STA	OUTBUF+8 - 1 (1 : 1)	") INSTEAD OF LINE NUMBER
37 43		00 00 47	1097		JMP	4F	
37 44	+	37 12 00 00 47	1098	3H:	JMP	BUFFER LN	" 'PRINT' THE LINE NUMBER
*37 40	+	37 44 00 04 50					
37 45	+	00 08 00 02 64	1099	4H:	ENTA	8	") POSITION IN BUFFER
*37 43	+	37 45 00 00 47					
37 46	+	00 11 00 02 65	1100		ENT5	1 : 1	") FOR SYMBOL
37 47	+	36 24 00 00 47	1101		JMP	BUFFER SYMB	" 'PRINT' THE SYMBOL
37 48	+	29 16 00 00 47	1102		JMP	PRINT LINE	" PRINT THE LINE
37 49	+	00 00 00 00 47	1103	9H:	JMP	MOD	" RETURN
*37 28	+	37 49 00 01 50					
*37 26	+	37 49 00 02 40					
			1104				
			1105				
			1106	INITIALIZE:		" NO PARAMETERS	
37 50		00 02 40	1107		STJ	9F	" SAVE RETURN ADDRESS
37 51	+	28 28 00 05 11	1108		LDI	TABLE PTR1	") FIRST ADDRESS PRECEDING RELEVANT
37 52	+	13 28 01 02 61	1109		ENT1	TABLE , 1	") INFORMATION FOR ASSEMBLER
37 53		00 00 51	1110	1H:	JIN	2F	" JUMP IF ADDRESS ZERO TREATED
37 54	+	00 00 01 05 41	1111		STZ	0 , 1	" INITIALIZE ADDRESS ON ZERO
37 55	+	00 01 00 01 61	1112		DEC1	1	" NEXT ADDRESS
37 56	+	37 53 00 00 47	1113		JMP	1B	" CONTINUE
37 57	+	28 37 00 05 20	1114	2H:	LDAN	END MARKER	
*37 53	+	37 57 00 00 51					
37 58	+	12 21 00 05 30	1115		STA	DANGEROUS	" INITIALIZE DANGEROUS
37 59	+	28 37 00 05 10	1116		LDA	END MARKER	
37 60	+	12 25 00 05 30	1117		STA	LAST CONSTANT	" INITIALIZE LAST CONSTANT
37 61	+	12 38 00 05 30	1118		STA	HERE - 10	") INITIALIZE
37 62	+	12 39 00 02 61	1119		ENT1	HERE - 9	") ARRAYS 'H' AND 'F'
37 63	+	12 38 00 09 07	1120		MOVE	HERE - 10 (9)	
37 64	+	12 02 00 05 41	1121	INIT READING:	STZ	LINE NUMBER	" INITIALIZE LINE NUMBER
37 65	+	13 04 00 02 61	1122		ENT1	OUTBUF	
37 66	+	28 41 00 01 07	1123		MOVE	SPACES (1)	") INITIALIZE
37 67	+	13 04 00 23 07	1124		MOVE	OUTBUF (23)	") OUTPUT BUFFER
37 68	+	12 48 00 16 44	1125		IN	INBUF X (READER)	" READ THE FIRST CARD
37 69	+	00 01 00 02 62	1126		ENT2	1	") INITIALIZE
37 70	+	00 55 00 02 63	1127		ENT3	5 : 5	") READING
37 71	+	00 66 00 02 60	1128		ENTA	SPACE	") POINTERS
37 72	+	12 00 00 05 30	1129		STA	CHAR	")
37 73	+	00 00 00 00 47	1130	9H:	JMP	MOD	" RETURN
*37 50	+	37 73 00 02 40					
			1131				
			1132				
			1133	INIT TABLE:		" NO PARAMETERS	
37 74		00 02 40	1134		STJ	9F	" SAVE RETURN ADDRESS
37 75		00 02 61	1135		ENT1	1F	") RETURN ADDRESS AFTER INIT READING
37 76	+	37 73 00 02 31	1136		ST1	9B (0 : 2)	") INTO RETURN OF INIT READING
37 77	+	37 64 00 00 47	1137		JMP	INIT READING	" INITIALIZE READING POINTERS
37 78	+	14 99 00 02 61	1138	1H:	ENT1	LENGTH OF TABLE - 1	" POINTER TO ROOT OF TREE
*37 75	+	37 78 00 02 61					

37 79	+	28 28 00 05 31	1139		ST1	TABLE PTR1	" INITIALIZE TABLE POINTER
37 81	+	29 44 00 00 47	1140	OPERS:	JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT
37 81	+	01 03 00 01 60	1141		DECA	SYMBOL	" IS IT A SYMBOL ?
37 82		00 04 50	1142		JANZ	2F	" ELSE TRY SEPARATOR
37 83	+	31 49 00 00 47	1143		JMP	READ AT EXPRESSION	" READ THE NUMBER
37 84	+	12 06 00 05 11	1144		LD1	LAST SYMBOL	" POINTER TO INSTRUCTION SYMBOL
37 85	+	13 27 01 45 30	1145		STA	TABLE - 1 ,1 (4 : 5)	" STORE STANDARD FIELD AND CODE
37 86	+	29 44 00 00 47	1146		JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT
37 87	+	01 03 00 01 60	1147		DECA	SYMBOL	"
37 88	-	00 02 00 01 60	1148	2H:	DECA	SEP - SYMBOL	" IS IT A SEPARATOR ?
37 82	+	37 88 00 04 50					
37 89	+	37 91 00 02 52	1149		J2P	* + 2	" JUMP IF NOT A SEMICOLON
37 90	+	12 04 00 05 41	1150		STZ	TEXT	" PREVENT NEW LINE
37 91	+	37 80 00 01 50	1151		JAZ	OPERS	" CONTINUE IF SEPARATOR READ
37 92	-	00 30 00 01 60	1152		DECA	PAREN - SEP	" IS SYNTACTICAL UNIT A PAREN ?
37 93		00 01 50	1153		JAZ	PSEUDO INSTS	" THEN PROCEED WITH PSEUDO INSTRUCTIONS
37 94	-	00 01 00 02 66	1154		ENT6	01	" REPORT ERROR:
37 95	+	29 26 00 00 47	1155		JMP	ERRORM	" INSTRUCTION DEFINITIONS NOT O.K.
37 96	+	00 00 00 02 05	1156		HLT		" STOP THE MACHINE
37 97	+	12 06 00 05 11	1157	PSEUDO INSTS:	LD1	LAST SYMBOL	" LAST INSTRUCTION SYMBOL
37 93	+	37 97 00 01 50					
37 98	+	28 29 00 05 31	1158		ST1	LAST INST	" INITIALIZE LAST INST
37 99	+	29 44 00 00 47	1159		JMP	READ SYNT UNIT	
38 00	+	28 30 00 05 31	1160		ST1	ALF	" INITIALIZE ALF
38 01	+	29 44 00 00 47	1161		JMP	READ SYNT UNIT	
38 02	+	28 31 00 05 31	1162		ST1	CON	" INITIALIZE CON
38 03	+	29 44 00 00 47	1163		JMP	READ SYNT UNIT	
38 04	+	28 32 00 05 31	1164		ST1	END	" INITIALIZE END
38 05	+	29 44 00 00 47	1165		JMP	READ SYNT UNIT	
38 06	+	28 33 00 05 31	1166		ST1	EQU	" INITIALIZE EQU
38 07	+	29 44 00 00 47	1167		JMP	READ SYNT UNIT	
38 08	+	28 34 00 05 31	1168		ST1	ORIG	" INITIALIZE ORIG
38 09	+	29 44 00 00 47	1169		JMP	READ SYNT UNIT	" NEXT SYNTACTICAL UNIT
38 10	+	00 72 00 01 60	1170		DECA	THIS IS	" IS IT A THIS IS ?
38 11	+	29 16 00 00 47	1171		JMP	PRINT LINE	
38 12	+	00 00 00 01 50	1172	9H:	JAZ	MOD	" THEN RETURN
38 74	+	38 12 00 02 40					
38 13	+	00 03 00 02 66	1173		ENT6	03	" REPORT ERROR:
38 14	+	29 26 00 00 47	1174		JMP	ERRORM	" PSEUDO INSTR. DEFINITIONS NOT O.K.
38 15	+	00 00 00 02 05	1175		HLT		" STOP THE MACHINE
			1176				
			1177				
			1178	RESET TABLE:		" NO PARAMETERS	
38 16		00 02 40	1179		STJ	9F	" SAVE RETURN ADDRESS
38 55	+	38 16 00 00 47					
38 17	+	28 34 00 05 12	1180		LD2	ORIG	" LAST (PSEUDO) INSTRUCTION SYMBOL
38 18	+	00 03 00 01 62	1181		DEC2	3	" POINTER TO CHARACTERS
38 19	+	13 28 02 05 10	1182	3H:	LDA	TABLE ,2	" LAST WORD WITH CHARACTERS ?
38 20	+	00 01 00 01 62	1183		DEC2	1	" ELSE SEARCH FOR
38 21	+	38 19 00 02 50	1184		JAP	3B	" THE LAST ONE
38 22	+	28 28 00 05 32	1185		ST2	TABLE PTR1	" NEW FREE SPACE
38 23	+	14 99 00 02 62	1186		ENT2	LENGTH OF TABLE - 1	" ROOT OF TREE
38 24	+	13 28 02 23 14	1187	4H:	LD4	TABLE ,2 (2 : 3)	" LEFT BRANCH, EMPTY ?
38 25		00 01 54	1188		J4Z	VISIT	" THEN VISIT NODE TO BE SAVED
38 26	+	28 28 00 05 74	1189		CMP4	TABLE PTR1	" THIS NODE TO BE DELETED ?
38 27		00 04 47	1190		JLE	DELETE	" THEN DELETE IT
38 28	+	00 00 04 02 62	1191		ENT2	0 ,4	" TRAVERSE LEFT BRANCH
38 29	+	38 24 00 00 47	1192		JMP	4B	" AND CONTINUE
38 30	+	13 28 02 23 41	1193	DELETE:	STZ	TABLE ,2 (2 : 3)	" EMPTY LEFT BRANCHE NOW
38 27	+	38 30 00 09 47					

38 31	+	00 00 02 02 63	1194	VISIT:	ENT3	0 , 2	" POINTER TO SAVED NODE
*38 29	+	38 31 00 01 54					
38 32	+	13 28 03 05 10	1195		LDA	TABLE , 3	") THREAD ? CR
38 33	+	13 28 03 45 12	1196		LD2	TABLE , 3 (4 : 5)	") RIGHT BRANCH ?
38 34	+	13 27 03 03 41	1197		STZ	TABLE - 1 , 3 (0 : 3)	" RESET TYPE AND LINE NUMBER
38 35	+	28 37 00 05 17	1198		LDX	END MARKER	") RESET THE FIELDS
38 36	+	13 26 03 05 37	1199		STX	TABLE - 2 , 3	") VALUE1 AND VALUE2
38 37		00 00 00 50	1200		JAN	CMP END MARKER	" JUMP IF THREAD
38 38	+	28 28 00 05 72	1201		CMP2	TABLE PTR1	" RIGHT BRANCH
38 39	+	38 24 00 06 47	1202		JG	4B	" JUMP IF RIGHT NODE TO BE SAVED
38 40	+	13 28 02 05 10	1203	5H:	LDA	TABLE , 2	") THREAD ? CR
38 41	+	13 28 02 05 12	1204		LD2	TABLE , 2 (4 : 5)	") RIGHT BRANCH ?
38 42	+	38 40 00 02 50	1205		JAP	5B	" ELSE SEARCH FOR THE THREAD
38 43	+	00 00 02 03 60	1206		ENNA	0 , 2	" TAKE THE THREAD (NEGATIVE)
38 44	+	13 28 03 45 30	1207		STA	TABLE , 3 (4 : 5)	" STORE THE THREAD IN SAVED NODE
38 45	+	13 28 03 00 30	1208		STA	TABLE , 3 (0 : 0)	" THREAD INDICATION
38 46	+	28 37 00 05 72	1209	CMP END MARKER:	CMP2	END MARKER	" IS THIS THREAD THE LAST ONE ?
*38 37	+	38 46 00 00 50					
38 47	+	38 31 00 08 47	1210		JNE	V'SIT	" ELSE CONTINUE
38 48	+	00 00 00 00 47	1211	9H:	JMP	MOD	" RETURN
*38 10	+	38 48 00 02 40					
			1212				
			1213				
			1214	MIXAL ASSEMBLER:			" RA IS A BOOLEAN OUTPUT PARAMETER; RA = - ERRONEOUS,
			1215				" R1 IS AN OUTPUT PARAMETER FOR THE START ADDRESS
38 49		00 02 40	1216		STJ	9F	" SAVE RETURN ADDRESS
38 50	+	37 50 00 00 47	1217		JMP	INITIALIZE	" INITIALIZE THE MEMORY
38 51	+	00 00 00 18 43	1218		IOC	0 (PRINTER)	" NEW PAGE
*38 52	+	32 80 00 00 47	1219		JMP	READ PROGRAM	" ASSEMBLE THE PROGRAM
38 53	+	29 16 00 00 47	1220		JMP	PRINT LINE	" PRINT THE LAST LINE
38 54	+	12 27 00 05 17	1221		LDX	START ADDRESS	" START ADDRESS
38 55	+	33 70 00 00 47	1222		JMP	POSSIBLE	" IS IT A POSSIBLE ADDRESS ?
38 56		00 02 51	1223		J1P	1F	" THEN JUMP
38 57	+	00 91 00 02 66	1224		ENT6	91	" REPORT ERROR:
38 58	+	29 26 00 00 47	1225		JMP	ERRORM	" WRONG START ADDRESS
38 59	+	12 09 00 05 10	1226	1H:	LDA	ERRONEOUS	" ERROR RECOVERED ?
*38 56	+	38 59 00 02 51					
38 60		00 01 50	1227		JAZ	2F	" ELSE JUMP
38 61	+	00 92 00 02 66	1228		ENT6	92	" REPORT ERROR:
38 62	+	29 26 00 00 47	1229		JMP	ERRORM	" PROGRAM INCORRECT
38 63	+	00 00 00 18 43	1230	2H:	IOC	0 (PRINTER)	" NEW PAGE
*38 58	+	38 63 00 01 50					
38 64	+	37 26 00 02 64	1231		ENT4	DUMP	" PARAMETER FOR TRAVERSE TREE
38 65	+	36 01 00 00 47	1232		JMP	TRAVERSE TREE	" DUMP THE SYMBOL TABLE
38 66	+	29 16 00 00 47	1233		JMP	PRINT LINE	
38 67	+	38 16 00 00 47	1234		JMP	RESET TABLE	" INITIALIZE SYMBOL TABLE AGAIN
38 68	+	12 09 00 05 20	1235		LOAN	ERRONEOUS	" COMPLEMENT OF ERRONEOUS IN RA
38 69	+	12 27 00 45 11	1236		LD1	START ADDRESS (4 : 5)	" START ADDRESS IN R1
38 70	+	00 00 00 00 47	1237	9H:	JMP	MOD	" RETURN
*38 49	+	38 70 00 02 40					
38 71	+	00 99 00 02 66	1238	EXIT:	ENT6	99	" REPORT ERROR:
*38 10	+	38 71 00 00 54					
*29 01	+	38 71 00 02 56					
*29 08	+	38 71 00 02 56					
38 72	+	29 26 00 00 47	1239		JMP	ERRORM	" OVERFLOW SYMBOL TABLE
38 73	+	28 45 00 00 47	1240		JMP	READ CHAR	") SKIP (BUT PRINT)
38 74	+	38 73 00 00 47	1241		JMP	* - 1	") REST OF INPUT
			1242				
			1243				
			1244				


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1245 " NEXT FOLLOWS THE PIECE OF PROGRAM, SERVING AS A VERY SIMPLE MONITORING
1246 " PROGRAM, ACCEPTING SEVERAL MIXAL PROGRAMS, BEING SEPARATED FROM EACH OTHER
1247 " BY A JOB SEPARATOR (" " IN THE COLUMNS 79 AND 80),
1248 " THESE MIXAL PROGRAMS ARE TRANSLATED AND EXECUTED (IF CORRECT) CONSECUTIVELY.
1249 " THE JOB SEPARATOR IS DETECTED BY THE ROUTINE "READ CHAR" ONLY; SO DURING
1250 " EXECUTION OF A PROGRAM, THE JOB SEPARATOR WILL BE IGNORED.
1251 " AS THE ASSEMBLER ALWAYS READS ONE CARD IN ADVANCE, THE CARD SUCCEEDING
1252 " THE END INSTRUCTION IS READ TOO BY THE ROUTINE "MIXAL ASSEMBLER".
1253 " THIS FACT HAS TO BE TAKEN INTO ACCOUNT WHEN IN - INSTRUCTIONS ARE PERFORMED
1254 " BY A PARTICULAR MIXAL PROGRAM.
1255 " THE MONITOR MAKES USE OF THE AUXILIARY ROUTINES "SKIP INPUT" AND
1256 " "PUNCH TRAJECT".
1257
1258
1259 START0: LD2 TABLE PTR1
1260 ENT2 TABLE + 1 ,2 " START ADDRESS OF TRAJECT
1261 ENT3 0 ,2 " ) LENGTH OF
1262 DECB END OF ASS " ) TRAJECT (NEGATIVE)
1263 JMP PUNCH TRAJECT " PUNCH PERMANENT OF ASSEMBLER
1264 START1: JMP INIT TABLE " INITIALIZE THE SYMBOL TABLE
1265 NEXT PROGRAM: JMP MIXAL ASSEMBLER " TRANSLATE THE PROGRAM
1266
1267 IOC 0 (PRINTER) " NEW PAGE
1268 JBUS * (READER) " SKIP THE FIRST INPUT CARD
1269 JANM 0 ,1 " EXECUTE THE PROGRAM IF CORRECT
1270 JMP SKIP INPUT " SKIP REST OF INPUT OF PROGRAM
1271 JMP NEXT PROGRAM " TREAT THE NEXT PROGRAM
1272
1273 SKIP INPUT: " NO PARAMETERS
1274 STJ 9F " SAVE RETURN ADDRESS
1275
1276 1H: IN INBUF (READER) " READ NEXT CARD
1277 JBUS * (READER) " WAIT FOR CARD READER
1278 LDA INBUF + 15 " POSSIBLE JOB SEPARATOR
1279 CMPA JOB SEPARATOR (4 : 5) " IS IT A JOB SEPARATOR ?
1280 JNE 1B " ELSE CONTINUE
1281 JMP MOD " RETURN
1282
1283 PUNCH TRAJECT: " R12 CONTAINS THE START ADDRESS,
1284 " R13 CONTAINS THE LENGTH OF THE TRAJECT (NEGATIVE)
1285 STJ 9F " SAVE RETURN ADDRESS
1286
1287 ENT1 OUTBUF - 1 " )
1288 MOVE SPACES (1) " ) INITIALIZE
1289 MOVE OUTBUF - 1 (16) " ) OUTPUT BUFFER
1290 ENT5 PLUS " CODE OF PLUS SIGN
1291 ENT6 MINUS " CODE OF MINUS SIGN
1292 ENTA 0 ,2 " START ADDRESS
1293 CHAR " CONVERT IT TO CHARACTERS
1294 STX OUTBUF " STORE THEM INTO THE OUTPUT BUFFER
1295 ENNA 0 ,3 " LENGTH OF TRAJECT (POSITIVE)
1296 CHAR " CONVERT IT TO CHARACTERS
1297 STX OUTBUF + 2 " STORE THEM INTO THE OUTPUT BUFFER
1298 DEC2 0 ,3 " START ADDRESS + LENGTH OF TRAJECT
1299 ST LOAD (0 : 2) " MODIFY INSTRUCTION
1300 OUT OUTBUF (CPUNCH) " PUNCH A CARD
1301 JBUS * (CPUNCH) " WAIT FOR CARD PUNCH

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39 10	+	00 00 00 01 53	1301	9H1	J3Z	MOD	" RETURN IF TRAJECT COMPLETED
*38 94	+	39 10 00 02 40					
39 11	+	13 04 00 02 61	1302		ENT1	OUTBUF	") INITIALIZE
39 12	+	13 03 00 06 07	1303		MOVE	OUTBUF - 1 (16)	") OUTPUT BUFFER AGAIN
39 13	+	00 15 00 03 64	1304		ENN4	15	" INITIALIZE POSITION ON CARD
39 14	+	00 00 03 05 10	1305	LOAD:	LDA	MOD ,3	" NEXT WORD
*39 07	+	39 14 00 02 32					
39 15		00 00 50	1306		JAN	1F	" JUMP IF NEGATIVE
39 16	+	13 19 04 53 35	1307		ST5	OUTBUF + 15 ,4 (3 : 3)	" STORE PLUS SIGN INTO BUFFER
39 17		00 00 47	1308		JMP	2F	
39 18	+	13 19 04 53 36	1309	1H:	ST6	OUTBUF + 15 ,4 (3 : 3)	" STORE MINUS SIGN INTO BUFFER
*39 15	+	39 18 00 00 50					
39 19	+	00 00 00 01 05	1310	2H:	CHAR		" CONVERT WORD TO CHARACTERS
*39 17	+	39 19 00 00 47					
39 20	+	13 20 04 05 30	1311		STA	OUTBUF + 16 ,4	") STORE THEM
39 21	+	13 21 04 05 37	1312		STX	OUTBUF + 17 ,4	") INTO OUTPUT BUFFER
39 22	+	00 03 00 00 64	1313		INC4	3	" NEXT POSITION ON CARD
39 23	+	00 01 00 00 63	1314		INC3	1	" NEXT WORD FROM MEMORY
39 24	+	39 08 00 01 54	1315		J4Z	OUT	" JUMP IF CARD COMPLETED
39 25	+	39 14 00 04 53	1316		J3NZ	LOAD	" JUMP IF TRAJECT NOT COMPLETED
39 26	+	39 08 00 00 47	1317		JMP	OUT	" ELSE PUNCH LAST CARD
			1318				
			1319				
			1320	END OF ASS:	END	START0	
39 27	+	00 00 00 00 00		AUX:	CON	0	
*32 54	+	39 27 00 05 70					
*32 53	+	39 27 00 05 37					
*30 25	+	39 27 00 05 03					
*30 25	+	39 27 00 05 30					
39 28	+	00 00 00 00 11					
*36 19	+	39 28 00 05 04					
*35 78	+	39 29 00 01 63					
*33 78	+	39 29 00 00 67					
*33 76	+	39 29 00 01 67					

START ADDRESS: +3875

TRANSLATION TIME: 330.54 SEC.

VALUE AND LINE NUMBER OF DEFINED SYMBOLS:

+	00	00	00	12	01	57	5CHAR
+	00	00	00	32	14	518	ADD
+	00	00	00	12	14	70	ADDRESS
+	00	00	00	28	30	102	ALF
+	00	00	00	12	15	71	APART
+	00	00	00	32	25	529	ARITHMOV
+	00	00	00	00	39	19	ASTERISK
+	00	00	00	31	82	482	ASTSK
+	00	00	00	12	10	66	ATEXPR
+	00	00	00	31	52	452	ATEXPRESSION
+	00	00	00	39	27	+	AUX
+	00	00	00	30	66	362	B
+	00	00	00	12	38	88	BACKWARD
+	00	00	00	28	35	107	BASE
+	00	00	00	12	00	56	BEGINOFASS
+	00	00	00	36	15	949	BUFFERCHAR
+	00	00	00	36	45	986	BUFFERCONSTANT
+	00	00	00	37	12	1063	BUFFERLN
+	00	00	00	36	24	962	BUFFERSYMB
+	00	00	00	01	00	108	BYTESIZE
+	00	00	00	12	00	56	CHAR
+	00	00	00	34	49	769	CHECKCOMMA
+	00	00	00	34	55	775	CHECKPAREN
+	00	00	00	34	65	785	CHECKVALUES
+	00	00	00	38	46	1209	COMPENMARKER
+	00	00	00	12	03	59	CODE
+	00	00	00	00	63	25	COLON
+	00	00	00	00	60	22	COMMA
+	00	00	00	34	52	772	COMMACHECKED
+	00	00	00	34	05	725	COMMONEXPR
+	00	00	00	29	88	283	COMPL
+	00	00	00	28	31	103	CON
+	00	00	00	12	13	69	CONSTANT
+	00	00	00	30	47	343	CONTAINSLETTER
+	00	00	00	28	44	119	CORRECTION
+	00	00	00	00	17	7	CPUNCH
+	00	00	00	12	21	77	DANGEROUS
+	00	00	00	00	63	25	DEC
+	00	00	00	36	66	1010	DEFINE
+	00	00	00	36	72	1017	DEFINE1
+	00	00	00	31	72	472	DEFINED
+	00	00	00	00	01	40	DEFINEDSYMB
+	00	00	00	38	30	1193	DELETE
+	00	00	00	31	34	431	DIV
+	00	00	00	00	40	20	DIV1
+	00	00	00	01	04	36	DIV2
+	00	00	00	32	34	533	DIVA
+	00	00	00	32	29	533	DIVX
+	00	00	00	37	26	1060	DUMP
+	00	00	00	28	32	104	END
+	00	00	00	28	37	111	ENDMARKER
+	00	00	00	39	29	1320	ENDCFASS
+	00	00	00	28	33	105	EQU
+	00	00	00	10	43	21	EQUALSIGN
+	00	00	00	12	23	79	EQUIV
+	00	00	00	12	09	65	ERRONEOUS
+	00	00	00	28	43	118	ERROR
+	00	00	00	19	26	217	ERRORM

+	00	00	00	38	71	1238	EXIT
+	00	00	00	12	11	67	EXPR
+	00	00	00	30	64	360	F
+	00	00	00	31	09	406	FOUND
+	00	00	00	12	17	73	FPART
+	00	00	00	34	34	754	FUTREF
+	00	00	00	34	25	745	FUTSYMB
+	00	00	00	12	20	76	FUTURE
+	00	00	00	12	48	90	HERE
+	00	00	00	28	40	115	HUNCREDTHOUSAND
+	00	00	00	12	64	91	INBUF
+	00	00	00	12	48	90	INBUFx
+	00	00	00	37	50	1107	INITIALIZE
+	00	00	00	37	64	1121	INITREADING
+	00	00	00	37	74	1134	INITTABLE
+	00	00	00	36	86	1034	INSERTCONSTANTS
+	00	00	00	12	24	80	INST
+	00	00	00	12	18	74	INSTRUCTION
+	00	00	00	36	08	938	INVESTIGATE
+	00	00	00	12	16	72	IPART
+	00	00	00	28	42	117	JOBSEPARATOR
+	00	00	00	12	25	81	LASTCONSTANT
+	00	00	00	28	29	101	LASTINST
+	00	00	00	12	22	78	LASTLABEL
+	00	00	00	12	06	62	LASTSYMBOL
+	00	00	00	12	19	75	LC
+	00	00	00	30	86	383	LEFT
+	00	00	00	28	39	114	LENGTHOFMEMORY
+	00	00	00	12	00	48	LENGTHOFTABLE
+	00	00	00	29	50	245	LETGIT
+	00	00	00	00	11	12	LETTERB
+	00	00	00	00	15	13	LETTERF
+	00	00	00	00	17	14	LETTERH
+	00	00	00	00	35	15	LETTERZ
+	00	00	00	12	02	50	LINENUMBER
+	00	00	00	34	08	728	LITCONS
+	00	00	00	12	26	82	LITCONSPTR
+	00	00	00	39	14	1305	LOAD
+	00	00	00	34	29	749	LOC
+	00	00	00	00	03	42	LOCALBACKWARD
+	00	00	00	00	04	43	LOCALFORWARD
+	00	00	00	00	02	41	LOCALLABEL
+	00	00	00	31	74	474	LOCDEFINED
+	00	00	00	33	05	615	LOCLAB
+	00	00	00	30	70	367	LOOKFORSYMBOL
+	00	00	00	35	65	892	LOOP1
+	00	00	00	35	85	912	LOOP2
+	00	00	00	28	37	111	MAXADDRESS
+	00	00	00	28	36	109	MAXBYTE
+	00	00	00	28	38	112	MAXWORD
+	00	00	00	00	38	17	MINUS
+	00	00	00	38	49	1216	MIXALASSEMBLER
+	00	00	00	00	00	46	MOD
+	00	00	00	32	21	525	MUL
-	00	00	00	00	01	47	N
+	00	00	00	31	01	398	NEW
+	00	00	00	30	71	308	NEXT
+	00	00	00	32	81	591	NEXTINST
+	00	00	00	31	93	496	NEXTOPERAND
+	00	00	00	38	81	1265	NEXTPROGRAM

+ 00 00 00 32 42	549	NEXTSPEC
+ 00 00 00 32 83	593	NEXTSYMBOL
+ 00 00 00 30 92	30	NLCR
+ 00 00 00 31 80	480	NUMB
+ 00 00 00 01 02	36	NUMBER
+ 00 00 00 37 80	1140	OPERS
+ 00 00 00 28 34	106	ORIG
+ 00 00 00 31 44	441	OTHER
+ 00 00 00 39 08	1299	OUT
+ 00 00 00 13 04	93	OUTBUF
+ 00 00 00 12 80	92	OUTBUFX
+ 00 00 00 30 43	338	OVERFLOW
+ 00 00 00 00 71	28	PAREN
+ 00 00 00 34 58	778	PARENCHECKED
+ 00 00 00 00 37	16	PLUS
+ 00 00 00 00 61	23	POINT
+ 00 00 00 33 70	685	POSSIBLE
+ 00 00 00 00 18	8	PRINTER
+ 00 00 00 29 16	204	PRINTLINE
+ 00 00 00 34 98	822	PRODUCE
+ 00 00 00 37 97	1157	PSEUDOINSTS
+ 00 00 00 28 94	1285	PUNCHTRAJECT
+ 00 00 00 00 94	31	QUOTE
+ 00 00 00 29 63	258	READ
+ 00 00 00 28 94	179	READ5CHAR
+ 00 00 00 31 49	449	READATEXPRESSION
+ 00 00 00 28 76	158	READC
+ 00 00 00 28 45	127	READCHAR
+ 00 00 00 00 16	6	READER
+ 00 00 00 31 85	488	READEXPRESSION
+ 00 00 00 33 87	707	READMIXMACHINEINST
+ 00 00 00 32 80	590	READPROGRAM
+ 00 00 00 29 44	239	READSYNTUNIT
+ 00 00 00 32 39	546	READWVALUE
+ 00 00 00 38 16	1179	RESETTABLE
+ 00 00 00 30 95	392	RIGHT
+ 00 00 00 01 00	108	S
+ 00 00 00 00 64	26	SEMICOLON
+ 00 00 00 01 01	35	SEP
+ 00 00 00 31 19	416	SEPARATOR
+ 00 00 00 38 87	1274	SKIPINPUT
+ 00 00 00 00 66	27	SPACE
+ 00 00 00 28 41	116	SPACES
+ 00 00 00 32 71	578	STANDARD
+ 00 00 00 38 75	1257	START0
+ 00 00 00 38 80	1264	START1
+ 00 00 00 12 27	83	STARTADDRESS
+ 00 00 00 32 17	521	SUB
+ 00 00 00 31 64	464	SYMB
+ 00 00 00 01 03	37	SYMBOL
+ 00 00 00 12 05	61	SYNTUNIT
+ 00 00 00 13 28	54	TABLE
+ 00 00 00 28 28	100	TABLEPTR1
+ 00 00 00 12 04	60	TEXT
+ 00 00 00 00 72	29	THESIS
+ 00 00 00 00 39	19	TIMES
+ 00 00 00 32 10	514	TIMESBASEPLUS
+ 00 00 00 36 01	931	TRAVERSE TREE
+ 00 00 00 33 14	825	TREATINSTRUCTION
+ 00 00 00 12 07	63	TYPE

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+	00 00 00 00 00	39	UNDEFINESYMB
+	00 00 00 33 01	611	UNDEFSYMB
+	00 00 00 35 64	891	UPDATE
+	00 00 00 12 12	68	VALUE
+	00 00 00 00 02	44	VALUE1
+	00 00 00 00 45	45	VALUE2
+	00 00 00 12 08	64	VALUEOFNUMBER
+	00 00 00 38 31	1194	VISIT
+	00 00 00 32 68	575	WRONGFIELD

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1	J3NP	553;	ENT5	265;	LD4	514;	DEC6	166;	INC6	066
2	JL	447;	SLAX	206;	CMP6	576;	ENN5	365;	FSUB	602
3	J1P	251;	J5N	056;	JXP	257;	LDX	517;	ST4	534
4	CMP2	572;	DEC2	162;	ENN1	361;	ENT1	261;	FADD	601
5	INC2	062;	J1N	051;	J2NP	552;	J4P	254;	JANP	550
6										
7	JXN	057;	LD2	512;	LD6	516;	NOP	000;	SRC	506
8	STJ	240;	CHAR	105;	CMP4	574;	CMPX	577;	DEC4	164
9	DECX	167;	ENN3	363;	ENNA	360;	ENT3	263;	ENTA	260
10	FDIV	604;	IN	044;	INC4	064;	INCX	067;	J1NP	551
11	J2N	052;	J2P	252;	J4N	054;	J5NP	555;	J5P	256
12										
13	JBUS	042;	JNOV	347;	JXNP	557;	LD1	511;	LD3	513
14	LD5	515;	LDA	510;	MOVE	107;	OUT	045;	SRA	106
15	ST2	532;	ST6	536;	STZ	541;	ADD	501;	CMP1	571
16	CMP3	573;	CMP5	575;	CMPA	570;	DEC1	161;	DEC3	163
17	DEC5	165;	DECA	160;	DIV	504;	ENN2	362;	ENN4	364
18										
19	ENN6	366;	ENNX	367;	ENT2	262;	ENT4	264;	ENT6	266
20	ENTX	267;	FCMP	670;	FMUL	603;	HLT	205;	INC1	061
21	INC3	063;	INC5	065;	INCA	060;	IOC	043;	J1NN	351
22	J1NZ	451;	J12	151;	J2NN	352;	J2NZ	452;	J3N	053
23	J3P	253;	J4NP	554;	J5N	055;	J5P	255;	J6NP	556
24										
25	JAN	050;	JAP	250;	JG	647;	JMP	047;	JRED	046
26	JXNN	357;	JXNZ	457;	JXZ	157;	LD1N	521;	LD2N	522
27	LD3N	523;	LD4N	524;	LD5N	525;	LD6N	526;	LDAN	520
28	LDXN	527;	MUL	503;	NUM	005;	SLA	006;	SLC	406
29	SRAX	306;	ST1	531;	ST3	533;	ST5	535;	STA	530
30										
31	STX	537;	SUB	502;	J2Z	152;	J3NN	353;	J3NZ	453
32	J3Z	153;	J4NN	354;	J4NZ	454;	J4Z	154;	J5NN	355
33	J5NZ	455;	J5Z	155;	J6NN	356;	J6NZ	456;	J6Z	156
34	JANN	350;	JANZ	450;	JAZ	150;	JE	547;	JGE	747
35	JLE	947;	JNE	847;	JOV	247;	JSJ	147		
36	(ALF		CON		END		EQU		ORIG)	

8.1 Table of error messages.

error number:	description:
ER 01	instruction definitions not o.k.
ER 03	pseudo instruction definitions not o.k.
ER 10	. required.
ER 11	value of number too large.
ER 21	atomic expression required.
ER 22	undefined local backward in expression.
ER 23	undefined symbol in expression.
ER 31	overflow by arithmetic operation.
ER 41	illegal field.
ER 42) required.
ER 51	symbol required.
ER 52	illegal label.
ER 53	(local) label used twice.
ER 54	unknown instruction.
ER 55	separator required.
ER 61	= required.
ER 62) required.
ER 63	value of a part not o.k.
ER 64	value of i part not o.k.
ER 65	value of f part not o.k.
ER 71	wrong address.
ER 72	address already used.
ER 81	value for future reference too large.
ER 91	wrong start address.
ER 92	program incorrect.
ER 98	end of input.
ER 99	overflow symbol table.

9. A loading program.

Before discussing the loading program, we have to treat a feature of the MIX simulator, concerning the start of the computer. With respect to the MIXAL system, consisting of simulator and assembler (see introduction), there is no question of a start of the simulator from the point that no relevant information is in the memory; the proper MIX simulator, without assembler, has the problem of how to load the first instructions into memory. For this purpose, the proper simulator has been provided with a simulated "start button", being "pressed" at the beginning of the execution of the MIX simulator. Pressing this start button, the MIX instruction "IN 0(16)" is executed, which causes the activation of the card reader; after one card has been read, and the information of it has been stored into the locations zero upwards, the normal cycle of the computer is started at location zero.

Preparing this first card by hand, we thus can get the first simple loading program, consisting of at most 16 instructions, into the machine. Each position on this card corresponds to one byte; this byte is filled with the internal code of the character on that position (this is the usual effect of an IN instruction).

Unfortunately, this function on the collection of characters into the numbers $\{n \mid 0 \leq n < 100\}$ is not a surjective one; e.g. the instruction codes of the IN and JBUS instructions (44 and 42, resp.) do not correspond to any MIX character. As these two instructions are indispensable in a loading program, we have to find another way to get these instructions into the memory. An easy solution is: constructing these instructions programmatically (e.g. by means of the instructions: "ENT2 40; INC2 4; ST2 X(5:5)", when an IN instruction has to be loaded at location X). Till so far, we have discussed the start of this special MIX simulator, and noticed some difficulties and possibilities of it.

The loading program of this section consists of three cards, the proper loader being contained in the second and third card. The first card serves to construct the instructions for reading the next two cards, and, after these cards have been read, to construct the IN and JBUS instructions of the proper loader. Next, control is passed on to the proper loader, and the

locations 0 - 15 are used as an input buffer by the loader. Much attention had to be paid to arranging the instructions of the proper loader in such a way, that the value of the bytes of all instructions did correspond to an existing MIX character.

Eventually, the code of the information to be loaded is given. The input of the loader consists of a number of blocks of information, each block consisting of a number of information cards, preceded by one traject card. On the traject card, two numbers have been punched: a startaddress (columns 1 - 5), and a quantity indicating the number of consecutive locations (beginning at the startaddress) to be loaded (columns 11 - 15). Each information card contains the information to be loaded into (the next) five consecutive locations. Column 3 contains the sign (+ or -) of the first location, columns 6 - 15 contain the numerical value for this location in decimal notation (two columns for each byte); columns 18 and 21 - 30 contain the information for the second location, etc..

The last of the information cards contains as much information as is necessary to complete this particular block.

A traject card, which contains a zero as second number, indicates the final block. In this case, the startaddress on the traject card is interpreted as the location at which the loaded program has to be started.

This kind of code is produced by the routine called "punch traject" of the MIXAL version of the assembler. This routine punches one block of information, including the traject card. The startaddress is given by the value of index register 2, and the number of consecutive locations to be punched out, is given by the negative value of index register 3.

Next follows the MIXAL text of the entire loading program (some instructions have been added, from location 100 upwards, to punch the three initial cards, mentioned before, which contain the loading system), in the way it has been printed by the MIXAL version of the MIXAL assembler.

0000	+	1640000261	1	INBUF:	ENT1	1640
0001	+	0002010262	2		ENT2	2 ,1
0002	+	0004000061	3		INC1	4
0003		004531	4		ST1	INCARD1 (4:5)
0004		004531	5		ST1	INCARD2 (4:5)
0005		004532	6		ST2	JBUS (4:5)
0006	+	0016000000	7	INCARD1:	NOP	16
*0007	+	0006000031				
0007	+	0032000000	8	INCARD2:	NOP	32
*0008	+	0007000031				
0008	+	0008000000	9	JBUS:	NOP	*
*0009	+	0008000032				
0009		004531	10		ST1	READINFCARDS (4:5)
0010		004532	11		ST2	WAITINFCARDS (4:5)
0011		004531	12		ST1	READTRAJECT (4:5)
0012		004532	13		ST2	WAITTRAJECT (4:5)
0013		000047	14		JMP	READTRAJECT
			15			
			16		ORIG	15
0015	+	0000000000	17	AUX:	CON	0
0016	+	0000000038	18	MINUS:	ALF	.0000-
0017	+	0003000063	19	1H:	INC3	3
0018		000453	20		J3NZ	2F
0019	+	0015000063	21		ENN3	15
0020	+	0000000000	22	READINFCARDS:	NOP	INBUF
*0009	+	0020000031				
0021	+	0021000000	23	WAITINFCARDS:	NOP	*
*0010	+	0021000032				
0022	+	001033311	24	2H:	LD1	INBUF + 15 ,3 (3:3)
*0018	+	0022000053				
0023	+	0016030010	25		LDA	INBUF + 16 ,3
0024	+	0017030017	26		LDX	INBUF + 17 ,3
0025	+	0000000005	27		NUM	
0026	+	0016000071	28		CMP1	MINUS
0027		000047	29		JNE	STORE
0028	+	0015000030	30		STA	AUX
0029	+	0015000020	31		LDAN	AUX
0030	+	0000000030	32	STORE:	STA	0 ,2
*0027	+	0030000047				
0031	+	0001000062	33		INC2	1
0032	+	0017000052	34		J2NZ	1B
0033	+	0000000000	35	READTRAJECT:	NOP	INBUF
*0013	+	0033000047				
*0011	+	0033000031				
0034	+	0034000000	36	WAITTRAJECT:	NOP	*
*0012	+	0034000032				
0035	+	00000000260	37		ENTA	0
0036	+	0002000017	38		LDX	INBUF + 2
0037	+	0000000005	39		NUM	
0038	+	00390000230	40		STA	* + 1 (0:2)
0039	+	00000000362	41		ENN2	0
0040	+	00000000260	42		ENTA	0
0041	+	00000000517	43		LDX	INBUF
0042	+	0000000005	44		NUM	
0043	+	00450000230	45		STA	* + 2 (0:2)
0044	+	00000000160	46		DECA	0 ,2
0045	+	00000000152	47		J2Z	0
0046	+	00300000230	48		STA	STORE (0:2)
0047	+	0019000047	49		JMP	READINFCARDS - 1
			50			

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0100	000240	51	ORIG	100
0101	+ 0000001745	52	STJ	9F
0102	+ 0016001745	53	CUT	0 (17)
0103	+ 0032001745	54	CUT	16 (17)
0104	+ 0104001742	55	CUT	32 (17)
0105	+ 0107000147	56	JBUS	* (17)
0106	+ 0105000240	57	JMP	*
		58	END	100

9H:

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55

04

+ 0000000015	17	AUX
+ 0000000000	1	INBUF
+ 0000000006	7	INCARD1
+ 0000000007	8	INCARD2
+ 0000000008	9	JBUS
+ 0000000016	10	MINUS
+ 0000000020	22	READINFCARDS
+ 0000000033	35	READTRAJECT
+ 0000000030	32	STORE
+ 0000000021	23	WAITINFCARDS
+ 0000000034	36	WAITTRAJECT

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