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SOME STATISTICS ON ALGOL 68 PROGRAMS

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Some Statistics on ALGOL 68 Programs*)

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ABSTRACT

An attempt is made to assess some static and dynamic properties of ALGOL 68 programs, which are useful for optimization decisions. The results indicate that slicing and assignation are the most important candidates for optimization, and that optimization efforts need to be directed to the simple cases only.

KEY WORDS & PHRASES: compiler construction, optimization, ALGOL 68

*) This report will be submitted for publication elsewhere.

1. THE PROBLEM

For the design of the code generator of the MC ALGOL 68 Compiler we are interested in the frequency of language constructs in normal run-of-the-mill ALGOL 68 programs [1, 2]. Knowledge of these frequencies can guide us on what to optimize, or, if we do not want to optimize now, at least prevent us from making decisions which would rule out useful optimizations later on.

The 'frequency of language constructs in normal programs' not a very precise notion and it is not easy to determine. is There is no good definition of a 'normal program' and we need a full parser to identify and count 'language constructs'.

We can, however, try to get an approximation. Rather than defining 'normal programs' and a distribution, we can take a number of existing real-world programs. ALGOL 68 is used extensively at our installation (Control Data Cyber 72), where 7 % of all compilations are ALGOL 68, so we have the opportunity. And rather than tinkering with the existing compiler (which cannot do) we can do statistical analysis on the texts of we the programs and try to interpret the results.

Much of the philosophy developed by Knuth in his study of FORTRAN programs [3] applies to this work as well.

Similar investigations have been done for ALGOL 60 [4]。 PL/I [5] and COBOL [6].

2. THE STATIC BEHAVIOUR

2.1. Simplifying transformations

We collected 53 real-world user programs (in total 8131 lines) by asking users. These programs were subjected to the following transformations (through editing, UNIX-commands and devious means):

- comments and pragmats were deleted; 1.
- 2. mode- and priority-declarations were removed;
- 3. all tags were replaced by 'tag',
 - all denotations by 'denotation',
 - all user operators by 'user operator',
 - all user mode indications by 'user_mode' and all colons by 'label_token', 'colon_token' (in specifi-cations), 'up_to_token' (in rowers and in trimmers) or 'routine_token', as appropriate;

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- 4. SKIP and NIL were replaced by 'denotation';
- 5. parentheses in parameter-packs in calls were recognized;
- 6. brackets were split in indexers and rowers;
- 7. symbols that come in pairs or triples were taken together (like (), [], IF THEN FI, etc.);
- 8. all different representations of the same operator were taken together (e.g. +:= and PLUSAB), except those for = and EQ.

2.2. Counting symbols

The symbols were then counted and sorted in descending frequency, which yields the following table.

	Table I, Symbol	Count		
17209	tag		104	ABS
5916	denotation		101	<
5200	3		81	ELIF THEN
3457	•		77	CASE IN ESAC
2362	()		76	<=
1892	* 		72	AND
1850	indexer		69	-:=
1488	=		65	STRING
1242	call		65	~
1061	user_mode		59	OR
961			55	IS
793	OF		54	* *
714	*		53	ВҮ
662	REAL		48	ISNT
626	DO OD		46	:
618	+		41	*:=
572	INT		41	>=
547	up_to_token		39	label_token
540	REF		37	e _
501	routine token		37	CHAR
499	то		22	/:=
480	FOR		20	LOC
461	1		13	FILE
413	rower		12	%
396	IF THEN FI		12	+=:
361	UPB		11	OUT
355	/		11	STRUCT
326	PROC		10	SIGN
278	user_operator		9	ELEM
228	FROM		8	MOD
211	+:=		6	ENTIER
202	ELSE		6	%:=
197	VOID		6	ROUND
150	LWB		4	BITS
148	HEAP		4	EXIT
147	BEGIN END		4	UNION
145	OP		2	FLEX

. Table I. Symbol Count

145	WHILE	1	GOTO	
143	>	1	ODD	
132	/=	1	OUSE	IN
121	colon_token	1	REPR	
108	BOOL			

This table gives rise to some observations. The meaning of some symbols is very unclear. Prime example is the = , which may be a dyadic operator or an is-defined-astoken; only profound analysis can tell the difference.

The first two items in the list correspond to loading a value, which can also be considered part of the operator that uses the result; and the next three items are not connected to any semantic action at all in a reasonable implementation. It is true that the semicolon signifies 'voiding' which technically would amount to discarding a result, but in practice no code needs to be generated. The first to require real action is the := . So it might be useful to weed from the list all symbols that are not directly connected to a run-time action (however, the above list does not contain the "invisible" actions in-volved in coercions). This yields:

	Table II, Action Count		
1892	:=	72	AND
1850	indexer	69	-:=
1488	=	65	~
1242	call	5 9	OR
961	-	55	IS
793	OF	54	* *
714	*	53	BY
618	+	48	ISNT
499	ТО	46	:
461		41	*:=
413	rower	41	>=
396	IF THEN FI	37	0
361	UPB	22	/:=
355	/	20	LOC
278	user_operator	12	%
228	FROM	12	+=:
211	+:=	10	SIGN
150	LWB	9	ELEM
148	HEAP	8	MOD
145	WHILE	6	ENTIER
143	>	6	%:=
132	/ =	6	ROUND
104	ABS	1	GOTO
101	<	1	ODD
81	ELIF THEN	1	OUSE IN
77	CASE IN ESAC	1	REPR

3

76. <=

It is tempting to put percentages into this list and say that "13 % of all semantic actions are assignations", but this is meaningful only if all the symbols given above correspond to actions of the same complexity, which is, of course, not true. Our objective is to find constructions which merit our attention in optimization; it is clear that assignations and slicing are the great winners.

Other constructions can be identified which do not show up directly in the tables. One is the 'boolean-enquiry-clause'; its frequency can be found by adding those of IF-THEN-FI, ELIF-THEN, WHILE and a percentage of | (which may represent THEN, ELSE, IN or OUT), and of |: | (which may be ELIF-THEN or OUSE-IN). If we make the only reasonable but totally unwarranted assumption that the brief symbols occur in the same ratio as the bold symbols, we find that 270 |'s are THEN's and 45 |:'s are ELIF's.

Another construction is 'standard-operator', which can be identified but is of doubtful use: the field is too wide for determined optimization. On the other hand, they are so numerous that not identifying them would also give a false impression. We then arrive at the following table.

	Table III, Summary				
4420	standard_operator	228	FROM		
1892	:=	148	HEAP		
1850	indexer	128	CASE	ΙŃ	ESAC
1488	=	55	IS		
1242	call	53	BY		
937	boolean_enquiry	48	ISNT		
793	OF	37	0		
499	ТО	20	LOC		
413	rower	2	OUSE	IN	
278	user_operator	1	GOTO		

The main constructs of interest are assignations, slices and calls. A further analysis (through more editing etc.) is given in the following tables ('simple' means 'identifier or denotation', s.slice means 'slice with simple indexers only', s.selection means 'selection on an identifier' and s.formula means 'formula with one standard operator and one or two simple operands').

Assignations. destination: source: 71 % simple: 45 % simple: s.slice: 15 % 5 % s.slice: 4 % 5 % s.selection: s.selection: s.formula: 8 % 37 % 10 % rest: rest: Slices. primary: indexer: 89 % simple: one, simple: 58 % more, simple: 20 % s.slice: 4 % 4 % 8 % s.selection: trimmer: 3 % rest: rest: 14 % Calls. primary: parameters: 100 % simple: one, simple: 22 % more, simple: 19 % 'print' etc: 17 % 42 % rest:

All this suggests very strongly that it is most efficient to direct the optimization effort to the simple cases only.

2.3. Denotations

The denotations extracted from the text in point 3 in paragraph 2.1 were distributed as follows,

3912	int	194	nil
813	real	148	skip
596	string	18	format
233	bool	2	bits

whereas the integral-denotations were classified thus:

value (range)	freq.
0	601
1	1628
2:3	663
4:15	664
16:255	299
256:4095	53
>4095	6

One conclusion from this is that a reasonable implementation on the IBM 370 may put integers smaller than 4096 in the instruction (LA) and use horrible code for the rest.

2.4. Identifiers

The distribution of identifier-lengths was as follows:

freq.	length	freq.	length
6539	1	29	14
3200	2	11	15
1985	3	12	16
1792	4	14	17
1345	5	2	18
690	6	2	20
731	7	3	21
190	8	3	24
350	9	2	27
168	10	4	34
129	11	1	42
138	12	1	50
52	13	1	52

or, if we consider different identifiers only:

freq.	length	freq.	length
26	1	13	14
295	2	6	15
232	3	4	16
270	4	· 6	17
168	5	2	18
153	6	1	20
102	7	2	21
61	8	2	24
88	9	1	27
69	10	2	34
43	11	. 1	42
35	12	1	50
14	13	1	52

This may provide trade-off information for the identifiertable algorithm.

The 10 most frequent identifiers were:

976	i	359	r
581	n	339	S
564	k	324	Ъ
558	а	305	x
376	j	302	newline

3. THE DYNAMIC BEHAVIOUR

All the above measurements pertain to the static text of the program. We would, however, like to get some insight in the dynamic importance of the various constructs. Now such results are hard to come by and have a inherently large inaccuracy. We therefore decided to accept a static (textual) analysis of the innermost do-parts as a reasonable estimate of the dynamic behaviour of the program, on the (not too well founded) assumption that these parts are the most heavily executed pieces of code.

The same process as above yields the following tables:

Tab	le IV,	Symbol	Count	in	Inner	: Do-p	parts	
4021	tag				13	CASE	IN ESAC	
916	indexe	r			10	ELIF	THEN	
907	denota	tion			9	AND		
879	,				9	<		
504	:=				8	OR		
406	;				7	IS		
355	()				7	/:=		
269	call				7	~	*	
222	*				6	ELEM		
193	OF				6	ISNT		
166	-				6	<=		
159	+				6	rowen	5	
131	=				5	0		
113	+:=				4	UPB		
92	up_to_	token			4	>=		
86	/				3	ENTIE	ZR	
83	IF THE	N FI			3	%		
74	1				3	PROC		
68	user_m	ode			3	1: 1		
44	-:=				2	BEGIN	N END	
39	REF				2	MOD		
37	INT				1	BITS		
35	user_o	perator			1	CHAR		
33	REAL				1	OUT		
30	* *				1	%:=		
28	/=				1	ROUNI)	
27	ELSE				1	STRIM	NG	
24	ABS				1	STRUC	СТ	
22	*:=				1	UNION	N	
22	>				1	+=:		
18	HEAP				1	routi	ine_token	1
14	colon_	token						

Tal	ble V, Action Count in In	nner	Do-parts
916	indexer	10	ELIF THEN
504	:=	9	AND
269	call	9	<
222	*	8	OR
193	OF	7	IS
166	-	7	/:=
159	+	7	~
131	=	6	ELEM
113	+:=	6	ISNT
86	/	6	<=
83	IF THEN FI	6	rower
74		5	Q
44	-:=	4	UPB
35	user operator	4	>=
30	**	3	ENTIER
28	/=	3	%
24	ABS	3	:
22	*:=	2	MOD
22	>	1	%:=
18	HEAP	1	ROUND
13	CASE IN ESAC	1	+=:
Tabi	le VI, Summary of Counts	in]	Inner Do-parts
987	standard_operator	35	user operator
916	indexer	18	HEAP
504	:=	1.3	CASE IN ESAC
269	call .	7	IS
193	OF	6	ISNT

Although the overall picture remains the same, certain shifts in emphasis can be discerned. The slice is now clearly the most important construct, but assignation is still a powerful second. The call has lost much of its weight.

6 rower

5@

Analysis of slice and assignation gives:

146 boolean_enquiry

131 =

Slice in Inner Do-parts.

primary:			indexer:		
simple:	86	%	one, simple:	62	%
s.slice:	7	%	more, simple:	22	%
s.selection:	4	%	trimmer:	0	%
rest:	3	%	rest:	16	%

Ass	signati	ions	in Inner Do-parts.			
destinat	ion:		source:			
simple:	50	%	simple:	32	%	
s.slice:	43	%	s.slice:	9	%	
s.selectior	n: 2	%	s.selection:	9	%	
			s.formula:	37	%	
rest:	-5	%	rest:	13	%	

We see that the assignations tend to have simpler sources now, which again suggests that optimizing the simple cases only will lead to considerable gain. The slices themselves show no real difference.

4. CONCLUSION

The main candidates for optimization efforts are slices, assignations and calls; there are indications that the first two are the most important from a dynamical point of view.

Optimization efforts need to be directed to the simple variants of the above constructions only.

This conclusion is in full agreement with the results obtained by Knuth for FORTRAN [3].

5. ACKNOWLEDGEMENT

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