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A COMPILER COMPILER

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0. Introduction.

The purpose of this report is to give an informal description of a compiler compiler for languages defined by means of a Compiler Description Language based on affix grammars, together with a description of that compiler compiler itself in that Compiler Description Language (CDL). Affix grammars have been described formally in [2].

In the first chapters of this report, an explanation is given of the use of this compiler compiler. This explanation is neither very formal nor very complete, but should enable the reader to understand the description given in the last chapter of the report.

The compiler compiler is a syntax-directed compiler of the top-to-bottom variety equipped with a macro mechanism, which accepts input written in a two-level extension of CF grammar, giving as output a program in ALGOL 60. It is intended as a tool in the study and production of compilers for high level languages, such as ALGOL 68.

0.1 Notation and terminology.

A Context-Free grammar (CF grammar) consists of two distinct finite collections of symbols (, termed the collections of "nonterminal" and "terminal" symbols respectively), and furthermore of one specific non-terminal, its "starting symbol", and a finite collection of "rules".

Each rule associates with some nonterminal its "alternatives", a number of possibly empty sequences of nonterminals and terminals.

In writing down a grammar, we will write down one rule for each non-terminal, with that nonterminal as its "left-hand-side", followed by a colon (:), followed by the "right-hand-side" of the rule, followed by a point (.). The right-hand-side consists of the alternatives of the left-hand-side, each separated from the next by a semicolon (;). The symbols within one alternative are separated by commas (,). As symbols we will use only "tags", i.e., sequences of letters and digits beginning with a letter.

The notation just described, the "Van Wijngaarden notation" [3], is a variation on the more common Backus-Naur form [4], with as advantages ease of writing and reading, and the fact that the end of a rule is clearly indicated.

By a "direct production" of a nonterminal we mean any of its alternatives. By a "production" of a nonterminal x we mean a sequence of symbols, which is either a direct production of x , or is the result of replacing, in some production of x , a nonterminal y by a direct production of y .

By a "terminal production" of a nonterminal x we mean a production of x that contains no nonterminals.

By a "sentence" of a CF grammar we mean a terminal production of its starting symbol.

By the "language" of a CF grammar G we mean the collection of its sentences. This language is denoted by $L(G)$.

0.2 CF grammars of type LL(1).

By the productions of a sequence of symbols y we mean the productions of a new nonterminal x with y as its only direct production, with the exclusion of y itself.

For an alternative x , we will indicate by $\text{first}(x)$ the set of terminal symbols which are the first symbol of some production of x , and by $\text{follow}(x)$ the set of terminal symbols which immediately follow x in any production of the starting symbol.

We then say a CF grammar is "of type LL(1)" [6] if it satisfies the following 3 conditions:

- 1) For any rule $A: a_1; a_2; \dots; a_n$, the sets $\text{first}(a_1)$, $\text{first}(a_2)$, \dots , $\text{first}(a_n)$ are mutually disjoint.
- 2) At most one of a_1, \dots, a_n can produce an empty string ϵ .
- 3) If a_p produces ϵ , then $\text{first}(a_q)$ has no elements in common with $\text{follow}(A)$.

It is decidable whether a given CF grammar is of type LL(1). Without loss of generality we may assume that only the last alternative can produce ϵ (condition 2).

1. Syntax-directed compiling.

1.1 What is a compiler compiler?

A "program" is a text in some programming language more or less directly executable on a computer, which, when executed with some input, gives some output.

A "compiler" for some high level language is a program which, when executed with a program in that language as input, gives as output a program in some other language (preferably more directly executable).

A "compiler compiler" is a compiler which, when executed with a formal description of a compiler for some language as input, gives as output a compiler for that language.

A "family of compiler compilers" is a collection of compiler compilers which, when executed with identical input, give as output equivalent programs, possibly in different languages.

Ideally, a formal description of some programming language should be capable of yielding also a formal description of a compiler for it; certainly a formal description of a compiler may be used to define a language. The formalism for definition of a language consists of two coupled mechanisms, its syntax and semantics. Strictly speaking, the syntax should indicate whether a given sequence of symbols is a program in the language, and if so, derive a parsing tree for it, while the semantics should attach a meaning to the parsing tree of a correct program. In the definition of most languages, for instance that of ALGOL 60, this separation is not so clean cut: one can argue that part of the syntax is given under the heading of semantics. In ALGOL 68, the "context-conditions" are an example of syntactical restrictions on programs appearing as semantics, so that the syntax by itself is not sufficient to define what is a proper program.

Formalization of the syntax of programming languages is now well enough understood to obtain part of a compiler by mechanical means, using "syntax-directed" techniques. On the other hand, there exists no formalization of the semantics which could usefully give the "semantics-directed" parts of a compiler. For this, manual techniques are used, for instance by translating to some set of macros which then reflect the semantics.

Our goal is to extend the applicability of syntax-directed techniques so far that all but the very kernel of the semantics of a compiler can be defined by syntax, so that only that kernel remains to be defined by other means. To this end, a notational system, "Compiler Description Language" (CDL) will be introduced, strongly based on an extension of CF grammars ("affix grammars"), a kind of grammars which, on the one hand, are powerful enough to define major programming languages, and, on the other hand, lend themselves to syntax-directed compiling techniques.

In the light of the foregoing discussion the Compiler Description Language can best be seen as a programming language, which allows one, at a high level of abstraction, to write compilers in a machine independent way.

Although it is our intention to describe a family of compiler compilers, the discussion will center around one particular member of that family, whose input is in Compiler Description Language, and whose output is in ALGOL 60. ALGOL 60 as an output language provides a clear enough framework for the notions involved while being broadly available. ALGOL 68 would certainly provide a better framework, but its availability is still very limited.

Other members of the same family of compiler compilers are being constructed, giving output in various assembly languages and in ALGOL 68.

1.2 Parsing according to a CF grammar.

The basis of a syntax-directed compiler is a parser, so we must first go into matters of parsing.

By a "parse" of a sentence according to a grammar we mean a sequence of substitutions of a direct production for some nonterminal, which leads from the starting symbol of the grammar to that sentence. The problem of parsing a given sequence of symbols then consists of

- 1) deciding whether that sequence of symbols is a sentence of the grammar;
- 2) if so, then constructing a parse for it.

Using a CF grammar as a generating device is simple enough, but it takes more trouble to find a practical algorithm for parsing, even though

the parsing problem can easily be shown to be decidable. Still, at the moment the parsing problem for CF grammars is wholly solved, and literature abounds with parsing methods.

These methods can be classified along various criteria:

- a) bottom up or top down
- b) general algorithm & table for the particular grammar or particular parser for that grammar only
- c) parser for general CF grammars or parser for a reduced class obtained by imposing restrictions on the grammars.

The compiler compiler to be described generates a top down particular parser, written in ALGOL 60, for a CF grammar which satisfies the LL(1) conditions.

From a grammar, the particular parser is obtained by a transcription process where each rule is turned into a <declaration> for a parsing procedure corresponding to the nonterminal in its left-hand-side; these procedures are then embedded in an environment, which contains at least some as yet unspecified means of input and output. Several such transcriptions are possible, and we will investigate two of them.

1.2.1 The first parsing method.

The transcription process is best introduced by an example.

Assume the following part of a grammar:

```
number: digit, numbertail.                                G1
numbertail: digit, numbertail; endmarker.
```

Its transcription would run:

```
Boolean procedure number;                                P1
begin if  $\neg$  digit then goto 1;
      if  $\neg$  numbertail then goto 0;
      number := true; goto end;
1:0: number := false; end;
end;
```

```

Boolean procedure numbertail;
begin if  $\neg$  digit then goto 1;
    if  $\neg$  numbertail then goto 0;
    numbertail := true; goto end;
  1: if  $\neg$  endmarker then goto 2;
    numbertail := true; goto end;
  2: 0: numbertail := false; end;
end;

```

We will assume *digit* and *endmark* to be procedures declared in the environment, such that they return true if the current symbol in the input-stream is a digit or an endmark respectively, simultaneously advancing the input by one symbol, and return false otherwise.

When we now call *number* it returns either true or false, depending on whether the input stream, starting with the current symbol, contained a sequence of digits followed by an endmarker.

Now consider another CF grammar for number 1:

```
number 1: digit, number 1; digit, endmarker. G2
```

with transcription:

```

Boolean procedure number 1; P2.a
begin if  $\neg$  digit then goto 1;
    if  $\neg$  number 1 then goto 0;
    number 1 := true; goto end;
  1: if  $\neg$  digit then goto 2;
    if  $\neg$  endmarker then goto 0;
    number 1 := true; goto end;
  2: 0: number 1 := false; end;
end;

```

The first transcription method fails to parse any number 1. Consider, e.g., *9 #*, where *9* is a digit and *#* the endmarker. It finds the string begins with a digit, so advances the input, and then tries to parse the remainder as a number 1, which fails; so number 1 returns false.

Clearly, there are CF grammars for which method 1 does not work satisfactorily.

1.2.2 A second parsing method.

The trouble with grammar G2 had to do with the fact that, upon meeting a digit, the input could not safely be advanced. Assume the environment to contain an input pointer *pin*, indicating at every moment the current symbol of an input array. Advancing the input is done by incrementing *pin*, so backtracking the input can be done by restoring *pin* to its original value.

The second transcription method transforms G2 into

```

Boolean procedure number 1;                                P2.b
begin integer pold; pold:= pin;
    if  $\neg$  digit then goto 1;
    if  $\neg$  number 1 then goto 1;
    number 1:= true; goto end;
1: pin:= pold;
    if  $\neg$  digit then goto 2;
    if  $\neg$  endmarker then goto 2;
    number 1:= true; goto end;
2: pin:= pold; number 1:= false; end:
end;

```

This second version will recognize the string 9 # as follows: it finds 9 is a digit and increments *pin*; it finds # is not a number 1, so control goes to the label 1; *pin* is restored to its old value; a digit is found, incrementing *pin*; an endmarker is found, again incrementing *pin* and the procedure returns true.

Clearly this second transcription method works for a wider class of grammars than the first did.

1.2.3 Comparison of the two methods

In this section we will compare the merits of the two methods and the conditions under which they can be used.

First some terminology: We will term a procedure x , yielding a Boolean value, and obtained from a nonterminal y by one of the transcription processes, a "parsing procedure" for y .

We say a parsing procedure x "recognizes" the nonterminal x if it satisfies the following two requirements:

- 1) x always terminates;
- 2) if the inputstream, starting with the current symbol, contains a terminal production of x , then x yields true, simultaneously advancing the input to the first symbol after that terminal production; otherwise it yields false.

We will say x "exactly recognizes" x when x recognizes x and advances the inputstream only when returning true.

In this terminology *digit* is a parsing procedure exactly recognizing a digit. By induction on the length of the inputstream, one can deduce that *number* and *numbertail* recognize number and numbertail, though not exactly, as shown by the inputstring $1a$ where a is not an endmarker: number 1 will return false, but has then already advanced the input.

According to the first method, number 1 (P2.a) fails to recognize anything, but the second version (P2.b) exactly recognizes number 1.

For convenience we will term the first transcription method "non-restoring" and the second "restoring".

We can make a number of observations concerning the two transcription methods:

- 1) For a CF grammar its nonrestoring parser recognizes its starting symbol if and only if that grammar satisfies the LL(1) conditions. The first condition assures that, upon recognizing a terminal symbol, the input can safely be advanced without a chance of taking a wrong alternative. The second and third condition assure that no problems arise because of empty terminal productions.
- 2) For a CF grammar, its restoring parser exactly recognizes its starting symbol if that grammar satisfies the LL(1) conditions. These conditions assure that an alternative taken and found present is indeed the only correct one. They are sufficient, but not necessary: the class of grammars recognized exactly contains the LL(1) grammars.

The difference between the restoring and nonrestoring parser for a LL(1) grammar comes out only for input which does not form a sentence: while the restoring parser rejects it (recognizing exactly), the non-restoring parser may reject it but still have advanced the input, or even accept it, as is clear from the example:

```
statement: ass stat; if stat.                                G3.a
ass stat:  identifier, becomes symbol, expression.
if stat:   if symbol, bool expr, then symbol, statement.
```

With *statement* as starting symbol, and appropriate exact recognizers for *identifier*, *becomes symbol*, etc., the nonrestoring parser for G3 would accept not only

```
      v:= 0
and   if v > 2 then v:= v + 1
but also v if v > 2 then v:= v + 1
```

without giving any error message. A nonrestoring parser may accept erroneous sentences without warning. This deficiency must be mended by explicitly taking possible errors into account while constructing the syntax:

```
ass stat:      identifier, rest ass stat.                    G3.b
rest ass stat: becomes symbol, expression; errormessage.
```

Errormessage should print out some appropriate warning.

Advantages of the nonrestoring parser over the restoring parser are:

- 1) Time-efficiency: the parser has less work to do than the restoring parser, which may have to perform lots of backtracking.
- 2) Memory-efficiency: because backtracking can never occur, there is no need for any array to hold all of the input, and no need for temporary storage to retain old values of *pin*.

The disadvantage of the nonrestoring parser lies in the restrictions it imposes on the grammar.

In a practical situation, the advantages of the nonrestoring parser may outweigh its disadvantage: when constructing a syntax, one has good

grounds to make it LL(1).

The compiler compiler allows one to choose between the two parsing methods, translating some rules in a restoring, others in a nonrestoring fashion. This facility is essential for making compilers as efficient as possible, but makes explaining the workings of the compiler compiler somewhat complicated. For this reason we will show mainly nonrestoring parsers in the rest of this chapter, leaving the construction of the restoring version to the reader.

1.3 Extensions to CF syntax.

In this section we will give an account of the extensions that have to be made to CF syntax in order to turn it into an appropriate input-language for the compiler compiler.

1.3.1 Embedding actions into the syntax.

In order to turn a parser into a compiler, one has to provide it with means to perform, as a side-effect of parsing, some semantic actions. To accomplish this, these actions are embedded in the syntax as follows: Consider the parsing of a number with as translation its value:

```
number 2: digit, action 1, numbertail 2.                                G4
numbertail 2: digit, action 2, numbertail 2; endmarker.
```

where *action 1* assigns the value of the last digit read to some global variable *s* while *action 2* multiplies *s* by 10 and then adds the value of the last digit read to *s*.

We divide the nonterminals into "predicates" and "actions", where predicates are transcribed as described in 1.2, while actions are transcribed as insertion of a procedure call at the corresponding place of the parser:

```

Boolean procedure number 2;
begin if  $\neg$  digit then goto 1;
    action 1;
    if  $\neg$  numbertail 2 then goto 0;
    number 2 := true; goto end;
1:0: number 2 := false; end:
end;

```

```

Boolean procedure numbertail 2;
begin if  $\neg$  digit then goto 1;
    action 2;
    if  $\neg$  numbertail 2 then goto 0;
    numbertail 2 := true; goto end;
1: if  $\neg$  endmarker then goto 2;
    numbertail 2 := true; goto end;
2:0: numbertail 2 := false; end:
end;

```

We will allow "primitive" actions and predicates to be defined not by a rule of the grammar but by some other means, e.g., macros. If we take as such primitive actions:

```

alpha = s := 0
beta  = s := 10 * s
and gamma = s := s + last digit read

```

then we can define the actions action 1 and action 2 in terms of those primitive actions:

```

action 1: alpha, gamma.
action 2: beta, gamma.

```

with transcription:

```

procedure action 1; begin s := 0; s := s + last digit read end;
procedure action 2; begin s := 10 * s; s := s + last digit read end;

```

An example of a primitive predicate could be digit.

The transcription of a rule for an action is the <declaration> of a procedure (not a Boolean procedure) which may involve again actions and predicates, primitive or not primitive. More involved examples will follow.

1.3.2 Affixes.

The primitive actions alpha, beta and gamma mentioned in 1.3.1 are neither really primitive nor very practical: they involve a global variable *s* which may not be used for other purposes by other actions, e.g., some action taken by digit.

We want to provide parsing procedures with parameters, as a mechanism for communicating information to and from other parsing procedures.

To this end, a nonterminal may be accompanied by a number of "affixes" in the form of symbols written after it in the syntax and each preceded by a plus (+), e.g.:

```
number 3 + value: G6
    digit, action 3, numbertail 3 + value.
numbertail 3 + value:
    digit, action 4, numbertail 3 + value;
    endmarker.
```

Where action 3 and action 4 are primitive actions:

```
action 3 = value := last digit read
and action 4 = value := 10 * value + last digit read
```

with transcription:

```
Boolean procedure number 3 (value); integer value; P6
begin if ¬digit then goto 1;
    value := last digit read;
    if ¬numbertail 3 (value) then goto 0;
    number 3 := true; goto end;
1:0: number 3 := false; end;
end;
```



```

Boolean procedure numbetail 3 (value); integer value;
begin if  $\neg$  digit then goto 1;
    value := 10 * value + last digit read;
    if  $\neg$  numbetail 3 (value) then goto 0;
    numbetail 3 := true; goto end;
1: if  $\neg$  endmarker then goto 2;
    numbetail 3 := true; goto end;
2:0: numbetail 3 := false; end:
end;

```

Note that *value* is an output parameter. All parameters are called by name. Of course, also primitive actions and predicates can have parameters.

Apart from parameters, we will also need local variables, which we will indicate in the left-hand-side of the rule as symbols each preceded by a minus (-), e.g.:

```

number 4 + value - d:                                     G7
    digit 4 + d, action 5 + value + d, numbetail 4 + value.
numbetail 4 + value - d:
    digit 4 + d, action 6 + value + d, numbetail 4 + value;
    endmarker.
digit 4 + d: digit, action 5 + d + last digit read.

```

where

```

action 5 + x + y = x := y
and action 6 + x + y = x := 10 * x + y

```

with transcription

```

Boolean procedure number 4 (value); integer value;                                     P7
begin integer d;
    if  $\neg$  digit 4 (d) then goto 1;
    value := d;
    if  $\neg$  numbetail 4 (value) then goto 0;
    number 4 := true; goto end;
1:0: number 4 := false; end:
end;

```

```

Boolean procedure numbertail 4 (value); integer value;
begin integer d;
    if  $\neg$  digit 4 (d) then goto 1;
    value := 10 * value + d;
    if  $\neg$  numbertail 4 (value) then goto 0;
    numbertail 4 := true; goto end;
    1: if  $\neg$  endmarker then goto 2;
    numbertail 4 := true; goto end;
    2: 0: numbertail 4 := false; end:
end;

```

```

Boolean procedure digit 4 (d); integer d;
begin if  $\neg$  digit then goto 1;
    d := last digit read;
    digit 4 := true; goto end;
    1: digit 4 := false; end:
end;

```

Those symbols which occur in the grammar to indicate a parameter or local variable are affixes, and must be distinct from all other symbols.

Affixes occurring in the left-hand-side of a rule we term "bound" affixes of that rule if they are preceded by a plus and "free" affixes if preceded by a minus. Affixes occurring in a rule which are neither bound nor free are "terminal".

For sake of simplicity we assume the bound, free and terminal affixes to be disjoint sets. If a bound affix is used as an input parameter to a parsing procedure we will term it "inherited"; otherwise, it is "derived".

A nonterminal must, when occurring in a right-hand-side, always be accompanied by as many affixes as there are bound affixes in the left-hand-side of its defining rule.

Of course, care has to be taken that an inherited affix should get a value before being used. In an article of D.E. Knuth [5] it is shown how to investigate by graph-theoretical means whether a system of inherited and derived affixes is consistent (in his terminology, whether a system of inherited and synthesized attributes is well defined).

A nonterminal together with its affixes is termed an "affix expression", of which the nonterminal is the "handle". The handle with its appended affixes, in a definition of an affix grammar as a generating device, takes the same place as a nonterminal in the definition of a CF grammar. Affixes derive their name from the fact that they are considered to be attached to the handle of an affix expression.

1.3.3 Grouping; jumps and labels.

Consider again G2:

number 1: digit, number 1; digit, endmarker.

When recognizing a number 1, a restoring parser for this grammar has to recognize the the last digit twice. The grammar can be rewritten to G1 in order to obviate this backtracking, but then we must introduce an extra nonterminal numbertail. When we have recognized the first digit of a number 1, we expect either an endmarker or another number 1; this could be denoted as

number 5: digit, G8
(number 5; endmarker).

with as transcription:

Boolean procedure number 5; P8.a
begin if \neg digit then goto 1;
 begin if \neg number 5 then goto 2;
 number 5:= true; goto end;
 2: if \neg endmarker then goto 3;
 number 5:= true; goto end;
 3: goto 0;
 end;
1:0: number 5:= false; end:
end;

The restoring version of this parser runs:

```

Boolean procedure number 5;                                P8.b
begin integer pold; pold:= pin;
  if  $\neg$  digit then goto 1;
  begin integer pold; pold:= pin;
    if  $\neg$  number 5 then goto 2;
    number 5:= true; goto end;
  2: pin:= pold;
    if  $\neg$  endmarker then goto 3;
    number 5:= true; goto end;
  3:
  end;
  1: pin:= pold; number 5:= false; end:
end;

```

As a restoring parser, P8.b recognizes the same language as P2.b but without recognizing the last digit twice. As a nonrestoring parser, P8.a succeeds where P2.a failed; consequently G8 can be considered as of type LL(1). The introduction of grouping brackets allows one to construct grammars which can be parsed more efficiently, and to enlarge the applicability of nonrestoring parsers.

Still, the parser for number 5 is inefficient because it is an unnecessarily recursive procedure; the recursion in number 5 is used only to effectuate repetition: a number consists of a first digit, optionally followed by some number of digits, followed by an endmarker. We will need a device for performing repetition. Therefore we will allow both labels and jumps within one rule of a grammar, indicating a label as a symbol acting as label-identifier followed by a colon, and a jump by a colon followed by such a symbol, e.g. (with label *rep*):

```

number 6:                                                    G9
  digit,
  rep: (digit, :rep; endmarker).

```

with transcription:

```

Boolean procedure number 6;
begin if  $\neg$  digit then goto 1;
  rep: begin if  $\neg$  digit then goto 2;
    goto rep;
    2: if  $\neg$  endmarker then goto 3;
      number 6 := true; goto end;
    3: goto 0;
  end;
  1:0: number 6 := false; end;
end;

```

P9

A jump is an action and may only occur at the end of some alternative, to a label within that same rule.

Under these conditions, the introduction of jumps and labels can easily be understood as shorthand for a corresponding recursive definition, so that the language accepted by a parser is unchanged, and none of the problems associated with the unrestricted introduction of jumps and labels is encountered. As an example, the recursive counterpart of number 6 is simply:

```

number 7: digit, rep.
rep: digit, rep; endmarker.

```

1.3.4 Some more syntactic sugar.

In order to turn affix grammars into a Compiler Description Language, some more features are needed.

Apart from parameters and local variables, parsing procedures may require global variables. The variable *pin* used by restoring parsers is an example. For the sake of simplicity we will use integer variables only, as a very general and useful datatype. Thus, for a pointer *pin*, a <declaration> will appear in the parser, as

```
integer pin;
```

Furthermore, a facility is needed to declare arrays, e.g. for holding input, and, as we will see later, to be used as stacks. Lastly, a facility

is needed for declaring a Boolean variable, termed a "flag", which may then occur instead of a predicate. In 2.4 the precise notation for such declarations is given. In 2.3 it is described how one can indicate to the compiler compiler that a specific symbol is an action or predicate, defined globally, and thus to be used without further definition.

A simple mechanism for specifying macros is also described in 2.4.

1.4 What is the use of a compiler compiler?

The conventions introduced in the previous sections, and described in chapter 2, allow one to program the input to the compiler compiler in a reasonable level detail giving a reasonable efficient compiler, where especially one can leave many administrative details to the system. Of course, ALGOL 60 is not the most practical language to write compilers in. But also in the hardest machine code, subroutines (with local variables and parameters kept on a stack), global variables, labels and jumps etc. have their counterparts, so a compiler compiler can be written towards every decent machine code.

The primitive actions and predicates are not defined in the Compiler Description Language itself, but are borrowed from some other language in the form of macros. These primitive actions and predicates therefore have to be chosen with particular care, so that they can be implemented in an equivalent way on any reasonable computer.

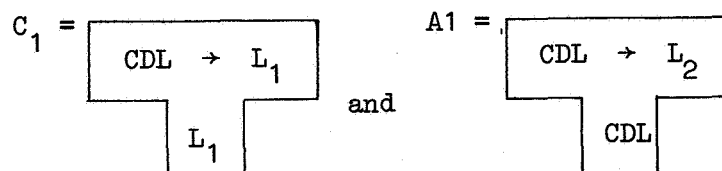
The Compiler Description Language allows one to define, in a machine independent fashion, how to construct recursively more involved actions and predicates out of the primitive ones, and out of other actions and predicates.

A compiler description in CDL describes not one single compiler, but a collection of these, one for every language to which a compiler compiler exists.

In particular, because the compiler compiler itself is defined in CDL (chapter 4) it is possible to obtain a compiler compiler to another language, and written in that other language, by a process of bootstrapping [7]:

Preparation Assume we have a compiler compiler C_1 , written in L_1 , which turns a compiler description (= text in CDL) into a compiler written in L_1 again. We then construct a compiler description A1 of a compiler compiler which turns a compiler description into a compiler written in L_2 .

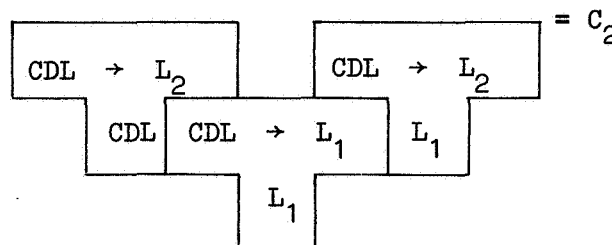
In pictures:



Step 1

We execute C_1 with A1 as input, giving as output C_2 , which is a compiler compiler written in L_1 , turning a compiler description into a compiler written in L_2 .

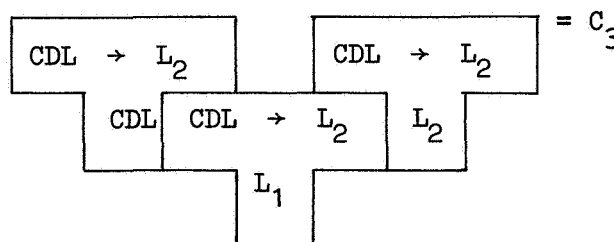
In a picture:



Step 2

We execute C_2 with A1 as input, giving as output C_3 , which is a compiler compiler written in L_2 , turning a compiler description into a compiler written in L_2 .

In a picture:



Using CDL for writing some compiler is preferable both to using machine-code and to using a general purpose high-level language as, e.g., ALGOL 60, for the following reasons:

- Compiler Description Language forms a very high framework for writing compilers, where one is little bothered with the administrative details of the parsing process.

- A Compiler Description is highly machine-independent and transferable, since a compiler can be obtained in any object language for which a compiler compiler exists. Because of their simplicity, compiler compilers for CDL can easily be made compatible.
- There is the possibility of first testing the compiler description using a high-level object code with good diagnostics, which is guaranteed never to blow up the system, and later to use a version in machine-code for production.
- By forcing oneself to use only specific programming tools, there is less possibility for tricky programming, using local features of some system which make transfer impossible.
- A Compiler Description is well enough readable to be of help in documentation.

Of course a carefully handwritten program can be more efficient both in time and in space than the mechanically obtained result of a compiler compiler. But one must not exaggerate the importance of this point: by a suitable choice of macros and careful construction of the crucial part of the syntax, one can get as near to machine-code as one should wish.

On a more general level, the Compiler Description Language can be seen as a mechanism for the syntactical composition of semantics, which should be suitable for the definition of programming languages, not by an in principle generative grammar plus verbally formulated semantics, but as a system describing the meaning of a program in terms of very simple and basic semantics actions, i.e., a machine-independent compiler.

2. About this compiler compiler.

In this chapter, a description is given of the Compiler Description Language, in the form of a context-free grammar exhibiting its syntax, accompanied by semantic remarks. For a precise definition, the description of the compiler compiler itself is to be found in chapter 4.

2.1 Compiler description.

compiler description:

specification, compiler description;
 declaration, compiler description;
 command, compiler description;
 comment, compiler description;
 rule, compiler description;
 starting symbol.

A compiler description consists of various building stones; some of them provide only information to the compiler compiler (as specifications and commands), others give rise to a translation. The compiler compiler treats those building stones in one scan as it goes along, at the same time printing out a copy of the text with possibly some diagnostic messages, producing a translation and collecting information, some of which is displayed afterwards as a diagnostic aid (see section 2.9).

Since all work is done in one forward scan, as a general rule defining occurrences of symbols must precede applied occurrences: if one wants to use, e.g., some symbol as an action, one has to specify it as such before its defining rule. Very limited use is made of block structure: there are only two levels; symbols are either global, or local to some rule; in the latter case they may be redefined in a later rule.

2.2 Symbols.

The input is seen as a sequence of symbols, between which layout characters are ignored. A symbol is of one of four kinds:

- 1) tag, consisting of a letter, possibly followed by a number of letters or digits, between which spaces are ignored.
- 2) constant, a sequence of digits, between which spaces are ignored.
- 3) spec, one of the following reserved special symbols:

plus	+
minus	-
times	*
semicolon	;
comma	,
sub	[
sup]
open	(
close)
colon	:
point	.
equals	=

- 4) bold, a sequence of characters other than accents, enclosed between accents. The following bolds are reserved bold symbols:

external symbol	'external'
action symbol	'action'
predicate symbol	'predicate'
pointer symbol	'pointer'
flag symbol	'flag'
macro symbol	'macro'
list symbol	'list'
restore symbol	'restore'
unrestore symbol	'unrestore'
short symbol	'short'
long symbol	'long'
trace symbol	'trace'
untrace symbol	'untrace'
first parameter symbol	'1'
second parameter symbol	'2'

third parameter symbol '3'
 fourth parameter symbol '4'
 fifth parameter symbol '5'
 result symbol 'result'

Tags are used as nonterminals, terminals, affixes, labels, in macros, etc. Constants may occur, e.g. as inherited affixes and in macros. Non-reserved bold symbols may occur in macros.

2.3 Specifications.

specification:

external specification;
 internal specification;
 macro specification.

external specification:

external symbol, type, tag list, point.

type:

action symbol; predicate symbol;
 pointer symbol; flag symbol.

tag list:

tag, (comma, tag list;).

internal specification:

type of internal, tag list, point.

type of internal:

action symbol; predicate symbol.

macro specification:

macro symbol, type, rest macro spec, point.

rest macro spec:

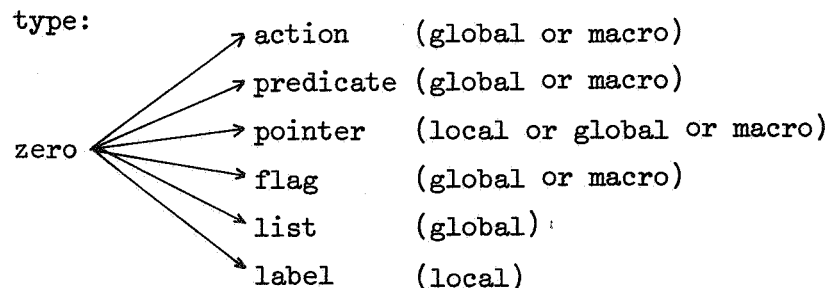
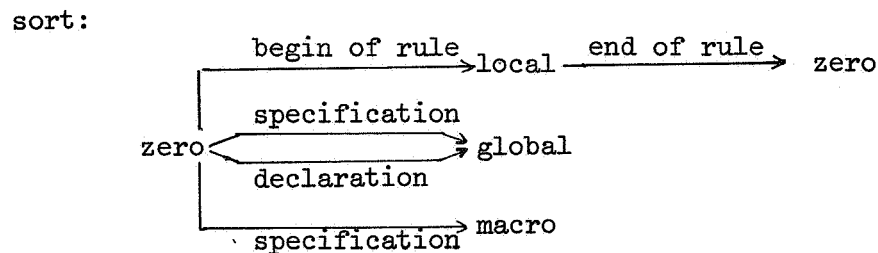
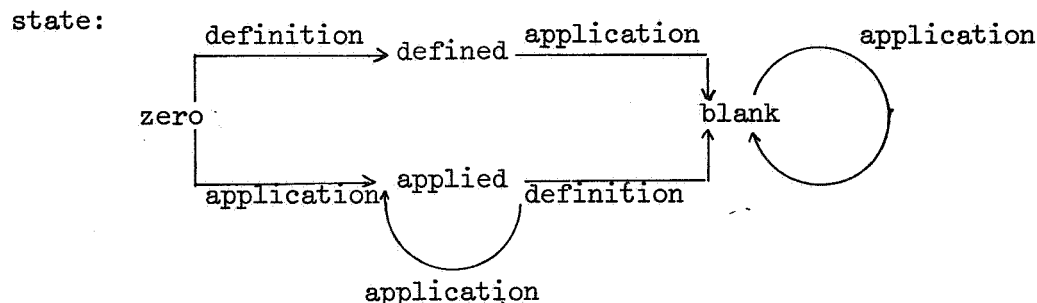
macro def,
 (comma, rest macro spec;).

macro def:

tag, equals, macrotext.

During execution of a compiler compiler with a compiler description as input, every tag occurring in the compiler description has three attributes, its "state", "sort" and "type", which may change according to fixed rules from occurrence to occurrence. Before its first occurrence, the attributes are all assumed to be zero. At every occurrence, the context may prescribe a state, sort and type, and an attempt is made to change the attributes of the tag accordingly. If such a change is not allowed, a diagnostic message is printed. Thus, it is not allowed to redefine a predicate, or use it as an affix, etc. Affixes local to a rule get, at the end of that rule, attributes zero, so that they can be used afterwards with other attributes. The attributes also serve to influence the output of the compiler compiler, so that it translates an action in a different way than a predicate.

Allowed changes of the attributes are



At the end of each rule, the state of each local should be blank (i.e. both defined and applied; if not then an error message is given); all locals again obtain state, sort and type zero.

At the end of the compiler description, the state of each global and macro should be blank, with the exception that an applied but not defined global pointer is treated as a terminal symbol, i.e. a declaration is produced for it, and also a read statement. (See 2.8.)

Specifications serve to prescribe the attributed of a tag, as follows:

		<u>state</u>	<u>sort</u>	<u>type</u>
'external'	'action'	defined	global	action
'external'	'predicate'	defined	global	predicate
'external'	'pointer'	defined	global	pointer
'external'	'flag'	defined	global	flag
'action'		applied	global	action
'predicate'		applied	global	predicate
'macro'	'action'	defined	macro	action
'macro'	'predicate'	defined	macro	predicate
'macro'	'flag'	defined	macro	flag
'macro'	'pointer'	defined	macro	pointer

An external specification indicates that the tags mentioned are supposed to have been declared externally by other means; such a tag may not be redefined and must be applied later on in the compiler description. In this way, e.g., system procedures can be used.

An internal specification indicates that the tags mentioned are to be defined later in in a specific way; it serves as an application: the tags mentioned are henceforth applied globals. Example: 'action' x.

If now a rule follows with x in the left-hand-side, then it will be translated as a <declaration> for an action-procedure. Had the rule not been preceded by this specification, then x would have been treated as a global predicate: an internal specification 'predicate' is always assumed by default.

A macro specification not only establishes the tags mentioned as

defined macros of some type, but also stores their macrotext. A macrotext is any sequence of symbols not containing a comma or a point.

Examples of parameterless macros:

```
'macro' 'pointer' max list = 5000.
```

Whenever now in the text `max list` appears, it will be expanded to `5000`.

If a macrotext must contain a point or comma, then this symbol can be turned into a different (bold) symbol, by enclosing it between accents:

```
'macro' 'pointer' pi = 3'.141592654.
```

On output, the accents will disappear. If one wants to output an accent, then the quote may serve as accent-image, being output as an accent. (There is, unfortunately, no quote image, so a macrotext cannot contain a quote on output). Example:

```
'macro' 'flag' true = '"true"'.
```

which, expanded on output, gives `'true'`.

Macros can also be parameterized, with up to 5 parameters. (Another upper limit might have been chosen.) In the macrotext, the reserved bold symbols `'1'`, `'2'`, `'3'`, `'4'` and `'5'` stand for first, second, up to fifth parameter respectively. For a parameter may be substituted any affix, i.e. a local, global or macro, pointer, flag or list. In 2.7 is described how and when a macro is expanded. The macro system described here is admittedly simpleminded and limited; another system might be used instead.

2.4 Declarations.

declaration:

pointer declaration;

flag declaration;

list declaration.

pointer declaration:

pointer symbol, tag list, point.

flag declaration:

flag symbol, tag list, point.

list declaration:

list symbol, rest list declaration.

rest list declaration:

tag, sub, expr, comma, expr, bus,
(comma, rest list declaration; point).

expr: tag, rest expr; constant, rest expr.

rest expr: plus, expr; minus, expr; .

A declaration has an effect on the attributes of the tags mentioned, making them defined globals.

Furthermore, a pointer declaration is translated as a <declaration> for a number of integer variables, a flag declaration as one for a number of Boolean variables, and a list declaration as one for a number of integer arrays.

Macros occurring in an expr are expanded. Example:

'macro' 'pointer' min text = 1, max text = 10000.

'list' text [min text : max text].

'pointer' stackpointer.

with translation:

integer array text [1 : 10000];

integer stackpointer;

After these declarations, text is a defined global list, min text is a macro pointer (both defined and applied), and stackpointer a defined global pointer.

2.5 Commands.

command: restore symbol; unrestore symbol;

short symbol; long symbol;

trace symbol; untrace symbol.

The command 'restore' causes the compiler to translate rules in the restoring mode, 'unrestore' in the nonrestoring mode. The default

mode is nonrestoring.

The command 'short' causes the compiler compiler to produce compact output by suppressing most layout; 'long' causes more legible output to be produced. The default command is 'long'.

The command 'trace' causes, at the end of each rule, and at the end of the compiler description, tracing output to appear indicating attributes of all locals and globals respectively; the normal situation is 'untrace'.

2.6 Comments.

comment: sub, rest comment.

rest comment: bus; nonbus, rest comment.

(Here, nonbus stands for any symbol except].)

A comment is translated as <comment>, e.g. [input] is translated as comment input;

2.7 Rules.

rule: left hand side, middle, right hand side, point.

left hand side:

handle, optional bound affixes.

handle: tag.

optional bound affixes:

plus, bound affix, optional bound affixes;

times, bound affix, optional bound affixes; .

bound affix: tag.

middle:

optional free affixes, colon.

optional free affixes:

minus, free affix, optional free affixes; .

free affix: tag.

right hand side:

alternative,

(semicolon, right-hand-side;); .

alternative:
 affix expression,
 (comma, alternative;);
 group;
 jump; .
 affix expression:
 label, colon, affix expression;
 handle, optional affixes.
 label: tag.
 optional affixes:
 plus, affix, optional affixes; .
 affix: tag; constant.
 group: open, right hand side, close.
 jump: colon, label.

A rule is translated as a <declaration> for a parsing procedure corresponding to its handle, where the left-hand-side furnishes the <procedure heading>.

Bound affixes that occur preceded by a plus appear in the <specification part> specified integer, whereas those preceded by a times remain unspecified (so as to allow, e.g., arrays to be passed as parameters).

The middle of a rule is translated as begin followed by <declarations> for the free affixes, if any, possibly followed by integer pold; pold:= pin;.

The translation of the right-hand-side consists of the translation of its alternatives, followed by

0: handle:= false; end:
end;

where *handle* is an <identifier> corresponding to the handle of the rule. The <label> *0* labels the only way in which the parsing procedure can return false.

The translation of an alternative consists of the translation of its constituents, if any, followed by *handle:= true; goto end;* followed by an unique <label>.

The translation of a jump is the corresponding <go to statement>. A jump furnishes an applied occurrence of a local label.

The translation of a group is

- 1) In the restoring mode:

begin integer pold; pold:= pin; followed by the translation of its constituent alternatives, followed by end;

- 2) In the nonrestoring mode:

begin followed by the translation of its constituent alternatives, followed by goto 0 ; end;

The translation of an affix expression depends on sort and type of its handle, as follows:

Let λ stand for:

- i) in the restoring mode: the <label> at the very end of the alternative of which the affix expression is a constituent (termed the current <label>);
- ii) in the nonrestoring mode: if this affix expression is the first one of the alternative with a predicate as handle, then the current <label>, and, otherwise, the <label> 0.

Let α stand for:

- i) if the handle is a flag, then the <identifier> corresponding to it;
- ii) if the handle is global, then a call of the procedure corresponding to it, with its affixes translated as <actual parameter>s;
- iii) if the handle is a macro, then its expansion, with its first, second, etc. affix substituted for the corresponding parameter in its macro text. If the handle is a flag or predicate, then this expansion is furthermore enclosed between brackets (and).

For an affix occurring as a parameter. the following translation is given:

- i) if the affix is a constant, then the corresponding <unsigned integer>;
- ii) if the affix is a tag but not a macro, then the corresponding <identifier>;
- iii) if the affix is a macro, then its expansion (without parameter substitution; a parameterized macro is not, for the moment, allowed as an affix).

The expansion of a macro with some number of parameters is its macro text after substitution of parameters and expansion of all parameterless macros occurring in it. (A parameterized macro occurring in the macro text is not, for the moment, expanded.)

The expansion of a parameterless macro is its macro text. (Clearly, recursive macros cannot occur.)

The translation of an affix expression is:

- 1) If the handle is a flag or predicate, then

if $\neg \alpha$. then goto λ ;

- 2) If the handle is an action, then

α ;

The compiler compiler may delete some elements from this most general translation:

- 1) it deletes the assignations to the procedure identifier if the handle of the rule is an action, correspondingly changing the Boolean procedure into a procedure;
- 2) it deletes the <label>s to which no jumps occur, together with those parts of the procedure that can only be reached via those <label>s.

As an example, consider a part of a compiler description:

```
'action' set to zero.
'macro' 'action' put = '1' ['2']:= '3', incr = '1':= '1' + 1,
      make = '1':= '2'.
'macro' 'predicate' equal = '1' = '2'.
set to zero * list + min + max - p:
      make + p + min,
rep: put + list + p + 0,
      (equal + p + max;
      incr + p, :rep).
```

Translation:

```
procedure set to zero (list, min, max); integer max, min;
begin integer p;
      p:= min;
      rep: list [p]:= 0;
      begin if  $\neg$  (p = max) then goto 2;
          goto end;
          2: p:= p + 1; goto rep;
      end;
end;
```

A tag with type zero, when applied as an affix, becomes by context an applied global pointer. If no declaration or external specification for some applied global pointer occurs in the compiler description, then it will be treated at the end as a terminal (2.8). This is the only provision for terminal symbols: in a compiler description they occur as affixes without defining occurrence.

2.8 Starting symbol.

```
starting symbol:
      result symbol, handle, point.
```

The starting symbol is translated as a call for the procedure corres-

ponding to its handle. In front of this call, for each applied but not defined global pointer a <declaration> and a read <statement> is given. Between these <declaration>s and the read <statement>s a preparation procedure initialize for reading is called once. If there are no terminals, then also this call is obviated.

Example: assume one terminal, *t*, and starting symbol sentence:

```
integer t;
initialize for reading;
read (t);
sentence

end
```

Consequences:

- 1) in executing the obtained ALGOL 60 program, a representation for each terminal is to be provided as input.
- 2) If there are terminals, the following actions have to be defined in the compiler description:

read, out and *initialize for reading*.
- 3) In restoring parsers, the pointer *pin* must be defined. (Other representations may be taken if the relevant input to the compiler compiler is changed.)

2.9 Diagnostics.

For debugging a compiler description, various diagnostic aids are provided.

In the first place, the whole text is printed out while it is read, with two empty lines between building stones such as rules, declarations, etc. Input is assumed to come from cards, so lines have a maximum width of 80 positions. Output is assumed to proceed on cards, with a maximum of 72 positions used. A \$ in column 72 of an output card is used to indicate continuation on the next card. Continuation cards are not counted. The number of the input line appears on the left of its printout, the number of the current output card on its right.

If, at the end of a rule, some local has not been both defined and applied, then a warning message is printed.

If, at the end of the compiler description, some global has not been both defined and applied, then also a warning message appears. Various other errors are signalled by messages which should speak for themselves.

The compiler compiler makes use of a number of tables of fixed size. Whenever the upper bound of one of those tables is exceeded, but also at the end of successful execution, a "post mortem" is printed, giving the lower and upper bound of each table, with the space occupied at that moment, and finally the number of cards read and punched.

These tables, with lower and upper bounds are:

ttag	100001	110000	(tags)
tbold	200001	200300	(bolds)
tspec	300001	300100	(specs)
tcons	400001	400300	(constants)
lloc	500001	500200	(locals)
lglob	600001	601000	(globals)
lmacr	700001	701100	(macrottext)
ltext	800001	804000	(line numbers of occur- rences of globals).

The bounds are such that a compiler of the size and complexity of the compiler compiler itself needs only 50% of each of those tables.

Finally, at the end of the execution, a list is printed in alphabetical order of all the globals and macros occurring in the grammar, with their linenumbers of occurrence.

3. Examples.

In this chapter some examples are given and explained of texts in CDL. They are all assumed to occur in an environment containing:

```
'unrestore'
'macro' 'action'
get = '3':= '1' ['2'], put = '1' ['2']:= '3',
make = '1':= '2',
mark = '1':= -'1',
add = '3':= '1' + '2', subtr = '3':= '1' - '2',
addmult = '4':= '1' * '2' + '3',
divrem = '3':= '1' ÷ '2'; '4':= '1' - '2' * '3',
incr = '1':= '1' + 1, decr = '1':= '1' - 1.
'macro' 'predicate'
marked = '1' < 0,
less = '1' < '2', equal = '1' = '2', lseq = '1' ≤ '2'.
'global' 'action' resym, prsym, exit.
'global' 'pointer' newlinechar, spacechar.
```

This is a rather minimal environment with which a lot can be done. The macro actions *get* and *put* are for accessing array elements; *make* performs the assignation; *mark* inverts the sign of an integer variable, which we will use as marking bit; *add* and *subtr* allow arithmetic; *addmult* and *divrem* serve for packing and unpacking small integers into a word; and *incr* and *decr* serve to increment and decrement a variable by one.

The system procedure *resym* reads one character from input (without possibility of backtracking) and assigns to its only parameter an integer, viz. the internal value of that character, which does not exceed 255. We will assume that the digits have internal value 0 - 9 respectively, and the letters the internal values 10 - 35. Complementarily, *prsym* prints out the character corresponding to the value of its parameter. The global variables *newlinechar* and *spacechar* contain the internal values of the suggested characters. A call of *exit* terminates execution.

The following set of examples introduces one by one some building stones for a compiler for an ALGOL like language. Each example may make use of objects defined in a previous example.

3.1 Reading a number.

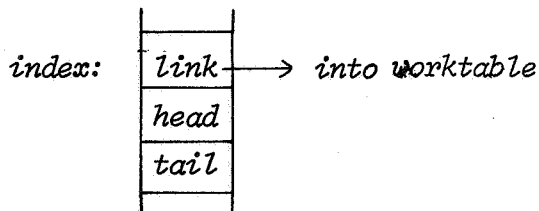
By now we have seen so many versions of number that it is time to see a good one:

```
'action' nextchar.
'pointer' char.
number + val - d:
    digit + d, make + val + d,
    rep: (digit + d, addmult + val + 10 + d + val, :rep;).
digit + d:
    lseq + char + 9, make + d + char, nextchar.
nextchar:
    resym + char,
    rep: (equal + char + spacechar, resym + char, :rep;).
```

Before the first number is read, *initialize for reading* has to take care of calling *nextchar* once.

3.2 Updating a chain.

In a worktable, we want to keep a number of chains of two elements of the form:



Each element points, via its link, to the next of the chain, the last element having link 0. Each chain is accessible by a variable which either points to its first element or contains 0 to indicate an empty chain. We want to be able to

- 1) search a chain for an element whose head has a specific value, and, if such an element exists, obtain its tail.
- 2) insert an element at the head of a chain.
- 3) see whether a specific element is already there, and in that case

obtain its index.

- 4) detach an element, if any, from the head of a chain, but leave it in the worktable (no garbage-collection).

An updating system with these properties is needed in a compiler, e.g., for collecting declarations: The most recently found declaration for a specific identifier is always the first to be looked at and the first to lose its significance, although the information obtained is needed again afterwards and may therefore not disappear from the worktable.

```
'macro' 'pointer' min work = 1, max work = 4095.
'list' work [min work : max work].
'pointer' pwork.
'action' insert 2, initialize chain adm, errormessage.
'macro' 'action'
get link = '2':= work ['1'], put link = work ['1']:= '2',
get head = '2':= work ['1' + 1], put head = work ['1' + 1]:= '2',
get tail = '2':= work ['1' + 2], put tail = work ['1' + 2]:= '2'.

search 2 + p + hd + tl - q - x:
    make + q + p,
    fnd: (lseq + min work + q, get head + q + x,
        (equal + x + hd, get tail + q + tl;
         get link + q + q, :fnd)).
insert 2 + p + hd + tl - q:
    make + q + p, make + p + pwork, add + pwork + 3 + pwork,
    (less + max work + pwork, errormessage + work full;
     put head + p + hd, put tail + p + tail, put link + p + q).
already there + p + hd + tl + q - x:
    make + q + p,
    fnd: (lseq + min work + q,
        (get head + q + x, equal + x + hd, get tail + q + x,
         equal + x + tl;
         get link + q + q, :fnd)).
```

```
detach 2 + p + hd + tl:
    lseq + min work + p, get head + p + hd, get tail + p + tail,
        get link + p + p.
'pointer' access.
initialize chain adm:
    make + pwork + min work, make + access + 0.
```

After one call of *initialize chain adm*, *access* gives access to an empty list. If the worktable becomes too small, *errormessage* is called (see 3.4).

If the compiler is to be run on a computer where the wordlength is such that the head and tail may fit together in one word, then one may simply replace the macro text of put head, get head, put tail and get tail by a version which manipulates with halfwords.

An advantage of using macros for accessing fields of a datastructure is that one can change the datastructures without invalidating the whole grammar. One could, e.g., in testing the compiler perform no packing, and in a later production version introduce packing. This allows the development of quite machine-independent compilers without a resulting price in efficiency.

As an initialization, *initialize chain adm* has to be called.

Note that work full is a terminal symbol (being an affix without defining occurrence) and that a representation of that symbol will have to be at the beginning of the input to the compiler resulting from this compiler description.

3.3 Reading tags.

We want to read <identifier>s, i.e. tags, and obtain a unique key for each different tag. We want to be able to list the tags in alphabetical order. Accompanying each tag we must remember a pointer, originally zero.

We choose the following storage organization: the information is kept in cells linked together in a text table, one cell for each tag, a cell consisting of a fixed and a variable part.

The fixed part consists of 3 pointer:

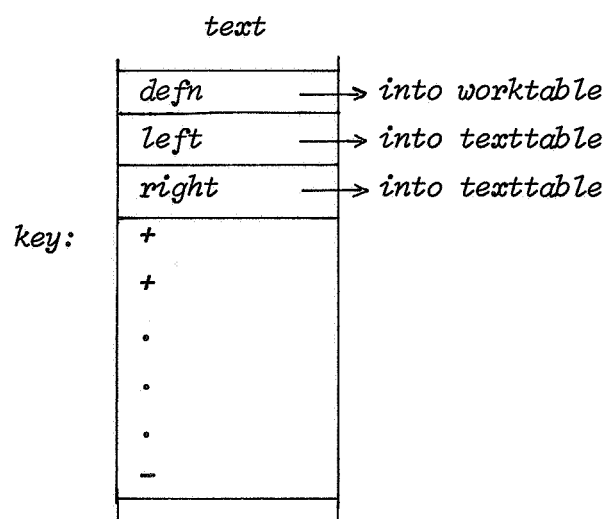
defn an access to a definition chain;
left a pointer to the chain of all tags alphabetically preceding the tag, zero if empty;
right idem for the alphabetically following tags.

The variable part contains the successive characters of the tag, packed three to a word, the last word being filled out with dummies (= 255) and provided with a - sign, the others with a + sign. (We will assume that the wordlength of the computer on which the compiler will run is sufficient, i.e. at least 25 bits.)

We will give out the index (in the text table) of the first word of the variable part as a key.

The chainstructure just described is a doubly linked namelist. It is more efficient to use than a linear namelist, and less efficient than a namelist using hash coding, over which again it has the advantage of keeping the tags in alphabetical order.

In a picture:



When reading a tag, we will first build up a cell for it in the text table, and then enter it, checking whether the same text already occurs in the table; if so, then the key of the old occurrence is delivered and the new cell removed; if not, then the new cell is fitted into the chain.

```

'macro' 'pointer' mintext = 1, max text = 4095.
'list text [min text : max text].
'pointer' ptext.
'macro' 'action'
put defn = text ['1' - 3]:= '2', get defn = '2':= text ['1' - 3],
put left = text ['1' - 2]:= '2', get left = '2':= text ['1' - 2],
put right= text ['1' - 1]:= '2', get right= '2':= text ['1' - 1].
'macro' 'predicate'
    is letter = '1' > 9 ^ '1' < 36,
    is letgit = '1' < 36.

[treatment of text]
'action' add to text, stack text, stack last, enter, reserve fixed part.
'pointer' word, c.
'macro' 'pointer' dummy = 255, bytesize = 256, maxbytes = 3.
'macro' 'predicate' left = abs('1') < abs('2').
add to text + x:
    equal + c + maxbytes, stack text + word, make + word + x,
                                                make + c + 1;
    addmult + word + bytesize + x + word, incr + c.
stack text + x:
    lseq + ptext + max text, put + text + ptext + x, incr + ptext;
    errormessage + text full.
stack last:
    rep: (less + c + maxbytes, addmult + word + bytesize + dummy + word,
        incr + c, :rep;
    mark + word, stack text + word).

```

Now follows the *pièce de résistance* of the treatment of text. The action enter checks whether the topmost element of the texttable, which is pointed to by x, already occurred somewhere else in the texttable; if such is the case, then x is made to point to that older element, and the new element is deleted. Also, this action takes care of the alphabetic ordering of the tags.

```

enter + x + pchain - y - x1 - y1 - wx - wy:
    equal + pchain + 0, make + pchain + x, reserve fixed part;
    make + y + pchain,
nxy: make + x1 + x, make + y1 + y,
nxw: get + text + x1 + wx, get + text + y1 + wy,
    (equal + wx + wy,
      (marked + wx,
        (equal + x + y, reserve fixed part;
          make + ptext + x, make + x + y);
        incr + x1, incr + y1, :nxw);
      left + wx + wy, get left + y + wy,
        (equal + wy + 0, put left + y + x, reserve fixed part;
          make + y + wy, :nxy);
      get right + y + wy,
        (equal + wy + 0, put right + y + x, reserve fixed part;
          make + y + wy, :nxy)).
reserve fixed part:
    stack text + 0, stack text + 0, stack text + 0.

[reading a tag]
'pointer' ptag.
read tag + x - t:
    letter + t, make + word + t, make + c + 1, make + x + ptext,
    nxt: (letgit + t, add to text + t, :nxt;
      stack last, enter + x + ptag).
letter + t:
    is letter + char, make + t + char, nextchar.
letgit + t:
    is letgit + char, make + t + char, nextchar.
initialize text adm:
    make + ptext + min text, reserve fixed part, make + ptag + 0.

```

Of course the left- and right-pointer of a tag can be packed together again in one word. As initialization, *initialize text adm* has to be called. The pointer ptag gives access to the chain of tags, text full is a terminal.

3.4 Printing, errormessages.

We are now in a position to define the printing of tags, and the treatment of errormessages.

```
'action' print, print 1, print 2, nlcr, errormessage.
print + x - p - el:
    make + p + x,
        rep: get + text + p + el,
            (marked + el, mark + el, print 1 + el;
            print 1 + el, incr + p, :rep).
print 1 + t - c:
    prsym + spacechar,
    make + c + max bytes, print 2 + c + t.
print 2 + c + t - el:
    equal + c + 1, prsym + el;
    divrem + t + bytesize + t + el, decr + c, print 2 + c + t,
                                                prsym + el.

nlcr: prsym + newlinechar.
errormessage + x: nlcr, print + x, exit.
```

The definition of print 2 has to be recursive in order to get the characters out in the right order.

The procedure *errormessage* prints a message and then terminates execution. A more subtle reaction might be envisaged.

3.5 An input administration.

Assuming that there also exist texts, similar to that in 3.3, for reading bold symbols, special symbols and constants, a complete input administration can be constructed:

```

'macro' 'pointer' min inpt = 1, max inpt = 2047.
'list' inpt [min inpt : max inpt].
'pointer' pin, xin, symb.
'action' next symbol, skip layout.
next symbol:
  nxt: skip layout,
    (read tag + symb;
     read bold + symb, add + symb + 20000 + symb;
     read spec + symb, add + symb + 30000 + symb;
     read const+ symb, add + symb + 40000 + symb;
     errormessage + incorrect char, nextchar, :nxt).
skip layout:
  skp: (equal + char + new line char, nextchar, :skp;).
rq + x:
  equal + symb + x, incr + pin,
    (equal + pin + xin, next symbol, put + inpt + pin + symb,
     incr + xin;
     get + inpt + pin + symb).
identifier + x:
  less + symb + 20000, make + x + symb, rq + x.
initialise for reading:
  nextchar, initialize chain adm, initialize text adm,
  make + xin + min inpt, make + pin + xin, nextsymbol, put + inpt +
  pin + symb.

```

The pointer *pin* serves as reading pointer, *xin* indicates up to what index the inputarray has been filled with symbol-keys. In *next symbol*, a multiple of 10000 is added to some keys as a type-indication, allowing discrimination between, e.g., tags and bolds. By means of *rq*, one can ask whether the current symbol (kept in *symb*), is equal to a given symbol, while *identifier* returns true only if the current symbol is a tag. A restoring parser can now work on the basis of this input administration.

3.6 Reading declarers.

We will now show how to recognize declarers not unlike those of ALGOL 68. During recognition of a declarer, some (head, tail) pairs may be added to the worktable (3.2), and the mode of the declarer is represented by the index of one such pair.

Let μ stand for a declarer, $\bar{\mu}$ for its index, τ for a tag, $\bar{\tau}$ for its key.

int gives an index to (0, int)
real gives (0, real)
long μ gives (long, $\bar{\mu}$)
ref μ gives (ref, $\bar{\mu}$) and
struct ($\mu_1\tau_1, \mu_2\tau_2, \dots, \mu_n\tau_n$) gives
 (struct,) (,) (,) ... (, 0)
 ↖ ↖ ↖
 ($\bar{\mu}_1, \bar{\tau}_1$) ($\bar{\mu}_2, \bar{\tau}_2$) ... ($\bar{\mu}_n, \bar{\tau}_n$)

In adding a pair to the worktable, we take care that no two copies of a pair are ever inserted: in adding a pair that was already there, the index of the old copy is obtained. This implies that equivalent modes automatically get equal indices.

'action' add to decl.

'pointer' pdecl.

declarer + mode:

primitive declarator + mode;

long declarator + mode;

ref declarator + mode;

struct declarator + mode.

'restore'

primitive declarator + mode:

rq + int, add to decl + 0 + int + mode;

rq + real, add to decl + 0 + real + mode.


```

long declarator + mode:
    rq + long, declarer + mode, add to decl + long + mode + mode.
ref declarator + mode:
    rq + ref, declarer + mode, add to decl + ref + mode + mode.
struct declarator + mode:
    rq + struct, rq + open, fields + mode, add to decl + struct +
                                                mode + mode.
fields + mode - m1 - m2:
    rq + close, make + mode + 0;
    field + m1, fields + m2, add to decl + m1 + m2 + mode.
field + mode - tag:
    declarer + mode, identifier + tag, add to decl + mode + tag + mode.

'unrestore'
add to decl + hd + tl + mode - oldp:
    already there + hd + tl + pdecl + oldp, make + mode + oldp;
    make + oldp + pwork, insert 2 + hd + tl + pdecl, make + mode +
                                                oldp.

```

Note that for the declarators we want restoring parsers. Also note that affixes that have no defining occurrence, e.g., *int*, *real*, *long*, etc., are terminals, so that a representation for them has to be given as input to the resulting compiler.

For *already there* and *insert 2*, see 3.2.

3.7 Collecting defining occurrences.

We will now show how to collect defining occurrences of identifiers. We assume the existence of a global pointer *blocknumber* that is automatically updated by some other part of the compiler description.

A defining occurrence of an identifier is an occurrence in a declaration. For each defining occurrence we want to store its *blocknumber* and *mode* in the *worktable*. We will not allow an identifier to be defined twice in one block.

```
'global' 'pointer' blocknumber.  
'action' define identifier  
declaration - mode - idf:  
  declarer + mode, identifier + idf,  
  rep: (define identifier + idf + mode,  
        (rq + comma,  
          (identifier + idf, :rep;  
            declarer + mode, identifier + idf, :rep;  
            errormessage + incorrect declaration);  
          rq + semicolon)).  
define identifier + idf + mode - def - dummy:  
  get defn + idf + def,  
  (search 2 + def + blocknumber + dummy, errormessage +  
    defined twice;  
  insert 2 + def + blocknumber + mode, put defn + idf + def).
```

A more elaborate version of this last example can be found in [2].

COMPILER COMPILER TOWARDS ALGOL60

```

1
2
3
4 [4 COMPILER COMPILER DESCRIBED IN CDL *****] 2
5
6 'SHORT' 2
7
8 [4.1 GENERAL ENVIRONMENT] 2
9
10 [4.1.1 INTERFACE WITH MACHINE] 3
11 'EXTERNAL' 'ACTION' 3
12 NEW PAGE,EXIT,PRSYM,CSYM. 4
13
14 'EXTERNAL' 'POINTER' 'RESYM,' 4
15
16 'MACRO' 'FLAG' 4
17 WAS LETTER='1'>9^'1'<36, 4
18 WAS LETGIT='1'<36, 4
19 WAS DIGIT='1'<10, 4
20 WAS SPECCH='1'>63. 4
21
22 'MACRO' 'POINTER' 4
23 NIX=63,MINUSCODE=65,SPACECODE=93,TABCODE=118,NLCRCODE=119,ACCC=120,QUOTE=121. 4
24
25 [4.1.2 STACKS] 4
26 'MACRO' 'POINTER' 5
27 MIN TAG =100001,MAX TAG =110000, 5
28 MIN BOLD=200001,MAX BOLD=200600, 5
29 MIN SPEC=300001,MAX SPEC=300100, 5
30 MIN CONS=400001,MAX CONS=400300, 5
31 MIN LOC =500001,MAX LOC =500200, 5
32 MIN GLOB=600001,MAX GLOB=601000, 5
33 MIN MACR=700001,MAX MACR=701100, 5
34 MIN TEXT=800001,MAX TEXT=804000. 5
35
36 'LIST' 5
37 TTAG [MIN TAG :MAX TAG ], 6
38 TROLL [MIN BOLD:MAX BOLD], 6
39 TSPEC [MIN SPEC:MAX SPEC], 6
40 TCONS [MIN CONS:MAX CONS], 6
41 LLOC [MIN LOC :MAX LOC ], 6
42 LGLOB [MIN GLOB:MAX GLOB], 6
43 LTEXT [MIN TEXT:MAX TEXT], 6
44 LMACR [MIN MACR:MAX MACR]. 6
45
46 'MACRO' 'FLAG' 6
47 WAS TAG =MIN TAG "LE" '1' ^ '1' <PTAG , 6
48 WAS BOLD=MIN BOLD "LE" '1' ^ '1' <PBOLD, 6
49 WAS SPEC=MIN SPEC "LE" '1' ^ '1' <PSPEC, 6
50 WAS CCNS=MIN CONS "LE" '1' ^ '1' <PCONS, 6

```

```

46
47 [4.1.3  MACROS]
48 'MACRO' 'ACTION'
49 GET  = '3':='1'['2'],
50 PUT  = '1'['2']:= '3',
51 MAKE= '1':='2',
52 SET  = '1':='2',
53 MARK= '1':=- '1',
54 ADD  = '3':='1'+ '2',
55 SUBTR= '3':='1'- '2',
56 ARDMULT= '4':='1'* '2'+ '3',
57 D VREM= '3':='1'/' '2'; '4':='1'- '2'* '3',
58 PACK3= '4':=( '1'*128+ '2')*128+ '3',
59 UNPACK3= '2':='1'/' '16384; '4':='1'-16384* '2'; '3':='4'/' '128; '4':='4'-128* '3',
60 PACK2= '3':='192*' '1'+ '2',
61 UNPACK2= '2':='1'/' '8192; '3':='1'-8192* '2',
62 GET TAIL= '2':='1'- '1'/' '8192*8192,
63 INCR  = '1':='1'+1,
64 DECR  = '1':='1'-1.

65 'MACRO' 'FLAG'
66 MARKED= '1'<0,
67 LEFT=ABS('1')<ABS('2'),
68 LESS= '1'<'2',
69 LSEQ= '1'≤'2',
70 EQUAL= '1'='2'.

71 'MACRO' 'POINTER'
72 FALSE= "FALSE",
73 TRUE= "TRUE".

74 [DATA STRUCTURES AND THEIR ACCESS]
75 'MACRO' 'ACTION'
76 GET NPARS= '2':=TTAG['1'-8],PUT NPARS=TTAG['1'-8]:='2',
77 GET MTEXT= '2':=TTAG['1'-7],PUT MTEXT=TTAG['1'-7]:='2',
78 GET PLACE= '2':=TTAG['1'-6],PUT PLACE=TTAG['1'-6]:='2',
79 GET STATE= '2':=TTAG['1'-5],PUT STATE=TTAG['1'-5]:='2',
80 GET SORT  = '2':=TTAG['1'-4],PUT SORT  =TTAG['1'-4]:='2',
81 GET TYPE  = '2':=TTAG['1'-3],PUT TYPE  =TTAG['1'-3]:='2',
82 GET LEFT  = '2':=TTAG['1'-2],PUT LEFT  =TTAG['1'-2]:='2',
83 GET RIGHT = '2':=TTAG['1'-1],PUT RIGHT =TTAG['1'-1]:='2'.

84
85 [4.1.4  POINTERS AND FLAGS]
86 'POINTER'
87 PTAG,PBOLD,PSPEC,PCONS,
88 P.LOC,P.GLOB,P.TEXT,P.MACR,X.LOC.

89 'FLAG' GIVE TEXT,GIVE TRACE,LEGIBLE.

90 'POINTER' FIRST TAG.

91

```

92		12
93		12
94		12
95		12
96	[4.2 OUTPUT]	13
97		13
98	[4.2.1 PRINTER SECTION]	14
99	'ACTION'OUT,OUTINT,PRINT,PRINT1,NLCR,TAB,SPACE,OUTINT1,PRCHAR,SPACES,TABS,	14
100	POSITION,SHIFT2LINES.	
		14
101	'POINTER'POS.	
		15
102	OUT+X-EL:	16
103	WAS TAG +X,PRINT+TTAG +X;	16
104	WAS BOLD+X,PRINT+TBOLD+X;	16
105	WAS SPEC+X,PRINT+TSPEC+X;	16
106	WAS CONS + X, GET+TCONS+X+EL,OUTINT+EL;	16
107	OUTINT+X.	
		16
108	OUTINT+X-QUOT-REM:	17
109	SPACE,	17
110	LESS+X+0,MAKE+REM+X,MARK+REM,PRCHAR+MINUSCODE,OUTINT+REM;	17
111	EQUAL+X+0,PRCHAR+0,SPACE;	17
112	DIVREM+X+10+QUOT+REM,OUTINT1+QUOT,PRCHAR+REM,SPACE.	
		17
113	OUTINT1+X-QUOT-REM:	18
114	EQUAL+X+0;	18
115	DIVREM+X+10+QUOT+REM,OUTINT1+QUOT,PRCHAR+REM.	
		18
116	PRINT*LIST+Y-X-LX:	19
117	MAKE+X+Y,SPACE,	19
118	RST: GET+LIST+X+LX,	19
119	(MARKED+LX,MARK+LX,PRINT1+LX;PRINT1+LX,INCR+X,:RST).	
		19
120	PRINT1+X-X1-X2-X3:	20
121	UNPACK3+X+X1+X2+X3,	20
122	PRCHAR+X1,PRCHAR+X2,PRCHAR+X3,	
		20
123	POSITION+X-TBS-SPCES:	21
124	SUBTR+X+POS+SPCES,LSSEQ+0+SPCES,SPACES+SPCES;	21
125	NLCR,DIVREM+X+8+TBS+SPCES,TABS+TBS,	21
126	(EQUAL+SPCES+0;SPACES+SPCES).	
		21
127	SPACES+X-N:	22
128	MAKE+N+0,	22
129	SPC: LESS+N+X,SPACE,INCR+N,:SPC;	
		22
130	TABS+X-N:	23
131	MAKE+N+0,	23
132	TBS: LESS+N+X,TAB,INCR+N,:TBS;	

133	NLCR:		23
134		PRSYM +NLCRCODE,MAKE+POS+0.	24
135	SHIFT 2 LINES-OLD POS:		24
136		MAKE+OLDPOS+POS,NLCR,NLCR,SPACES+OLD POS.	25
137	TAB:PRSYM+TABCODE,ADD+POS+8+POS.		25
138	SPACE:		26
139		PRCHAR+SPACECODE.	27
140	PRCHAR+X:		27
141		EQUAL+X+NLCRCODE,NLCR;	28
142		EQUAL+X+TABCODE,TAB;	28
143		EQUAL+X+NIX;	28
144		PRSYM+X,INCR+POS.	28
145			28
146	[4.2.2 TRACE ADMINISTRATION]		28
147	'ACTION'TRACE,SIGNAL,INFORM,ERROR.		29
148	TRACE+X:		29
149		GIVE TRACE,POSITION+32,INFORM+X;	30
150	SIGNAL+X:		30
151		POSITION+16,INFORM+X.	31
152	INFORM+X-STATE-SORT-TYPE:		31
153		GET STATE+X+STATE,GET SORT+X+SORT,GET TYPE+X+TYPE,	32
154		OUT+STATE,OUT+SORT,OUT+TYPE,OUT+X.	32
155	ERROR+TEXT+INFO-OLD POS:		32
156		MAKE+OLD POS+POS,NLCR,OUT+TEXT,	33
157		(WAS TAG+INFO,INFORM+INFO,POSITION+OLD POS;	33
158		EQUAL+INFO+0,POSITION+OLD POS;	33
159		OUT+INFO,POSITION+OLD POS).	33
160			33
161	[4.2.3 CARD PUNCH SECTION]		33
162	'ACTION'WRITE,WRITEINT,WRITEINT1,PUNCH,PUNCH1,NEW CARD,TAB CARD,SPACE CARD,		34
163	PUCHAR.		34
164	'MACRO'POINTER!		34
165	MAX CARD=72,DOLLAR CODE=133.		34
166	'POINTER' CARD,CPOS.		34

167	WRITE +X - EL:	35
168	WAS TAG + X, PUNCH + TTAG + X;	36
169	WAS BOLD + X, PUNCH + TBOLD + X;	36
170	WAS SPEC + X, PUNCH + TSPEC + X;	36
171	WAS CONS + X, GET+TCONS+X+EL, WRITE INT+EL;	36
172	WRITE INT+X.	
173	WRITE INT+X-QUOT-REM:	36
174	LESS+X+0, MAKE+REM+X, MARK+REM, PUCHAR+MINUSCODE, WRITE INT+REM;	37
175	EQUAL+X+0, PUCHAR+0;	37
176	DIVREM+X+10+QUOT+REM, WRITE INT1+QUOT, PUCHAR+REM.	37
177	WRITE INT1+X-QUOT-REM:	38
178	EQUAL+X+0;	38
179	DIVREM+X+10+QUOT+REM, WRITE INT1+QUOT, PUCHAR+REM.	38
180	PUNCH+LIST+Y-X-LX:	38
181	MAKE+X+Y,	39
182	RST: GET+LIST+X+LX,	39
183	(MARKED+LX, MARK+LX, PUNCH1+LX;	39
184	PUNCH1+LX, INCR+X, :RST),	39
185	PUNCH1+X-X1-X2-X3:	39
186	UNPACK3+X+X1+X2+X3,	40
187	PUCHAR+X1, PUCHAR+X2, PUCHAR+X3.	40
188	NEW CARD:	40
189	SPC: LSEQ +CPOS + MAX CARD, CSYM+SPACECODE, INCR+CPOS, :SPC;	41
190	MAKE+CPOS+1, WRITE INT+CARD, CSYM+NLCR CODE, INCR+CARD, MAKE+CPOS+1.	41
191	TAB CARD-SPCES-TBS:	41
192	DIVREM+CPOS+8+TBS+SPCES,	42
193	SPC: LESS+SPCES +8, SPACE CARD, INCR+SPCES ,:SPC;	42
194	SPACE CARD:	42
195	PUCHAR+SPACECODE.	43
196	PUCHAR+X:	43
197	EQUAL + X + NLCRCODE, NEW CARD;	44
198	EQUAL + X + TABCODE, TAB CARD;	44
199	EQUAL + X + NIX;	44
200	LESS+CPOS+MAX CARD.CSYM+X, INCR+CPOS;	44
201	CSYM+DOLLARCODE, CSYM+NLCRCODE, MAKE+CPOS+1, PUCHAR+X.	44
202		44
203	[4.2.4 RESULT ADMINISTRATION]	44
204	'ACTION'G, LINE UP, PUTABS, U, L, NEW LINE, BLANK LINE, TAB LINE,	45
205	'FLAG' LINED UP.	45

206	'POINTER' INDENTATION.	46
207	G + X:	47
208	LINE UP, WRITE + X.	48
209	LINE UP:	48
210	LINED UP;	49
211	LEGIBLE,PUTABS+INDENTATION,SET+LINED UP+TRUE.	49
212	PUTABS + N - N1:	49
213	LSEQ + N + 0;	50
214	MAKE + N1 + N, DECR + N1, TAB CARD, PUTABS+ N1.	50
215	U: INCR + INDENTATION.	50
216	L: DECR + INDENTATION.	51
217	NEW LINE:	52
218	LEGIBLE,LINED UP,BLANK LINE.	53
219	BLANK LINE:	53
220	NEW CARD, SET + LINED UP + FALSE.	54
221	TAB LINE: LEGIBLE,TAB CARD.	54
222		55
223		55
224		55
225		55
226		55
227		55
228	[4.3 INPUT]	56
229	'ACTION' POST MORTEM.	56
230	'POINTER'	56
231	LINE,CHAR,INPT.	56
232		57
233	[4.3.1 READING CHARACTERS]	57
234	'ACTION'NEXTCHAR,DISPLAY CHARACTER,NEXT NON SPACE CHAR.	58
235	LETTER+X:	58
236	WAS LETTER+CHAR,MAKE+X+CHAR,NEXT NON SPACE CHAR.	59
237	LEFGIT+X:	59
		60

238 WAS LETGIT+CHAR,MAKE+X+CHAR,NEXT NON SPACE CHAR.

239 DIGIT+X:

240 WAS DIGIT +CHAR,MAKE+X+CHAR,NEXT NON SPACE CHAR.

241 BOLDCHAR+X:

242 EQUAL+CHAR+QUOTE,MAKE+X+ACCC,NEXTCHAR;

243 MAKE+X+CHAR,NEXTCHAR.

244 ACCENT:EQUAL+CHAR+ACCC,NEXTCHAR.

245 SPECCHAR + X:

246 WAS SPECCH+CHAR,MAKE+X+CHAR,NEXT NON SPACE CHAR.

247 NEXTCHAR:

248 DISPLAY CHARACTER,MAKE + CHAR + RESYM.

249 DISPLAY CHARACTER:

250 EQUAL+CHAR+NLCRCODE,INCR+LINE,

251 (GIVE TEXT,POSITION+133,OUT'NT1+CARD,

252 POSITION+48,OUT'NT1+LINE,POSITION+52;);

253 GIVE TEXT,PRCHAR+CHAR)

254 NEXT NON SPACE CHAR:

255 CHR: NEXTCHAR,

256 (EQUAL+CHAR+SPACECODE,:CHR;).

257 LAYOUT SYMBOL;

258 EQUAL + CHAR + SPACECODE;

259 EQUAL + CHAR + NLCRCODE;

260 EQUAL + CHAR + TABCODE .

261

262 [4,3,2 TAG SYMBOLS]

263 'ACTION'STACK TAG,ENTER TAG,RESERVE ADMIN SPACE.

264 READ TAG + X = T - T1 - T2 - T3 - T4:

265 LETTER+T1,MAKE+X+PTAG,

266 NXT: (LETGIT + T2,

267 (LETGIT + T3,

268 (LETGIT + T4, PACK3 + T1 + T2 + T3 + T,

269 STACK TAG+T,MAKE+T1+T4,:NXT;

270 PACK3 + T1 + T2 + T3 + T, :LST);

271 PACK3 + T1 + T2 + NIX + T, :LST);

272 PACK3 + T1 + NIX + NIX + T,

273 LST: MARK + T, STACK TAG + T).

274 STACK TAG + X:

275 LSEQ + PTAG + MAX TAG, PUT + TTAG + PTAG + X, INCR + PTAG;

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```
276 ERROR + TAG FULL+0,POST MORTEM,EXIT. 71
277 ENTER TAG + X - Y - X1 - Y1 - WX - WY: 72
278 MAKE + Y + FIRST TAG, 72
279 NXY: MAKE + X1 + X, MAKE + Y1 + Y, 72
280 NXW: GET+TTAG+X1+WX,GET+TTAG+Y1+WY, 72
281 (EQUAL + WX + WY, 72
282 (MARKED + WX, 72
283 (EQUAL + Y + X,RESERVE ADMIN SPACE; 72
284 MAKE+PTAG+X,MAKE+X+Y); 72
285 INCR+X1,INCR+Y1,:NXW); 72
286 LEFT+WX+WY,GET LEFT+Y+WY, 72
287 (EQUAL+WY+0,PUT LEFT+Y+X,RESERVE ADMIN SPACE; 72
288 MAKE+Y+WY,:NXY); 72
289 GET RIGHT+Y+WY, 72
290 (EQUAL+WY+0,PUT RIGHT+Y+X,RESERVE ADMIN SPACE; 72
291 MAKE+Y+WY,:NXY)). 72
292 RESERVE ADMIN SPACE: 72
293 ADD+PTAG+8+PTAG,LESS+PTAG+MAX TAG,PUT NPARS+PTAG+999, 73
294 PUT MTEXT+PTAG+0,PUT PLACE+PTAG+0,PUT STATE+PTAG+0,PUT SORT+PTAG+0, 73
295 PUT TYPE+PTAG+0,PUT LEFT+PTAG+0,PUT RIGHT+PTAG+0; 73
296 ERROR+TAG FULL+0,POST MORTEM,EXIT. 73
297 73
298 [4,3,3 BOLD SYMBOLS] 73
299 'ACTION'STACK BOLD,ENTER BOLD. 74
300 READ BOLD + X - B1-B2-B3-D: 74
301 ACCENT, MAKE+X+PBOLD, 75
302 (ACCENT,PACK3+NIX+NIX+NIX+D,MARK+D,STACK BOLD+D; 75
303 BOLDCHAR+B1, 75
304 RST: (ACCENT,PACK3+B1+NIX+NIX+D,MARK+D,STACK BOLD+D; 75
305 BOLDCHAR+B2, 75
306 (ACCENT,PACK3+B1+B2+NIX+D,MARK+D,STACK BOLD+D; 75
307 BOLDCHAR+B3,PACK3+B1+B2+B3+D, 75
308 (ACCENT,MARK+D,STACK BOLD+D; 75
309 BOLDCHAR+B1,STACK BOLD+D,:RST))). 75
310 STACK BOLD + X: 75
311 LSER + PBOLD + MAX BOLD, PUT + TBOLD + PBOLD + X, INCR+PBOLD; 76
312 ERROR + BOLD FULL+0,POST MORTEM,EXIT. 76
313 ENTER BOLD + X - Y - X1 - Y1 - WX - WY: 76
314 MAKE+Y+MIN BOLD, 77
315 NXY: MAKE + X1 + X, MAKE + Y1 + Y, 77
316 NXW: GET+TBOLD+X1+WX,GET+TBOLD+Y1+WY, 77
317 (EQUAL + WX + WY, 77
318 (MARKED + WX, 77
319 (EQUAL+Y+X; 77
320 MAKE+PBOLD+X,MAKE+X+Y); 77
321 INCR+X1,INCR+Y1,:NXW); 77
322 SKP: MARKED + WY, MAKE + Y + Y1, INCR + Y, :NXY; 77
323 INCR + Y1, GET + TBOLD + Y1 + WY, :SKP), 77
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324		77
325	[4.3.4 SPECIAL SYMBOLS]	77
326	'ACTION'STACK SPEC,ENTER SPEC.	78
327	READ SPEC + X - S - S1:	78
328	SPECCHAR+S1,MAKE+X+PSPEC,PACK3+S1+NIX+NIX+S,MARK+S,STACK SPEC+S.	79
329	STACK SPEC + X:	79
330	LSEQ + PSPEC + MAX SPEC, PUT + TSPEC + PSPEC + X, INCR+PSPEC;	80
331	ERROR + SPEC FULL+0,POST MORTEM,EXIT.	80
332	ENTER SPEC + X - Y - WX - WY:	80
333	MAKE+Y+MIN SPEC,GET+TSPEC+X+WX,	81
334	NXY: GET+TSPEC+Y+WY,	81
335	(EQUAL + WX + WY,	81
336	(EQUAL+Y+X;	81
337	MAKE+PSPEC+X,MAKE+X+Y);	81
338	INCR + Y, :NXY).	81
339		81
340	[4.3.5 CONSTANTS]	81
341	'ACTION'STACK CONS,ENTER CONS.	82
342	READ CONS + X - D - S:	82
343	DIGIT + D, MAKE + S + D,	83
344	RST: (DIGIT + D, ADDMULT + S + 10 + D + S,:RST;	83
345	MAKE + X + PCONS, STACK CONS + S).	83
346	S*ACK CONS + X:	83
347	LSEQ + PCONS + MAX CONS,PUT + TCONS + PCONS+ X, INCR + PCONS;	84
348	ERROR + CONS FULL+0,POST MORTEM,EXIT.	84
349	ENTER CONS + X - Y - S - T:	84
350	GET + TCONS + X + S, MAKE + Y + MIN CONS,	85
351	NXT: GET + TCONS + Y + T,	85
352	(EQUAL + S + T,	85
353	(EQUAL + X + Y;	85
354	MAKE + PCONS + X, MAKE + X + Y);	85
355	INCR + Y, :NXT).	85
356		85
357	[4.3.6 INPUT ADMINISTRATION]	85
358	'ACTION'INITIALIZE FOR READING,NEXT SYMBOL,READ.	86
359	INITIALIZE FOR READING:	86
360	MAKE+GIVE TEXT+FALSE,MAKE+PCONS+MIN CONS,	87
361	MAKE+PTAG+MIN TAG,MAKE+PBOLD+MIN BOLD,MAKE +PSPEC+MIN SPEC,MAKE+LINE+0,	87
362	MAKE+CHAR+RSYM,RESERVE ADMIN SPACE,MAKE+FIRST TAG+PTAG.	87

363 IS TAG + X: WAS TAG + INPT, MAKE + X + INPT, NEXT SYMBOL.

87

364 IS BOLD+X: WAS BOLD+INPT, MAKE+X+INPT, NEXT SYMBOL.

88

365 IS CONS + X: WAS CONS + INPT, MAKE + X + INPT, NEXT SYMBOL.

89

366 R + X: AHEAD + X, NEXT SYMBOL.

90

367 AHEAD + X: EQUAL + INPT + X.

91

368 NEXT SYMBOL:

369 READ+INPT.

92

93

370 READ+X:

371 SKP: LAYOUT SYMBOL, NEXTCHAR, :SKP;

372 READ TAG +X, ENTER TAG+X;

373 READ BOLD+X, ENTER BOLD+X;

374 READ SPEC+X, ENTER SPEC+X;

375 READ CONS+X, ENTER CONS+X;

376 ERROR+WRONG INPUT CODE+CHAR, NEXTCHAR, READ+X.

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382 [4.4 COMPILER COMPILER TO ALGOL60]

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383 'FLAG' GIVE RSTO, ADVANCED.

96

384 'POINTER' HANDLE, BOUNDX, SPECX, TO NEXT, TO LOST, TO END, CUR LAB, NEW LAB.

97

97

385

98

386 [4.4.1 AUXILIARY ACTIONS]

98

387 'ACTION'

98

388 FORGET LOCALS, WARN, ENTER LOCAL, ENTER GLOBAL, ENTER MACRO, ENTER TEXT,

98

389 ADD PLACE, REPLACE.

98

99

390 ADD PLACE+TAG-H=K-Q:

391 GET PLACE+TAG+H,

99

392 (EQUAL + H + 0, REPLACE + TAG + H;

99

393 GET+LTEXT+H+Q, GET TAIL+Q+K, LESS+K+LINE, REPLACE+TAG+H;).

99

100

394 REPLACE + TAG + H:

395 PUT PLACE+TAG+PTEXT,

100

396 (EQUAL+H+0, ENTER TEXT+LINE;

100

397 SUBTR+H+MIN TEXT+H, INCR+H, PACK2+H+LINE+H, ENTER TEXT+H).

100

398	FORGET LOCALS=LOC:	101
399	RST: EQUAL+PLOC+MIN LOC:	101
400	DECR+PLOC,GET+LLOC+PLOC+LOC,WARN+LOC,	101
401	PUT STATE+LOC+0,PUT SORT+LOC+0,PUT TYPE+LOC+0,IRST.	
402	WARN+X=STATE:	101
403	GET STATE+X+STATE,	102
404	(EQUAL+STATE+BLANK,TRACE+X;	102
405	EQUAL+STATE+0,TRACE+X;	102
406	SIGNAL+X).	
407	ENTER LOCAL + HEAD:	102
408	LSEQ + PLOC + MAX LOC, PUT + LLOC + PLOC + HEAD, INCR + PLOC,	103
409	(LSEQ + PLOC + XLOC; MAKE + XLOC + PLOC);	103
410	ERROR + LOC FULL+0,POST MORTEM,EXIT.	
411	ENTER GLOBAL + HEAD:	103
412	LSEQ + PGLOB + MAX GLOB, PUT + LGLOB + PGLOB + HEAD, INCR + PGLOB;	104
413	ERROR + GLOB FULL+0,POST MORTEM,EXIT.	104
414	ENTER MACRO + HEAD:	104
415	LSEQ + PMACR + MAX MACR, PUT + LMACR + PMACR + HEAD, INCR + PMACR;	105
416	ERROR + MACR FULL+0,POST MORTEM,EXIT.	105
417	ENTER TEXT + X:	105
418	LSEQ + PTEXT + MAX TEXT, PUT + LTEXT + PTEXT + X, INCR + PTEXT;	106
419	ERROR + TEXT FULL+0,POST MORTEM,EXIT.	106
420		106
421	[4.4.2 TREATMENT OF TYPES]	106
422	'ACTION'	107
423	DEFINE ACTION,DEFINE LABEL,DEFINE PREDICATE,DEFINE AFFIX,APPLY ACTION,	107
424	APPLY LABEL,REDEFINE,NEW PLACE,NEW NPARS,APPLY,APPLY PREDICATE.	107
425	DEFINE ACTION+HEAD:	107
426	REDEFINE+HEAD+DEFINED+GLOBAL+ACTION.	108
427	DEFINE PREDICATE+HEAD:	108
428	REDEFINE+HEAD+DEFINED+GLOBAL+PREDICATE.	109
429	DEFINE AFFIX+HEAD:	109
430	REDEFINE+HEAD+DEFINED+LOCAL+POINTER.	110
431	DEFINE LABEL+LAB:	110
432	WAS TAG +LAB,REDEFINE+LAB+DEFINED+LOCAL+LABEL;	111
433	WAS CONS+LAB.	111
434	APPLY ACTION+HEAD:	111
435	REDEFINE+HEAD+APPLIED+GLOBAL+ACTION.	112

	436 APPLY PREDICATE+HEAD:	112
	437 REDEFINE+HEAD+APPLIED+GLOBAL+PREDICATE.	113
	438 APPLY LABEL+LAB:	113
	439 WAS TAG +LAB,REDEFINE+LAB+APPLIED+LOCAL+LABEL;	114
	440 WAS CONS+LAB.	114
	441 APPLY+X:	114
	442 NEW PLACE+X,TRY STATE+X+APPLIED.	115
	443 REDEFINE+X+NSTATE+NSORT+NTYPE=OLD POS:	115
	444 REDEFINE1+X+NSTATE+NSORT+NTYPE,NEW PLACE+X;	116
	445 MAKE+OLD POS +POS,NLCR,OUT+X IMPOSSIBLE,INFORM+X,	116
	446 OUT+NSTATE,OUT+NSORT,OUT+NTYPE,POSITION+OLD POS.	116
	447 REDEFINE1+X+NSTATE+NSORT+NTYPE:	116
	448 TRY TYPE+X+NTYPE,TRY SORT+X+NSORT,	117
	TRY STATE+X+NSTATE.	
BACKTRACK?		
BACKTRACK?		
	449 TRY STATE+X+S-Q:	117
	450 GET STATE+X+Q,EQUAL+Q+0,PUT STATE+X+S;	118
	451 EQUAL+Q+DEFINED,EQUAL+S+APPLIED,PUT STATE+X+BLANK;	118
	452 EQUAL+Q+APPLIED,	118
	453 (EQUAL+S+APPLIED;	118
	454 EQUAL+S+DEFINED,PUT STATE+X+BLANK);	118
	455 EQUAL+Q+BLANK ,EQUAL+S+APPLIED.	
	456 TRY SORT+X+P-Q:	118
	457 GET SORT+X+Q,EQUAL+Q+0,PUT SORT+X+P,	119
	458 (EQUAL+P+LOCAL,ENTER LOCAL+X;	119
	459 EQUAL+P+GLOBAL,ENTER GLOBAL+X;	119
	460 EQUAL+P+MACRO,ENTER GLOBAL+X);	119
	461 EQUAL+P+Q.	
	462 TRY TYPE+X+P-Q:	119
	463 GET TYPE+X+Q,EQUAL+Q+0,PUT TYPE+X+P;	120
	464 EQUAL+P+Q.	120
	465 NEW PLACE+X-SORT:	121
	466 GET SORT+X+SORT,EQUAL+SORT+LOCAL; ADD PLACE+X.	121
	467 NEW NPARS+X+P-Q:	121
	468 GET NPARS+X+Q,EQUAL+Q+P;	122
	469 EQUAL+Q+999,PUT NPARS+X+P;	122
	470 ERROR+WRONG NUMBER OF PARAMETERS+X.	122
	471 WAS ACTION+HEAD-T:	122
		123

	472	GET TYPE+HEAD+T,EQUAL+T+ACTION.	
	473	WAS AFFIX+HEAD-T:GET TYPE+HEAD+T,	123
	474	EQUAL+T+POINTER;	124
	475	EQUAL+T+FLAG;	124
	476	EQUAL+T+LIST,	
	477	WAS MACRO+HEAD-S:	124
	478	GET SORT+HEAD+S,EQUAL+S+MACRO.	125
	479	WAS PARAMLESS MACRO+HEAD:	125
	480	WAS TAG+HEAD,WAS MACRO+HEAD,WAS AFFIX+HEAD,	126
BACKTRACK?			
	481		126
	482	[4.4,3 SPECIFICATIONS]	126
	483	'ACTION'TREAT SPEC LIST,READ MACRO,	127
	484	SPECIFICATION:	127
	485	EXTERNAL SPECIFICATION;	128
	486	INTERNAL SPECIFICATION;	128
	487	MACRO SPECIFICATION,	128
	488	EXTERNAL SPECIFICATION-TYPE:	128
	489	R+EXTERNAL,IS TYPE+TYPE,	129
BACKTRACK?		TREAT SPEC LIST+DEFINED+GLOBAL+TYPE,	
	490	IS TYPE+X:MAKE+X+INPT,	129
	491	R+ACTION;	130
	492	R+PREDICATE;	130
	493	R+POINTER;	130
	494	R+FLAG,	130
	495	TREAT SPEC LIST+STATE+SORT+TYPE-HEAD:	130
	496	NXT: IS TAG+HEAD,REDEFINE+HEAD+STATE+SORT+TYPE,	131
	497	(R+COMMA,:NXT;SHIFT2LINES,R+POINT).	131
	498	INTERNAL SPECIFICATION-TYPE:	131
	499	IS TYPE OF LOCAL+TYPE,TREAT SPEC LIST+APPLIED+GLOBAL+TYPE,	132
	500	IS TYPE OF LOCAL+X:MAKE+X+INPT,	132
	501	R+ACTION;	133
	502	R+PREDICATE,	133
	503	MACRO SPECIFICATION-TYPE-HEAD:	133
	504	R+MACRO,IS TYPE+TYPE,	134
BACKTRACK?			134
BACKTRACK?	505	NXT: IS TAG+HEAD,	
		REDEFINE+HEAD+DEFINED+MACRO+TYPE,READ MACRO+HEAD,	134

	506	(R+COMMA,:NXT;SHIFT2LINES,R+POINT),	
	507	READ MACRO+HEAD-X:	134
	508	PUT MTEXT+HEAD+PMACR,	135
	509	NXT: (AHEAD + POINT, ENTER MACRO + POINT);	135
	510	AHEAD + COMMA, ENTER MACRO + POINT;	135
	511	IS TAG+X,TRY STATE+X+APPLIED,	
BACKTRACK?	512	ENTER MACRO + INPT, NEXT SYMBOL, :NXT);	135
	513		135
	514	[4.4.4 DECLARATIONS]	135
	515	'ACTION'TREAT DECL LIST,CONSTANT TEXT.	136
	516	'ACTION'PUT PARAMLESS MACRO.	136
	517	DECLARATION:	136
	518	POINTER DECLARATION;	137
	519	FLAG DECLARATION;	137
	520	LIST DECLARATION.	137
	521	POINTER DECLARATION:	137
	522	R+POINTER,TREAT DECL LIST+Q INTEGER+POINTER.	138
	523	TREAT DECL LIST+ALGOLTYPE+TYPE-HEAD:	138
	524	BLANK LINE,G+ALGOLTYPE,	139
	525	NXT: IS TAG+HEAD,REDEFINE+HEAD+DEFINED+GLOBAL+TYPE,G+HEAD,	139
	526	(R+COMMA,G+COMMA,:NXT;	139
	527	SHIFT2LINES,R+POINT,G+SEMICOLON).	
	528	FLAG DECLARATION:	139
	529	R+FLAG,TREAT DECL LIST+Q BOOLEAN+FLAG.	140
	530	LIST DECLARATION-HEAD:	140
	531	R+LIST,BLANK LINE,G+Q INTEGER,G+Q ARRAY,U,	141
	532	NXT: IS TAG+HEAD,	141
BACKTRACK?	533	R+SUB,	141
BACKTRACK?	534	G+SUB,CONSTANT TEXT,	141
BACKTRACK?	535	R+COLON,	141
BACKTRACK?	536	G+COLON,CONSTANT TEXT,	141
BACKTRACK?	537	R+BUS,	141
	538	G+BUS,	141
	539	(R+COMMA,G+COMMA,NEW LINE,:NXT;	141
	540	SHIFT2LINES,R+POINT,G+SEMICOLON,L).	141
	541		141
	542	CONSTANT TEXT - SYMB:	142
	543	NXT: IS TAG+ SYMB,APPLY+SYMB,	142
	544	(WAS PARAMLESS MACRO+SYMB,PUT PARAMLESS MACRO+SYMB,:NXT);	142
	545	G + SYMB, :NXT);	142
	546	IS CONS + SYMB, G + SYMB, :NXT;	142

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543 R + PLUS, G + PLUS, :NXT; 142
544 R + MINUS, G + MINUS, :NXT;
545 142
546 [4.4.5 AFFIX EXPRESSIONS] 142
547 'ACTION'PUTAFFIX EXPRESSION,PUT DIRECT,AFFIX PACK OPTION,TREAT AFFIX,PUT SINGLE, 143
548 PUT PARAMLESS MACRO,PUT MACRO,GET PARAMETERS,GET PAR,PUT PAR.
549 PUT AFFIX EXPRESSION+HEAD: 143
550 TRY SORT+HEAD+GLOBAL,PUT DIRECT+HEAD; 144
551 TRY SORT+HEAD+MACRO, 144
552 (WAS ACTION+HEAD,PUT MACRO+HEAD; 144
553 G+OPEN,PUT MACRO+HEAD,G+CLOSE),
554 PUT DIRECT+HEAD-NPARS: 144
555 MAKE+NPARS+0,G+HEAD,AFFIX PACK OPTION+NPARS,NEW NPARS+HEAD+NPARS, 145
556 AFFIX PACK OPTION+NPARS: 145
557 R + PLUS, G + OPEN, TREAT AFFIX, 146
558 RST: (INCR+NPARS,R+PLUS,G+COMMA,TREAT AFFIX,:RST; 146
559 G + CLOSE);
560 TREAT AFFIX - AFFX: 146
561 IS TAG + AFFX, PUT SINGLE + AFFX,APPLY+AFFX; 147
562 IS BOLD+AFFX,G+STRINGQUOTE,G+AFFX,G+STRINGQUOTE; 147
563 IS CONS + AFFX, G + AFFX; 147
564 ERROR+WRONG AFFIX+AFFX.
565 PUT SINGLE+X: 147
566 WAS AFFIX+X, 148
567 (WAS MACRO+X,PUT PARAMLESS MACRO+X; G+X); 148
568 REDEFINE+X+APPLIED+GLOBAL+POINTER,G+X.
569 PUT PARAMLESS MACRO+Y-X-M: 148
570 GET MTEXT+Y+X,NEW NPARS+Y+0, 149
571 NXT: INCR + X, GET + LMACR + X + M, 149
572 (EQUAL + M + POINT; G + M, :NXT).
573 PUT MACRO+X-P1-P2-P3-P4-P5-Q-M-NPARS: 149
574 MAKE+NPARS+0,GET MTEXT+X+Q, 150
575 GET PARAMETERS+NPARS+P1+P2+P3+P4+P5, NEW NPARS+X+NPARS, 150
576 NXT: INCR + Q, GET + LMACR + Q + M, 150
577 (EQUAL + M + POINT; 150
578 EQUAL+M+ONE ,PUT PAR+P1+1,:NXT; 150
579 EQUAL+M+TWO ,PUT PAR+P2+2,:NXT; 150
580 EQUAL+M+THREE,PUT PAR+P3+3,:NXT; 150
581 EQUAL+M+FOUR ,PUT PAR+P4+4,:NXT; 150
582 EQUAL+M+FIVE ,PUT PAR+P5+5,:NXT; 150
583 WAS PARAMLESS MACRO+M,PUT PARAMLESS MACRO+M,:NXT; 150
584 G + M, :NXT).

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585	GET PARAMETERS+P+P1+P2+P3+P4+P5:	151
586	GET PAR+P+P1,GET PAR+P+P2,GET PAR+P+P3,GET PAR+P+P4,GET PAR+P+P5.	
587	GET PAR+P+X:	151
588	R+PLUS, INCR+P,	152
589	(IS TAG+X, APPLY+X; IS CONS+X);	152
590	MAKE + X + X PARAMETER ERROR.	152
591	PUT PAR + X + Y:	152
592	WAS CONS+X, G+X;	153
593	EQUAL+X+X PARAMETER ERROR, ERROR+X+Y;	153
594	WAS BOLD+X, G+STRINGQUOTE, G+X, G+STRINGQUOTE;	153
595	PUT SINGLE + X.	
596		153
597	[4,4,6 BUILDING STONES OF A RULE]	153
598	'ACTION' PUT CONNECT AND, PUT CONNECT OR, END JUMP, END LABEL, LOST LABEL, PUT INIT,	154
599	PUT RESTORE, FRESH LABEL, TRUE RESULT, FALSE RESULT, PUT JUMP, PUT LABEL, SEMICOL,	154
600	BEGIN, ENDBR, BECOMES.	
601	PUT CONNECT AND:	154
602	GIVE RSTO, PUT JUMP+CUR LAB, INCR+TO NEXT;	155
603	ADVANCED, ERROR+POSSIBLY BACKTRACK NECESSARY+0, PUT JUMP+0, INCR+TO LOST;	155
604	PUT JUMP+CUR LAB, INCR+TO NEXT.	
605	PUT CONNECT OR:	155
606	AHEAD+SEMICOLON, PUT RESTORE, FRESH LABEL,	156
607	(EQUAL+TO NEXT+0, ERROR+ALTERNATIVE NEVER REACHED+0;)	156
608	AHEAD+CLOSE,	156
609	(EQUAL+TO NEXT+0;	156
610	PUT LABEL+CUR LAB, FRESH LABEL,	156
611	(GIVE RSTO; PUT JUMP+0, INCR+TO LOST));	156
612	PUT RESTORE, LOST LABEL.	
613	LOST LABEL:	156
614	EQUAL+TO LOST+0,	157
615	(EQUAL+TO NEXT+0; FALSE RESULT);	157
616	PUT LABEL+0, FALSE RESULT.	
617	END LABEL:	157
618	EQUAL+TO END+0; PUT LABEL+999.	158
619	END JUMP:	158
620	EQUAL+INPT+POINT, EQUAL+TO LOST+0, EQUAL+TO NEXT+0;	159
621	PUT JUMP+999, INCR+TO END.	159
622	PUT .INIT:	159
623	GIVE RSTO, G+INIT RESTORE;	160
624	PUT RESTORE:	160
		161

	625	EQUAL+TO NEXT+0:	161
	626	PUT LABEL+CUR LAB,	161
	627	(GIVE RSTO,G+DO RESTORE;),	
	628	FRESH LABEL:	161
	629	MAKE+CUR LAB+NEW LAB,INCR+NEW LAB,	162
	630	TRUE RESULT:	162
	631	WAS ACTION+HANDLE;	163
	632	G+HANDLE,BECOMES,G+Q TRUE,SEMICOL.	163
	633	FALSE RESULT:	163
	634	WAS ACTION+HANDLE;	164
	635	G+HANDLE,BECOMES,G+Q FALSE,SEMICOL.	164
	636	PUT JUMP + LAB:	164
	637	G+Q GOTO, G + LAB, G + SEMICOLON, NEW LINE,APPLY LABEL+LAB.	165
	638	PUT LABEL + LAB:	165
	639	NEW LINE,L,G+LAB,G+COLON,U,DEFINE LABEL+LAB,TAB LINE.	166
	640	SEMICOL: G+SEMICOLON.	166
	641	BECOMES:	167
	642	G+COLON,G+EQUALS.	168
	643	BEGIN: NEW LINE, G + Q BEGIN,U,TAB LINE.	168
	644	ENDBR: NEW LINE, L, G + Q END.	169
	645		170
	646	[4.4.7 GENERAL FORM OF A RULE]	170
	647	'ACTION'OPTIONAL BOUND AFFIXES,REST BOUND AFFIXES,OPTIONAL FREE AFFIXES,RESTLHS.	171
	648	RULE:	171
	649	LEFT HAND SIDE,MIDDLE, RIGHT HAND SIDE,	172
BACKTRACK?	650	FORGET LOCALS,SHIFT2LINES,R+POINT.	172
BACKTRACK?	651	LEFT HAND SIDE:	172
	652	IS TAG+HANDLE,BLANK LINE,MAKE+BOUNDX+0,MAKE+SPECX+0,	173
	653	(WAS ACTION+HANDLE,DEFINE ACTION+HANDLE,REST LHS;	173
	654	G+Q BOOLEAN,DEFINE PREDICATE+HANDLE,REST LHS).	173
	655	RFST LHS:	173
			174

	656	G+Q. PROCEDURE,G+HANDLE,OPTIONAL BOUND AFFIXES,NEW NPARS+HANDLE+BOUNDX,	174
	657	SEMICOL.	
	658	OPTIONAL BOUND AFFIXES-AFFX:	174
	659	BOUND AFFIX1+AFFX,G+OPEN,G+AFFX,REST BOUND AFFIXES;	175
	660	BOUND AFFIX2+AFFX,G+OPEN,G+AFFX,REST BOUND AFFIXES,G+AFFX;	175
	661	REST BOUND AFFIXES-AFFX:	175
	662	BOUND AFFIX1+AFFX,G+COMMA,G+AFFX,REST BOUND AFFIXES;	176
	663	BOUND AFFIX2+AFFX,G+COMMA,G+AFFX,REST BOUND AFFIXES,G+AFFX,DECR+SPECX,	176
	664	(EQUAL+SPECX+0;G+COMMA);	176
	665	G+CLOSE,	176
	666	(EQUAL+SPECX+0;SEMICOL,G+Q INTEGER).	
	667	BOUND AFFIX1+AFFX:	176
	668	R+TIMES,IS TAG+AFFX,	177
BACKTRACK?		DEFINE AFFIX+AFFX,INCR+BOUNDX.	
	669	BOUND AFFIX2+AFFX:	177
	670	R+PLUS,IS TAG+AFFX,	178
BACKTRACK?		DEFINE AFFIX+AFFX,INCR+BOUNDX,INCR+SPECX.	
	671	MIDDLE:	178
	672	BEGIN,OPTIONAL FREE AFFIXES,R+COLON,PUT INIT.	179
	673	OPTIONAL FREE AFFIXES - AFFX:	179
	674	R+MINUS,G+Q INTEGER,	180
	675	FFX; IS TAG+AFFX,	180
BACKTRACK?		G+AFFX,DEFINE AFFIX+AFFX,	180
	676	(R+MINUS,G+COMMA,:FFX;	180
	677	G+SEMICOLON);	
	678	RIGHT HAND SIDE:	180
	679	MAKE+CUR LAB+1,MAKE+NEW LAB+2,MAKE+TO LOST+0,MAKE+TO END+0,	181
	680	NXT: ALTERNATIVE,PUT CONNECT OR,	181
	681	(R+SEMICOLON,:NXT;	181
	682	END LABEL,END BR,SEMICOL).	
	683	ALTERNATIVE-LAB-OLD NEXT:	181
	684	MAKE+TO NEXT+0,SET+ADVANCED+FALSE,	182
	685	NXT: R+COMMA,:NXT;	182
	686	R+COLON,IS TAG+LAB,	
BACKTRACK?		PUT JUMP+LAB;	182
	687	R+OPEN,MAKE+OLD NEXT+TO NEXT,CHOICE,	
BACKTRACK?		R+CLOSE,	
BACKTRACK?			
	688	(GIVE RSTO,MAKE+TO NEXT+1;MAKE+TO NEXT+OLD NEXT);	182
	689	MEMBER,:NXT;	182
	690	TRUE RESULT,END JUMP.	
	691	CHOICE-LAB:	182
			183

692	MAKE+LAB+CUR LAB,BEGIN,PUT INIT,FRESH LABEL,	183
693	NXT: ALTERNATIVE,PUT CONNECT OR,	183
694	(R+SEMICOLON,;NXT;ENDBR,SEMICOL,MAKE+CUR LAB+LAB).	
695	MEMBER-HEAD:	183
696	IS TAG+HEAD,APPLY+HEAD,	184
697	(R+COLON,PUT LABEL+HEAD;	184
698	TRY TYPE+HEAD+PREDICATE,	184
699	G+QIF,G+NOT,PUT AFFIX EXPRESSION+HEAD,G+Q THEN,	184
700	PUT CONNECT AND,SET+ADVANCED+TRUE;	184
701	TRY TYPE+HEAD+FLAG,	184
702	G+QIF,G+NOT,PUT AFFIX EXPRESSION+HEAD,G+Q THEN,	184
703	PUT CONNECT AND;	184
704	TRY TYPE+HEAD+ACTION,	184
705	PUT AFFIX EXPRESSION+HEAD,SEMICOL).	
706		184
707	[4.4.8 OTHER BUILDING STONES OF A GRAMMAR]	184
708	'ACTION'SKIP UNTIL POINT,TERMINALS.	185
709	COMMAND:	185
710	R+RSTO ON,SET+GIVE RSTO+TRUE;	186
711	R+RSTO OFF,SET+GIVE RSTO+FALSE;	186
712	R+SHORT,MAKE+LEGIBLE+FALSE;	186
713	R+LONG,MAKE+LEGIBLE+TRUE;	186
714	R + TRACE ON, SET + GIVE TRACE + TRUE;	186
715	R + TRACE OFF, SET + GIVE TRACE + FALSE.	
716	COMMENT:	186
717	R+SUB,BLANK LINE,G+Q COMMENT,	187
718	RST: (R + BUS, G + SEMICOLON;	187
719	G+INPT,READ+INPT,:RST).	
720		187
721	STARTING SYMBOL-HEAD:	187
722	R+RESULT,IS TAG+HEAD,	188
723	NEW PAGE,	188
724	(TRY TYPE+HEAD+ACTION,APPLY ACTION+HEAD,TERMINALS,G+HEAD;	188
	APPLY PREDICATE+HEAD,TERMINALS,G+HEAD).	
725	SKIP UNTIL POINT:	188
726	NXT: ERROR + SKIPPED + INPT,	189
727	(R + POINT; NEXT SYMBOL,:NXT).	189
728		189
729	[4.4.9 TERMINALS]	190
730	'ACTION' TERMINALS,FORGET GLOBALS,MAY BE TERMINAL,PUT DECL,FORGET TERMINALS,	190
731	PUT CALL.	
732	TERMINALS:	190
733	BLANK LINE,FORGET GLOBALS,NEW LINE,EQUAL+PLOC+MIN LOC;	191
734	APPLY ACTION+XREAD,APPLY ACTION+XOUT,APPLY ACTION+INIT READ,	191

BACKTRACK?

735	BLANK LINE, G + INIT READ, G + SEMICOLON, NEW LINE,	191
736	BLANK LINE, FORGET TERMINALS, NEW LINE, BLANK LINE.	
737	FORGET GLOBALS-P-TERM:	191
738	MAKE+P+MIN GLOB, MAKE+PLOC+MIN LOC,	192
739	NXT: EQUAL + P + PGLOB;	192
740	GET+LGLOB+P+TERM, MAY BE TERMINAL+TERM, INCR+P, ;NXT.	192
741	MAY BE TERMINAL+X:	193
742	WAS TERMINAL+X, PUT DECL+X, ENTER LOCAL+X;	193
743	WARN+X.	
744	WAS TERMINAL+X-S:	193
745	GET STATE+X+S, EQUAL+S+APPLIED,	194
746	GET SORT +X+S, EQUAL+S+GLOBAL,	194
747	GET TYPE+X+S, EQUAL+S+POINTER.	194
748	PUT DECL + X:	195
749	NEW LINE, G+Q INTEGER, G+X, G+SEMICOLON, POSITION+64, INFORM+X.	
750	FORGET TERMINALS-P-TERM:	195
751	MAKE+P+MIN LOC,	196
752	NXT: EQUAL+P+PLOC;	196
753	GET+LLOC+P+TERM, PUT CALL+X READ+TERM, INCR+P, ;NXT.	196
754	PUT CALL +PROC + X:	197
755	NEW LINE, G + PROC, G + OPEN, G + X, G + CLOSE, G + SEMICOLON.	
756		197
757	[4,4,10 POST MORTEM]	197
758	'ACTION' POST MORTEM, DICTIONARY, LIST TAGS, ENTRIES.	198
759	POST MORTEM:	198
760	NLCR, OUTINT+MIN TAG , OUTINT+PTAG , OUTINT+MAX TAG ,	199
761	NLCR, OUTINT+MIN BOLD, OUTINT+PBOLD, OUTINT+MAX BOLD,	199
762	NLCR, OUTINT+MIN SPEC, OUTINT+PSPEC, OUTINT+MAX SPEC,	199
763	NLCR, OUTINT+MIN CONS, OUTINT+PCONS, OUTINT+MAX CONS,	199
764	NLCR, OUTINT+MIN LOC , OUTINT+XLOC , OUTINT+MAX LOC ,	199
765	NLCR, OUTINT+MIN GLOB, OUTINT+PGLOB, OUTINT+MAX GLOB,	199
766	NLCR, OUTINT+MIN MACR, OUTINT+PMACR, OUTINT+MAX MACR,	199
767	NLCR, OUTINT+MIN TEXT, OUTINT+PTEXT, OUTINT+MAX TEXT,	199
768	NLCR, OUTINT+LINE, OUTINT+CARD, NLCR, NLCR, NLCR, DICTIONARY.	
769	DICTIONARY:	199
770	LIST TAGS + FIRST TAG.	200
771	LIST TAGS + X - P-STATE-SORT-TYPE:	201
772	EQUAL + X + 0;	201
773	GET LEFT+X+P, LIST TAGS+P,	201
774	GET STATE+X+STATE, GET SORT+X+SORT, GET TYPE+X+TYPE.	201

775	EQUAL+SORT+0, :SKP;	201
776	NLCR, OUT+X, POSITION+32, OUT+STATE, OUT+SORT, OUT+TYPE,	201
777	POSITION+59, ENTRIES+X,	201
778	SKP: GET RIGHT+X+R, LIST TAGS+P.	
779	ENTRIES+P-Q-H-K:	201
780	GET PLACE+P+Q, EQUAL+Q+0, NLCR;	202
781	LNE: GET+LTEXT+Q+H, UNPACK2+H+Q+K, OUTINT+K,	202
782	(EQUAL+Q+0, NLCR;	202
783	ADD+Q+MIN TEXT+Q, DECR+Q, :LNE).	
784		202
785	[4.5... HEART OF COMPILER COMPILER]	202
786	'ACTION' SENTENCE, START SYSTEM, COMPILER DESCRIPTION.	203
787	'POINTER' INIT RESTORE, DO RESTORE, XREAD, XOUT, INIT READ, TITLE.	203
788	SENTENCE:	204
789	START SYSTEM, BLANK LINE,	205
790	BEGIN, COMPILER DESCRIPTION, ENDBR,	205
791	BLANK LINE, NEW PAGE, POST MORTEM.	205
792	START SYSTEM:	205
793	READ+INIT RESTORE, READ+DO RESTORE, READ+X READ, READ+X OUT,	206
794	READ+INIT READ, READ+TITLE,	206
795	SET+GIVERSTO+FALSE, SET+GIVE TRACE+FALSE, SET+GIVE TEXT+TRUE,	206
796	MAKE+PLOC+MIN LOC, MAKE+PGLOB+MIN GLOB, MAKE+PTEXT+MIN TEXT,	206
797	MAKE+PMACR+MIN MACR, MAKE+CARD+0, MAKE+CPOS+1, MAKE+XLOC+0,	206
798	SET+LINED UP+FALSE, SET+LEGIBLE+TRUE,	206
799	MAKE+INDENTATION+0, MAKE+LINE+0, MAKE+POS+0, OUT+TITLE, READ+INPT,	
800	COMPILER DESCRIPTION:	206
801	NXT: SPECIFICATION, :NXT;	207
802	DECLARATION, :NXT;	207
803	COMMAND, :NXT;	207
804	COMMENT, :NXT;	207
805	STARTING SYMBOL;	207
806	RULE, :NXT;	207
807	SKIP UNTIL POINT, :NXT,	
808		207
809	'RESULT' SENTENCE.	207

APPLIED GLOBAL POINTER TAGFULL
APPLIED GLOBAL POINTER BOLDFULL
APPLIED GLOBAL POINTER SPECFULL
APPLIED GLOBAL POINTER CONSFULL
APPLIED GLOBAL POINTER WRONGINPUTCODE
APPLIED GLOBAL POINTER BLANK
APPLIED GLOBAL POINTER LOCFULL
APPLIED GLOBAL POINTER GLOBFULL
APPLIED GLOBAL POINTER MACRFULL
APPLIED GLOBAL POINTER TEXTFULL
APPLIED GLOBAL POINTER DEFINED
APPLIED GLOBAL POINTER GLOPAL
APPLIED GLOBAL POINTER ACTION
APPLIED GLOBAL POINTER PREDICATE
APPLIED GLOBAL POINTER LOCAL
APPLIED GLOBAL POINTER POINTER
APPLIED GLOBAL POINTER LABEL
APPLIED GLOBAL POINTER APPLIED
APPLIED GLOBAL POINTER XIMPOSSIBLE
APPLIED GLOBAL POINTER MACRO
APPLIED GLOBAL POINTER WRONGNUMBEROFFPARAMETERS
APPLIED GLOBAL POINTER FLAG
APPLIED GLOBAL POINTER LIST
APPLIED GLOBAL POINTER EXTERNAL
APPLIED GLOBAL POINTER COMP'A
APPLIED GLOBAL POINTER POINT
APPLIED GLOBAL POINTER QINTEGER
APPLIED GLOBAL POINTER SEMICOLON
APPLIED GLOBAL POINTER QBOCLEAN
APPLIED GLOBAL POINTER QARPAY
APPLIED GLOBAL POINTER SUB
APPLIED GLOBAL POINTER COLON
APPLIED GLOBAL POINTER BUS
APPLIED GLOBAL POINTER PLUS
APPLIED GLOBAL POINTER MINUS
APPLIED GLOBAL POINTER OPEN
APPLIED GLOBAL POINTER CLOSE
APPLIED GLOBAL POINTER STRINGQUOTE
APPLIED GLOBAL POINTER WRONGAFFIX
APPLIED GLOBAL POINTER ONE
APPLIED GLOBAL POINTER TWO
APPLIED GLOBAL POINTER THREE
APPLIED GLOBAL POINTER FOUR
APPLIED GLOBAL POINTER FIVE
APPLIED GLOBAL POINTER XPARAMETERERROR
APPLIED GLOBAL POINTER POSSIBLYBACKTRACKNECESSARY
APPLIED GLOBAL POINTER ALTERNATIVEEVERREACHED
APPLIED GLOBAL POINTER QTRUE
APPLIED GLOBAL POINTER QFALSE
APPLIED GLOBAL POINTER QGOTO
APPLIED GLOBAL POINTER EQUALS
APPLIED GLOBAL POINTER QBEGIN
APPLIED GLOBAL POINTER QEND
APPLIED GLOBAL POINTER QPROCEDURE
APPLIED GLOBAL POINTER TIMES
APPLIED GLOBAL POINTER QIF
APPLIED GLOBAL POINTER NOT
APPLIED GLOBAL POINTER QTHEN
APPLIED GLOBAL POINTER RSTOON

APPLIED GLOBAL POINTER RSTOOF
APPLIED GLOBAL POINTER SHORT
APPLIED GLOBAL POINTER LONG
APPLIED GLOBAL POINTER TRACEON
APPLIED GLOBAL POINTER TRACEOFF
APPLIED GLOBAL POINTER QCOMMENT
APPLIED GLOBAL POINTER RESULT
APPLIED GLOBAL POINTER SKIPPED

NEWNPARS	GLOBAL ACTION	656	575	570	555	467	424												
NEWPAGE	GLOBAL ACTION	791	722	12															
NEWPLACE	GLOBAL ACTION	465	444	442	424														
NEXTCHAR	GLOBAL ACTION	376	371	255	247	244	243	242	234										
NEXTNONSPACECHAR	GLOBAL ACTION	254	246	240	238	236	234												
NEXTSYMBCL	GLOBAL ACTION	727	512	368	366	365	364	363	358										
NIX	MACRO POINTER	328	306	304	302	272	271	199	143	20									
NLCRCGDE	MACRO POINTER	259	250	201	197	190	141	134	20										
NLCR 125 99	GLOBAL ACTION	782	780	776	768	767	766	765	764	763	762	761	760	445	156	141	136	133	
NOT	APPLIED GLOBAL POINTER	702	699																
ONE	APPLIED GLOBAL POINTER	578																	
OPEN	APPLIED GLOBAL POINTER	755	687	660	659	557	553												
OPTIONALBOUNDRAFFIXES	GLOBAL ACTION	658	656	647															
OPTIONALFREEAFFIXES	GLOBAL ACTION	673	672	647															
CUT	GLOBAL ACTION	809	799	776	446	445	159	156	154	102	99								
CUTINT	GLOBAL ACTION	781	768	767	766	765	764	763	762	761	760	110	108	107	106	99			
CUTINT1	GLOBAL ACTION	252	251	115	113	112	99												
PACK2	MACRO ACTION	397	60																
PACK3	MACRO ACTION	328	307	306	304	302	272	271	270	268	58								
PBOLD	GLOBAL POINTER	761	361	320	311	301	87												
PCCNS	GLOBAL POINTER	763	360	354	347	345	87												
PGLOB	GLOBAL POINTER	796	765	739	412	88													
PLCC	GLOBAL POINTER	796	752	738	733	409	408	400	399	88									
PLUS	APPLIED GLOBAL POINTER	670	588	558	557	543													
PMACR	GLOBAL POINTER	797	766	508	415	88													
POINTERDECLARATION	GLOBAL PREDICATE	521	518																
POINTER	APPLIED GLOBAL POINTER	747	568	522	493	474	430												
POINT	APPLIED GLOBAL POINTER	727	650	620	577	572	537	527	510	509	506	497							
POSITION	GLOBAL ACTION	777	776	749	446	252	251	159	158	157	151	149	123	100					

POS	GLOBAL POINTER	799	445	156	144	137	136	134	124	101				
POSSIBLYBACKTRACKNECESSARY	APPLIED GLOBAL POINTER	603												
POSTMORTEM	GLOBAL ACTION	791	759	758	419	416	413	410	348	331	312	296	276	229
PRCHAR	GLOBAL ACTION	253	140	139	122	115	112	111	110	99				
PREDICATE	APPLIED GLOBAL POINTER	698	502	492	437	428								
PRINT	GLOBAL ACTION	120	119	99										
PRINT	GLOBAL ACTION	116	105	104	103	99								
PREYM	GLOBAL ACTION	144	137	134	12									
PSPEC	GLOBAL POINTER	762	361	337	330	328	87							
PTAG	GLOBAL POINTER	760	362	361	295	294	293	284	275	265	87			
PTEXT	GLOBAL POINTER	796	767	418	395	88								
PUCHAR	GLOBAL ACTION	201	196	195	187	179	176	175	174	163				
PUNCH1	GLOBAL ACTION	185	184	183	162									
PUNCH	GLOBAL ACTION	180	170	169	168	162								
PUT	MACRO ACTION	418	415	412	408	347	330	311	275	50				
PUTABS	GLOBAL ACTION	214	212	211	204									
PUTAFFIXEXPRESSION	GLOBAL ACTION	705	702	699	549	547								
PUTCALL	GLOBAL ACTION	754	753	731										
PUTCONNECTAND	GLOBAL ACTION	703	700	601	598									
PUTCONNECTCR	GLOBAL ACTION	693	680	605	598									
PUTDECL	GLOBAL ACTION	748	742	730										
PUTDIRECT	GLOBAL ACTION	554	550	547										
PUTINIT	GLOBAL ACTION	692	672	622	598									
PUTJUMP	GLOBAL ACTION	686	636	621	611	604	603	602	599					
PUTLABEL	GLOBAL ACTION	697	638	626	618	616	610	599						
PUTLEFT	MACRO ACTION	295	287	82										
PUTMACRO	GLOBAL ACTION	573	553	552	548									
PUTMTEXT	MACRO ACTION	508	294	77										
PUTNPARS	MACRO ACTION	469	293	76										
PUTPARAMLESSMACRO	GLOBAL ACTION	583	569	567	548	540	516							

STACKSPEC	GLOBAL ACTION	329	328	326					
STACKTAG	GLOBAL ACTION	274	273	269	263				
STARTINGSYMBCL	GLOBAL PREDICATE	805	721						
STARTSYSTEM	GLOBAL ACTION	792	789	786					
STRINGOJCTE	APPLIED GLOBAL POINTER	594	562						
SUPTR	MACRO ACTION	397	124	55					
SUP	APPLIED GLOBAL POINTER	717	533						
TABCARD	GLOBAL ACTION	221	214	198	191	162			
TABCCDE	MACRO POINTER	260	198	142	137	20			
TAB	GLOBAL ACTION	142	137	132	99				
TABLIE	GLOBAL ACTION	643	639	221	204				
TABS	GLOBAL ACTION	130	125	99					
TAGFULL	APPLIED GLOBAL POINTER	296	276						
TBOLD	GLOBAL LIST	323	316	311	169	104	34		
TCCNS	GLOBAL LIST	351	350	347	171	106	36		
TERMINALS	GLOBAL ACTION	732	730	724	723	708			
TEXTFULL	APPLIED GLOBAL POINTER	419							
THREE	APPLIED GLOBAL POINTER	580							
TINES	APPLIED GLOBAL POINTER	668							
TITLE	GLOBAL POINTER	799	794	787					
TOEND	GLOBAL POINTER	679	621	618	384				
TOLOST	GLOBAL POINTER	679	620	614	611	603	384		
TOEXT	GLOBAL POINTER	688	687	684	625	620	615	609	607
TRACEOFF	APPLIED GLOBAL POINTER	715							
TRACEON	APPLIED GLOBAL POINTER	714							
TRACE	GLOBAL ACTION	405	404	148	147				
TREATAFFIX	GLOBAL ACTION	560	558	557	547				
TREATDECLIST	GLOBAL ACTION	529	523	522	515				
TREATSPECLIST	GLOBAL ACTION	499	495	489	483				
TRIERESULT	GLOBAL ACTION	690	630	599					

TRUE	MACRO POINTER	798	795	714	713	710	700	211	73
TRYS CRT	GLOBAL PREDICATE	551	550	456	448				
TRYSTATE	GLOBAL PREDICATE	511	449	448	442				
TRYTYPE	GLOBAL PREDICATE	723	704	701	698	462	448		
TSPEC	GLOBAL LIST	334	333	330	170	105	35		
TTAG	GLOBAL LIST	280	275	168	103	33			
TWC	APPLIED GLOBAL POINTER	579							
UNPACK2	MACRO ACTION	781	61						
UNPACK3	MACRO ACTION	186	121	59					
U	GLOBAL ACTION	643	639	531	215	204			
WARN	GLOBAL ACTION	743	402	400	388				
WASACTION	GLOBAL PREDICATE	653	634	631	552	471			
WASAFFIX	GLOBAL PREDICATE	566	480	473					
WASBOLD	MACRO FLAG	594	364	169	104	43			
WASCCNS	MACRO FLAG	592	440	433	365	171	106	45	
WASDIGIT	MACRO FLAG	240	17						
WASLETGIT	MACRO FLAG	238	16						
WASLETTER	MACRO FLAG	236	15						
WASMACRO	GLOBAL PREDICATE	567	480	477					
WASPARAMLESSMACRO	GLOBAL PREDICATE	583	540	479					
WASSPECCH	MACRO FLAG	246	18						
WASSPEC	MACRO FLAG	170	105	44					
WASTAG	MACRO FLAG	480	439	432	363	168	157	103	42
WASTERMINAL	GLOBAL PREDICATE	744	742						
WRITEINT1	GLOBAL ACTION	179	177	176	162				
WRITEINT	GLOBAL ACTION	190	174	173	172	171	162		
WRITE	GLOBAL ACTION	208	167	162					
WRNGAFFIX	APPLIED GLOBAL POINTER	564							
WRNGINPUTCODE	APPLIED GLOBAL POINTER	376							
WRNGNUMBEROFPARAMETERS	APPLIED GLOBAL POINTER	470							

XIMPCSS:BLE.....APPLIED GLOBAL POINTER.....445.....
XLOC.....GLOBAL POINTER.....797 764 409 88.....
XOUT.....GLOBAL POINTER.....793 787 734.....
XPARAMETERERROR.....APPLIED GLOBAL POINTER.....593 590.....
XREAD.....GLOBAL POINTER.....793 787 753 734.....

A1262J.19,KCOSTER

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1  'BEGIN'
2  'COMMENT'4COMPILERCOMPILERDESCRIBEDINCOL*****;
3  'COMMENT'4,1GENERALENVIRONMENT;
4  'COMMENT'4,1,1INTERFACEWITHMACHINE;
5  'COMMENT'4,1,2STACKS;
6  'INTEGER'ARRAY'TTAG[100001:110000],TBOLD[200001:200600],TSPEC[300001:300100],TCONS[400001:400300],LLOC[500001:500200],LGLOBAL[600001:6010
00],LTEXT[800001:804000],LMACR[700001:701100];
7  'COMMENT'4,1,3MACROS;
8  'COMMENT'DATASTRUCTURESANDTHEIRACCESS;
9  'COMMENT'4,1,4POINTERSANDFLAGS;
10 'INTEGER'PTAG,PBOLD,PSPEC,PCONS,PLOC,PGLOBAL,PTEXT,PMACR,XLOC;
11 'EOCLEAN'GIVETEXT,GIVETRACE,LEGIBLE;
12 'INTEGER'FIRSTTAG;
13 'COMMENT'4,2OUTPUT;
14 'COMMENT'4,2,1PRINTERSECTION;
15 'INTEGER'POS;
16 'PROCEDURE'OUT(X);'INTEGER'X;'BEGIN''INTEGER'EL;'IF'-(100001'LE'X<X<PTAG)'THEN'GOTO'1;PRINT(TTAG,X);'GOTO'999;1:'IF'-(200001'LE'X<X<PBO
LD)'THEN'GOTO'2;PRINT(TBOLD,X);'GOTO'999;2:'IF'-(300001'LE'X<X<PSPEC)'THEN'GOTO'3;PRINT(TSPEC,X);'GOTO'999;3:'IF'-(400001'LE'X<X<PCONS)'THEN'
GOTO'4;EL:=TCONS[X];OUTINT(EL);'GOTO'999;4:OUTINT(X);999:'END';
17 'PROCEDURE'OUTINT(X);'INTEGER'X;'BEGIN''INTEGER'QUOT,REM;SPACE;'IF'-(X<0)'THEN'GOTO'1;REM:=X;REM:=-REM;PRCHAR(65);OUTINT(REM);'GOTO'999
;1:'IF'-(X=0)'THEN'GOTO'2;PRCHAR(0);SPACE;'GOTO'999;2:QUOT:=X/'10;REM:=X-10*QUOT;OUTINT1(QUOT);PRCHAR(REM);SPACE;999:'END';
18 'PROCEDURE'OUTINT1(X);'INTEGER'X;'BEGIN''INTEGER'QUOT,REM;'IF'-(X=0)'THEN'GOTO'1;'GOTO'999;1:QUOT:=X/'10;REM:=X-10*QUOT;OUTINT1(QUOT);
PRCHAR(REM);999:'END';
19 'PROCEDURE'PRINT(LIST,Y);'INTEGER'Y;'BEGIN''INTEGER'X,LX;X:=Y;SPACE;RST:LX:=LIST[X];'BEGIN''IF'-(LX<0)'THEN'GOTO'2;LX:=-LX;PRINT1(LX);'
GOTO'999;2:PRINT1(LX);X:=X+1;'GOTO'RST;'END';999:'END';
20 'PROCEDURE'PRINT1(X);'INTEGER'X;'BEGIN''INTEGER'X1,X2,X3;X1:=X/'16384;X3:=X-16384*X1;X2:=X3/'128;X3:=X3-128*X2;PRCHAR(X1);PRCHAR(X2);P
RCHAR(X3);'END';
21 'PROCEDURE'POSITION(X);'INTEGER'X;'BEGIN''INTEGER'TBS,SPCES;SPCES:=X-POS;'IF'-(0<SPCES)'THEN'GOTO'1;SPACES(SPCES);'GOTO'999;1:NLCR;TBS:
=X/'8;SPCES:=X-8*TBS;TABS(TBS);'BEGIN''IF'-(SPCES=0)'THEN'GOTO'3;'GOTO'999;3:SPACES(SPCES);'GOTO'999;'END';999:'END';
22 'PROCEDURE'SPACES(X);'INTEGER'X;'BEGIN''INTEGER'N;N:=0;SPC;'IF'-(N<X)'THEN'GOTO'1;SPACE;N:=N+1;'GOTO'SPC;1:'END';
23 'PROCEDURE'TABS(X);'INTEGER'X;'BEGIN''INTEGER'N;N:=0;TBS;'IF'-(N<X)'THEN'GOTO'1;TAB;N:=N+1;'GOTO'TBS;1:'END';
24 'PROCEDURE'NLCR;'BEGIN'PRSYM(119);POS:=0;'END';
25 'PROCEDURE'SHIFT2LINES;'BEGIN''INTEGER'OLDPOS;OLDPOS:=POS;NLCR;NLCR;SPACES(OLDPOS);'END';
26 'PROCEDURE'TAB;'BEGIN'PRSYM(118);POS:=POS+8;'END';
27 'PROCEDURE'SPACE;'BEGIN'PRCHAR(93);'END';
28 'PROCEDURE'PRCHAR(X);'INTEGER'X;'BEGIN''IF'-(X=119)'THEN'GOTO'1;NLCR;'GOTO'999;1:'IF'-(X=118)'THEN'GOTO'2;TAB;'GOTO'999;2:'IF'-(X=63)'
THEN'GOTO'3;'GOTO'999;3:PRSYM(X);PCS:=POS+1;999:'END';
29 'COMMENT'4,2,2TRACEADMINISTRATION;
30 'PROCEDURE'TRACE(X);'INTEGER'X;'BEGIN''IF'GIVETRACE'THEN'GOTO'1;POSITION(32);INFORM(X);'GOTO'999;1:999:'END';
31 'PROCEDURE'SIGNAL(X);'INTEGER'X;'BEGIN'POSITION(16);INFORM(X);'END';
32 'PROCEDURE'INFORM(X);'INTEGER'X;'BEGIN''INTEGER'STATE, SORT, TYPE;STATE:=TTAG[X-5];SORT:=TTAG[X-4];TYPE:=TTAG[X-3];OUT(STATE);OUT(SORT);OU
T(TYPE);OUT(X);'END';
33 'PROCEDURE'ERROR(TEXT,INFO);'INTEGER'INFO,TEXT;'BEGIN''INTEGER'OLDPOS;OLDPOS:=POS;NLCR;OUT(TEXT);'BEGIN''IF'-(100001'LE'INFO<INFO<PTAG)'
THEN'GOTO'2;INFORM(INFO);POSITION(OLDPOS);'GOTO'999;2:'IF'-(INFO=0)'THEN'GOTO'3;POSITION(OLDPOS);'GOTO'999;3:OUT(INFO);POSITION(OLDPOS);'GOTO'
999;'END';999:'END';
34 'COMMENT'4,2,3CARDPUNCHSECTION;
35 'INTEGER'CARD,CPOS;
36 'PROCEDURE'WRITE(X);'INTEGER'X;'BEGIN''INTEGER'EL;'IF'-(100001'LE'X<X<PTAG)'THEN'GOTO'1;PUNCH(TTAG,X);'GOTO'999;1:'IF'-(200001'LE'X<X<P
BOLD)'THEN'GOTO'2;PUNCH(TBOLD,X);'GOTO'999;2:'IF'-(300001'LE'X<X<PSPEC)'THEN'GOTO'3;PUNCH(TSPEC,X);'GOTO'999;3:'IF'-(400001'LE'X<X<PCONS)'THEN'
'GOTO'4;EL:=TCONS[X];WRITEINT(EL);'GOTO'999;4:WRITEINT(X);999:'END';
37 'PROCEDURE'WRITEINT(X);'INTEGER'X;'BEGIN''INTEGER'QUOT,REM;'IF'-(X<0)'THEN'GOTO'1;REM:=X;REM:=-REM;PUCHAR(65);WRITEINT(REM);'GOTO'999;1
:'IF'-(X=0)'THEN'GOTO'2;PUCHAR(0);'GOTO'999;2:QUOT:=X/'10;REM:=X-10*QUOT;WRITEINT1(QUOT);PUCHAR(REM);999:'END';
38 'PROCEDURE'WRITEINT1(X);'INTEGER'X;'BEGIN''INTEGER'QUOT,REM;'IF'-(X=0)'THEN'GOTO'1;'GOTO'999;1:QUOT:=X/'10;REM:=X-10*QUOT;WRITEINT1(QU
OT);PUCHAR(REM);999:'END';
39 'PROCEDURE'PUNCH(LIST,Y);'INTEGER'Y,LIST;'BEGIN''INTEGER'X,LX;X:=Y;RST:LX:=LIST[X];'BEGIN''IF'-(LX<0)'THEN'GOTO'2;LX:=-LX;PUNCH1(LX);'G
OTO'999;2:PUNCH1(LX);X:=X+1;'GOTO'RST;'END';999:'END';

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40 'PROCEDURE'PUNCH1(X); 'INTEGER'X; 'BEGIN' 'INTEGER'X1,X2,X3;X1:=X/'16384;X3:=X-16384*X1;X2:=X3/'128;X3:=X3-128*X2;PUCHAR(X1);PUCHAR(X2);P
UCHAR(X3); 'END';
41 'PROCEDURE'NEWCARD; 'BEGIN' SPC: 'IF' ~(CPOS<72) 'THEN' 'GOTO'1; CSYM(93); CPOS:=CPOS+1; 'GOTO' SPC; 1: CPOS:=1; WRITE INT(CARD); CSYM(119); CARD:=CARD+
1; CPOS:=1; 'END';
42 'PROCEDURE'TABCARD; 'BEGIN' 'INTEGER' SPCEs, TBS; TBS:=CPOS/'8; SPCEs:=CPOS-8*TBS; SPC: 'IF' ~(SPCEs<8) 'THEN' 'GOTO'1; SPACECARD; SPCEs:=SPCEs+1; 'G
OTO' SPC; 2: 'END';
43 'PROCEDURE'SPACECARD; 'BEGIN' PUCHAR(93); 'END';
44 'PROCEDURE'PUCHAR(X); 'INTEGER'X; 'BEGIN' 'IF' ~(X=119) 'THEN' 'GOTO'1; NEWCARD; 'GOTO'999; 1: 'IF' ~(X=118) 'THEN' 'GOTO'2; TABCARD; 'GOTO'999; 2: 'IF' ~
(X=63) 'THEN' 'GOTO'3; 'GOTO'999; 3: 'IF' ~(CPOS<72) 'THEN' 'GOTO'4; CSYM(X); CPOS:=CPOS+1; 'GOTO'999; 4: CSYM(133); CSYM(119); CPOS:=1; PUCHAR(X); 999: 'END';
45 'COMMENT'4.2.4RESULTADMINISTRATION;
46 'FOCLEAN'LINEDUP;
47 'INTEGER'INDENTATION;
48 'PROCEDURE'G(X); 'INTEGER'X; 'BEGIN'LINEUP;WRITE(X); 'END';
49 'PROCEDURE'LINEUP; 'BEGIN' 'IF' ~LINEDUP 'THEN' 'GOTO'1; 'GOTO'999; 1: 'IF' ~LEGIBLE 'THEN' 'GOTO'2; PUTABS(INDENTATION); LINEDUP:='TRUE'; 'GOTO'999; 2
:999: 'END';
50 'PROCEDURE'PUTABS(N); 'INTEGER'N; 'BEGIN' 'INTEGER'N1; 'IF' ~(N<0) 'THEN' 'GOTO'1; 'GOTO'999; 1: N1:=N; N1:=N1-1; TABCARD; PUTABS(N1); 999: 'END';
51 'PROCEDURE'U; 'BEGIN'INDENTATION:=INDENTATION+1; 'END';
52 'PROCEDURE'L; 'BEGIN'INDENTATION:=INDENTATION-1; 'END';
53 'PROCEDURE'NEWLINE; 'BEGIN' 'IF' ~LEGIBLE 'THEN' 'GOTO'1; 'IF' ~LINEDUP 'THEN' 'GOTO'1; BLANKLINE; 'GOTO'999; 1:999: 'END';
54 'PROCEDURE'BLANKLINE; 'BEGIN'NEWCARD; LINEDUP:='FALSE'; 'END';
55 'PROCEDURE'TABL; 'BEGIN' 'IF' ~LEGIBLE 'THEN' 'GOTO'1; TABCARD; 'GOTO'999; 1:999: 'END';
56 'COMMENT'4.3INPUT;
57 'INTEGER'LINE, CHAR, INPT;
58 'COMMENT'4.3.1READINGCHARACTERS;
59 'FOCLEAN' 'PROCEDURE'LETTER(X); 'INTEGER'X; 'BEGIN' 'IF' ~(CHAR>9&CHAR<36) 'THEN' 'GOTO'1; X:=CHAR; NEXTNONSPACECHAR; LETTER:='TRUE'; 'GOTO'999; 1:L
ETTER:='FALSE'; 999: 'END';
60 'FOCLEAN' 'PROCEDURE'LETGIT(X); 'INTEGER'X; 'BEGIN' 'IF' ~(CHAR<36) 'THEN' 'GOTO'1; X:=CHAR; NEXTNONSPACECHAR; LETGIT:='TRUE'; 'GOTO'999; 1:LETGIT:=
'FALSE'; 999: 'END';
61 'FOCLEAN' 'PROCEDURE'DIGIT(X); 'INTEGER'X; 'BEGIN' 'IF' ~(CHAR<10) 'THEN' 'GOTO'1; X:=CHAR; NEXTNONSPACECHAR; DIGIT:='TRUE'; 'GOTO'999; 1: DIGIT:='FA
LSE'; 999: 'END';
62 'FOCLEAN' 'PROCEDURE'BOLDCHAR(X); 'INTEGER'X; 'BEGIN' 'IF' ~(CHAR=121) 'THEN' 'GOTO'1; X:=120; NEXTCHAR; BOLDCHAR:='TRUE'; 'GOTO'999; 1: X:=CHAR; NEXT
CHAR; BOLDCHAR:='TRUE'; 999: 'END';
63 'FOCLEAN' 'PROCEDURE'ACCENT; 'BEGIN' 'IF' ~(CHAR=120) 'THEN' 'GOTO'1; NEXTCHAR; ACCENT:='TRUE'; 'GOTO'999; 1: ACCENT:='FALSE'; 999: 'END';
64 'FOCLEAN' 'PROCEDURE'SPECCHAR(X); 'INTEGER'X; 'BEGIN' 'IF' ~(CHAR>63) 'THEN' 'GOTO'1; X:=CHAR; NEXTNONSPACECHAR; SPECCHAR:='TRUE'; 'GOTO'999; 1: SPEC
CHAR:='FALSE'; 999: 'END';
65 'PROCEDURE'NEXTCHAR; 'BEGIN'DISPLAYCHARACTER; CHAR:=RESYM; 'END';
66 'PROCEDURE'DISPLAYCHARACTER; 'BEGIN' 'IF' ~(CHAR=119) 'THEN' 'GOTO'1; LINE:=LINE+1; 'BEGIN' 'IF' ~GIVETEXT 'THEN' 'GOTO'2; POSITION(133); OUTINT1(CAR
D); POSITION(48); OUTINT1(LINE); POSITION(52); 'GOTO'999; 2: 'GOTO'999; 'END'; 1: 'IF' ~GIVETEXT 'THEN' 'GOTO'4; PRCHAR(CHAR); 'GOTO'999; 4:999: 'END';
67 'PROCEDURE'NEXTNONSPACECHAR; 'BEGIN'CHR: NEXTCHAR; 'BEGIN' 'IF' ~(CHAR=93) 'THEN' 'GOTO'2; 'GOTO'CHR; 2: 'GOTO'999; 'END'; 999: 'END';
68 'FOCLEAN' 'PROCEDURE'LAYOUTSYMBOL; 'BEGIN' 'IF' ~(CHAR=93) 'THEN' 'GOTO'1; LAYOUTSYMBOL:='TRUE'; 'GOTO'999; 1: 'IF' ~(CHAR=119) 'THEN' 'GOTO'2; LAYOUT
SYMBOL:='TRUE'; 'GOTO'999; 2: 'IF' ~(CHAR=118) 'THEN' 'GOTO'3; LAYOUTSYMBOL:='TRUE'; 'GOTO'999; 3: LAYOUTSYMBOL:='FALSE'; 999: 'END';
69 'COMMENT'4.3.2TAGSYMBOLS;
70 'FOCLEAN' 'PROCEDURE'READTAG(X); 'INTEGER'X; 'BEGIN' 'INTEGER'T, T1, T2, T3, T4; 'IF' ~LETTER(T1) 'THEN' 'GOTO'1; X:=PTAG; NXT: 'BEGIN' 'IF' ~LETGIT(T2)
' THEN' 'GOTO'2; 'BEGIN' 'IF' ~LETGIT(T3) 'THEN' 'GOTO'3; 'BEGIN' 'IF' ~LETGIT(T4) 'THEN' 'GOTO'4; T:=(T1*128+T2)*128+T3; STACKTAG(T); T1:=T4; 'GOTO'NXT; 4: T:=(T1
*128+T2)*128+T3; 'GOTO'LST; 'END'; 3: T:=(T1*128+T2)*128+63; 'GOTO'LST; 'END'; 2: T:=(T1*128+63)*128+63; LST: T:=-T; STACKTAG(T); READTAG:='TRUE'; 'GOTO'999;
'END'; 1: READTAG:='FALSE'; 999: 'END';
71 'PROCEDURE'STACKTAG(X); 'INTEGER'X; 'BEGIN' 'IF' ~(PTAG<11000) 'THEN' 'GOTO'1; TTAG[PTAG]:=X; PTAG:=PTAG+1; 'GOTO'999; 1: ERROR(TAGFULL, 0); POSTMOR
TEM; EXIT; 999: 'END';
72 'PROCEDURE'ENTERTAG(X); 'INTEGER'X; 'BEGIN' 'INTEGER'Y, X1, Y1, WX, WY; Y:=FIRSTTAG; NXY: X1:=X; Y1:=Y; NXW: WX:=TTAG[X1]; WY:=TTAG[Y1]; 'BEGIN' 'IF' ~(W
X=WY) 'THEN' 'GOTO'2; 'BEGIN' 'IF' ~(WX<0) 'THEN' 'GOTO'3; 'BEGIN' 'IF' ~(Y=X) 'THEN' 'GOTO'4; RESERVEADMINSNPACE; 'GOTO'999; 4: PTAG:=X; X:=Y; 'GOTO'999; 'END'; 3: X
1:=X1+1; Y1:=Y1+1; 'GOTO'NXW; 'END'; 2: 'IF' ~(ABS(WX)<ABS(WY)) 'THEN' 'GOTO'7; WY:=TTAG[Y-2]; 'BEGIN' 'IF' ~(WY=0) 'THEN' 'GOTO'8; TTAG[Y-2]:=X; RESERVEADMINSP
ACE; 'GOTO'999; 8: Y:=WY; 'GOTO'NXY; 'END'; 7: WY:=TTAG[Y-1]; 'BEGIN' 'IF' ~(WY=0) 'THEN' 'GOTO'11; TTAG[Y-1]:=X; RESERVEADMINSNPACE; 'GOTO'999; 11: Y:=WY; 'GOTO'N
XY; 'END'; 'END'; 999: 'END';
73 'PROCEDURE'RESERVEADMINSNPACE; 'BEGIN'PTAG:=PTAG+8; 'IF' ~(PTAG<11000) 'THEN' 'GOTO'1; TTAG[PTAG-8]:=999; TTAG[PTAG-7]:=0; TTAG[PTAG-6]:=0; TTAG[
PTAG-5]:=0; TTAG[PTAG-4]:=0; TTAG[PTAG-3]:=0; TTAG[PTAG-2]:=0; TTAG[PTAG-1]:=0; 'GOTO'999; 1: ERROR(TAGFULL, 0); POSTMORTEM; EXIT; 999: 'END';
74 'COMMENT'4.3.3BOLDSYMBOLS;
75 'FOCLEAN' 'PROCEDURE'READBOLD(X); 'INTEGER'X; 'BEGIN' 'INTEGER'B1, B2, B3, D; 'IF' ~ACCENT 'THEN' 'GOTO'1; X:=PBOLD; 'BEGIN' 'IF' ~ACCENT 'THEN' 'GOTO'2;
D:=(63*128+63)*128+63; D:=-D; STACKBOLD(D); READBOLD:='TRUE'; 'GOTO'999; 2: 'IF' ~BOLDCHAR(B1) 'THEN' 'GOTO'3; RST: 'BEGIN' 'IF' ~ACCENT 'THEN' 'GOTO'4; D:=(B1*
128+63)*128+63; D:=-D; STACKBOLD(D); READBOLD:='TRUE'; 'GOTO'999; 4: 'IF' ~BOLDCHAR(B2) 'THEN' 'GOTO'5; 'BEGIN' 'IF' ~ACCENT 'THEN' 'GOTO'6; D:=(B1*128+B2)*128
+63; D:=-D; STACKBOLD(D); READBOLD:='TRUE'; 'GOTO'999; 6: 'IF' ~BOLDCHAR(B3) 'THEN' 'GOTO'7; D:=(B1*128+B2)*128+63; 'BEGIN' 'IF' ~ACCENT 'THEN' 'GOTO'8; D:=-D; S

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TACKBOLD(D);READBOLD:='TRUE';GOTO'999;8:'IF'~BOLDCHAR(B1)THEN'GOTO'9;STACKBOLD(D);GOTO'RST;9:'GOTO'0;'END';7:'GOTO'0;'END';5:'GOTO'0;'END';3
:'GOTO'0;'END';1:0:READBOLD:='FALSE';999:'END';
76 'PRCCEDURE'STACKBOLD(X);'INTEGER'X;'BEGIN''IF'~(PBOLD<200600)THEN'GOTO'1;TBOLD[PBOLD]:=X;PBOLD:=PBOLD+1;'GOTO'999;1:ERROR(BOLDFULL,0);
PCSTMORTEM;EXIT;999:'END';
77 'PRCCEDURE'ENTERBOLD(X);'INTEGER'X;'BEGIN''INTEGER'Y,X1,Y1,WX,WY;Y:=200001;NXY:X1:=X;Y1:=Y;NXW:WX:=TBOLD[X1];WY:=TBOLD[Y1];'BEGIN''IF'~(
WX=WY)THEN'GOTO'2;'BEGIN''IF'~(WX<0)THEN'GOTO'3;'BEGIN''IF'~(Y=X)THEN'GOTO'4;'GOTO'999;4:PBOLD:=X;X:=Y;'GOTO'999;'END';3:X1:=X1+1;Y1:=Y1+1
;'GOTO'XW;'END';2:SKP;'IF'~(WY<0)THEN'GOTO'7;Y:=Y1;Y:=Y+1;'GOTO'NXY;7:Y1:=Y1+1;WY:=TBOLD[Y1];'GOTO'SKP;'END';999:'END';
78 'COMMENT'4.3.4SPECIALSYMBOLS;
79 'BOCLEAN'PROCEDURE'READSPEC(X);'INTEGER'X;'BEGIN''INTEGER'S,S1;'IF'~SPECCHAR(S1)THEN'GOTO'1;X:=PSPEC;S1:=(S1*128+63)*128+63;S:=-S;STAC
KSPEC(S);READSPEC:='TRUE';GOTO'999;1:READSPEC:='FALSE';999:'END';
80 'PRCCEDURE'STACKSPEC(X);'INTEGER'X;'BEGIN''IF'~(PSPEC<300100)THEN'GOTO'1;TSPEC[PSPEC]:=X;PSPEC:=PSPEC+1;'GOTO'999;1:ERROR(SPECFULL,0);
PCSTMORTEM;EXIT;999:'END';
81 'PRCCEDURE'ENTERSPEC(X);'INTEGER'X;'BEGIN''INTEGER'Y,WX,WY;Y:=300001;WX:=TSPEC[X];NXY:WY:=TSPEC[Y];'BEGIN''IF'~(WX=WY)THEN'GOTO'2;'BEG
IN''IF'~(V=X)THEN'GOTO'3;'GOTO'999;3:PSPEC:=X;X:=Y;'GOTO'999;'END';2:Y:=Y+1;'GOTO'NXY;'END';999:'END';
82 'COMMENT'4.3.5CONSTANTS;
83 'BOCLEAN'PROCEDURE'READCONS(X);'INTEGER'X;'BEGIN''INTEGER'D,S;'IF'~DIGIT(D)THEN'GOTO'1;S:=D;RST:'BEGIN''IF'~DIGIT(D)THEN'GOTO'2;S:=
S*10+D;'GOTO'RST;2:X:=PCONS;STACKCONS(S);READCONS:='TRUE';GOTO'999;'END';1:READCONS:='FALSE';999:'END';
84 'PRCCEDURE'STACKCONS(X);'INTEGER'X;'BEGIN''IF'~(PCONS<400300)THEN'GOTO'1;TCONS[PCONS]:=X;PCONS:=PCONS+1;'GOTO'999;1:ERROR(CONSFULL,0);
PCSTMORTEM;EXIT;999:'END';
85 'PRCCEDURE'ENTERCONS(X);'INTEGER'X;'BEGIN''INTEGER'Y,S,T;S:=TCONS[X];Y:=400001;NXT:T:=TCONS[Y];'BEGIN''IF'~(S=T)THEN'GOTO'2;'BEGIN''IF
'~(X=Y)THEN'GOTO'3;'GOTO'999;3:PCONS:=X;X:=Y;'GOTO'999;'END';2:Y:=Y+1;'GOTO'NXT;'END';999:'END';
86 'COMMENT'4.3.6INPUTADMINISTRATCN;
87 'PRCCEDURE'INITIALIZEFORREADING;'BEGIN'GIVETEXT:='FALSE';PCONS:=400001;PTAG:=100001;PBOLD:=200001;PSPEC:=300001;LINE:=0;CHAR:=RESYM;RESE
RVEADMINSPACE;FIRSTTAG:=PTAG;'END';
88 'BOCLEAN'PROCEDURE'ISTAG(X);'INTEGER'X;'BEGIN''IF'~(100001'LE'INPT<INPT<PTAG)THEN'GOTO'1;X:=INPT;NEXTSYMBOL;ISTAG:='TRUE';GOTO'999;1
:ISTAG:='FALSE';999:'END';
89 'BOCLEAN'PROCEDURE'ISBOLD(X);'INTEGER'X;'BEGIN''IF'~(200001'LE'INPT<INPT<PBOLD)THEN'GOTO'1;X:=INPT;NEXTSYMBOL;ISBOLD:='TRUE';GOTO'99
9;1:ISBOLD:='FALSE';999:'END';
90 'BOCLEAN'PROCEDURE'ISCONS(X);'INTEGER'X;'BEGIN''IF'~(400001'LE'INPT<INPT<PCONS)THEN'GOTO'1;X:=INPT;NEXTSYMBOL;ISCONS:='TRUE';GOTO'99
9;1:ISCONS:='FALSE';999:'END';
91 'BOCLEAN'PROCEDURE'R(X);'INTEGER'X;'BEGIN''IF'~AHEAD(X)THEN'GOTO'1;NEXTSYMBOL;R:='TRUE';GOTO'999;1:R:='FALSE';999:'END';
92 'BOCLEAN'PROCEDURE'AHEAD(X);'INTEGER'X;'BEGIN''IF'~(INPT=X)THEN'GOTO'1;AHEAD:='TRUE';GOTO'999;1:AHEAD:='FALSE';999:'END';
93 'PRCCEDURE'NEXTSYMBOL;'BEGIN'READ(INPT);'END';
94 'PRCCEDURE'READ(X);'INTEGER'X;'BEGIN'SKP;'IF'~LAYOUTSYMBOLTHEN'GOTO'1;NEXTCHAR;'GOTO'SKP;1:'IF'~READTAG(X)THEN'GOTO'2;ENTERTAG(X);'G
OTO'999;2:'IF'~READBOLD(X)THEN'GOTO'3;ENTERBOLD(X);'GOTO'999;3:'IF'~READSPEC(X)THEN'GOTO'4;ENTERSPEC(X);'GOTO'999;4:'IF'~READCONS(X)THEN'G
OTO'5;ENTERCONS(X);'GOTO'999;5:ERROR(WRONGINPUTCODE,CHAR);NEXTCHAR;READ(X);999:'END';
95 'COMMENT'4.4.COMPIILERCOMPILERCALGOL60;
96 'BOCLEAN'GIVERSTO,ADVANCED;
97 'INTEGER'HANDLE,BOUNDX,SPECX,TONEXT,TOLOST,TOEND,CURLAB,NEULAB;
98 'COMMENT'4.4.1AUXILIARYACTIONS;
99 'PRCCEDURE'ADDPPLACE(TAG);'INTEGER'TAG;'BEGIN''INTEGER'H,K,Q;H:=TTAG[TAG-6];'BEGIN''IF'~(H=0)THEN'GOTO'2;REPLACE(TAG,H);'GOTO'999;2:H:=
LTEXT[H];K:=Q-'/'8192*8192;'IF'~(K<LINE)THEN'GOTO'3;REPLACE(TAG,H);'GOTO'999;3:'GOTO'999;'END';999:'END';
100 'PRCCEDURE'REPLACE(TAG,H);'INTEGER'H,TAG;'BEGIN''TAG[TAG-6]:=PTEXT;'BEGIN''IF'~(H=0)THEN'GOTO'2;ENTERTEXT(LINE);'GOTO'999;2:H:=H-80000
1;H:=H+1;H:=8192*H+LINE;ENTERTEXT(H);'GOTO'999;'END';999:'END';
101 'PRCCEDURE'FORGETLOCALS;'BEGIN''INTEGER'LOC;RST:'IF'~(PLOC=500001)THEN'GOTO'1;'GOTO'999;1:PLOC:=PLOC-1;LOC:=LLOC[PLOC];WARN(LCC);TTAG[
LCC-5]:=0;TTAG[LOC-4]:=0;TTAG[LOC-3]:=0;'GOTO'RST;999:'END';
102 'PRCCEDURE'WARN(X);'INTEGER'X;'BEGIN''INTEGER'STATE;STATE:=TTAG[X-5];'BEGIN''IF'~(STATE=BLANK)THEN'GOTO'2;TRACE(X);'GOTO'999;2:'IF'~(S
TATE=0)THEN'GOTO'3;TRACE(X);'GOTO'999;3:SIGNAL(X);'GOTO'999;'END';999:'END';
103 'PRCCEDURE'ENTERLOCAL(HEAD);'INTEGER'HEAD;'BEGIN''IF'~(PLOC<500200)THEN'GOTO'1;LLOC[PLOC]:=HEAD;PLOC:=PLOC+1;'BEGIN''IF'~(PLOC<XLOC)'T
HEN'GOTO'2;'GOTO'999;2:XLOC:=PLOC;'GOTO'999;'END';1:ERROR(LOCFULL,0);POSTMORTEM;EXIT;999:'END';
104 'PRCCEDURE'ENTERGLOBAL(HEAD);'INTEGER'HEAD;'BEGIN''IF'~(PLOB<601000)THEN'GOTO'1;LGLOBAL[PLOB]:=HEAD;PLOB:=PLOB+1;'GOTO'999;1:ERROR(G
LOBFULL,0);POSTMORTEM;EXIT;999:'END';
105 'PRCCEDURE'ENTERMACRO(HEAD);'INTEGER'HEAD;'BEGIN''IF'~(PMACR<701100)THEN'GOTO'1;LMACR[PMACR]:=HEAD;PMACR:=PMACR+1;'GOTO'999;1:ERROR(MA
CRFULL,0);PCSTMORTEM;EXIT;999:'END';
106 'PRCCEDURE'ENTERTEXT(X);'INTEGER'X;'BEGIN''IF'~(PTEXT<804000)THEN'GOTO'1;LTEXT[PTEXT]:=X;PTEXT:=PTEXT+1;'GOTO'999;1:ERROR(TEXTFULL,0);
PCSTMORTEM;EXIT;999:'END';
107 'COMMENT'4.4.2TREATMENTOFTYPES;
108 'PRCCEDURE'DEFINEACTION(HEAD);'INTEGER'HEAD;'BEGIN'REDEFINE(HEAD,DEFINED,GLOBAL,ACTION);'END';
109 'PRCCEDURE'DEFINEPREDICATE(HEAD);'INTEGER'HEAD;'BEGIN'REDEFINE(HEAD,DEFINED,GLOBAL,PREDICATE);'END';
110 'PRCCEDURE'DEFINEAFFIX(HEAD);'INTEGER'HEAD;'BEGIN'REDEFINE(HEAD,DEFINED,LOCAL,POINTER);'END';
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111 'PROCEDURE'DEFINELABEL(LAB);'INTEGER'LAB;'BEGIN''IF'-(100001'LE'LAB^LAB<PTAG)'THEN''GOTO'1;REDEFINE(LAB,DEFINED,LOCAL,LABEL);'GOTO'999;1
:'IF'-(400001'LE'LAB^LAB<PCONS)'THEN''GOTO'2;'GOTO'999;2;999:'END';
112 'PROCEDURE'APPLYACTION(HEAD);'INTEGER'HEAD;'BEGIN'REDEFINE(HEAD,APPLIED,GLOBAL,ACTION);'END';
113 'PROCEDURE'APPLYPREDICATE(HEAD);'INTEGER'HEAD;'BEGIN'REDEFINE(HEAD,APPLIED,GLOBAL,PREDICATE);'END';
114 'PROCEDURE'APPLYLABEL(LAB);'INTEGER'LAB;'BEGIN''IF'-(100001'LE'LAB^LAB<PTAG)'THEN''GOTO'1;REDEFINE(LAB,APPLIED,LOCAL,LABEL);'GOTO'999;1;
'IF'-(400001'LE'LAB^LAB<PCONS)'THEN''GOTO'2;'GOTO'999;2;999:'END';
115 'PROCEDURE'APPLY(X);'INTEGER'X;'BEGIN'NEWPLACE(X);'IF'-TRYSTATE(X,APPLIED)'THEN''GOTO'1;'GOTO'999;1;999:'END';
116 'PROCEDURE'REDEFINE(X,NSTATE,NSORT,NTYPE);'INTEGER'NTYPE,NSORT,NSTATE,X;'BEGIN''INTEGER'OLDPOS;'IF'-REDEFINE1(X,NSTATE,NSORT,NTYPE)'THEN
'GOTO'1;NEWPLACE(X);'GOTO'999;1;CLDPOS:=POS;NLCR;OUT(XIMPOSSIBLE);INFORM(X);OUT(NSTATE);OUT(NSORT);OUT(NTYPE);POSITION(CLDPOS);999:'END';
117 'BOCLEAN''PROCEDURE'REDEFINE1(X,NSTATE,NSORT,NTYPE);'INTEGER'NTYPE,NSORT,NSTATE,X;'BEGIN''IF'-TRYTYPE(X,NTYPE)'THEN''GOTO'1;'IF'-TRYSORT
(X,NSORT)'THEN''GOTO'0;'IF'-TRYSTATE(X,NSTATE)'THEN''GOTO'0;REDEFINE1:='TRUE';'GOTO'999;1;0:REDEFINE1:='FALSE';999:'END';
118 'BOCLEAN''PROCEDURE'TRYSTATE(X,S);'INTEGER'S,X;'BEGIN''INTEGER'Q;Q:=TTAG[X-5];'IF'-(Q=0)'THEN''GOTO'1;TTAG[X-5]:=S;TRYSTATE:='TRUE';'GOTO
999;1;'IF'-(Q=DEFINED)'THEN''GOTO'2;'IF'-(S=APPLIED)'THEN''GOTO'2;TTAG[X-5]:=BLANK;TRYSTATE:='TRUE';'GOTO'999;2;'IF'-(Q=APPLIED)'THEN''GOTO'3;
'BEGIN''IF'-(S=APPLIED)'THEN''GOTO'4;TRYSTATE:='TRUE';'GOTO'999;4;'IF'-(S=DEFINED)'THEN''GOTO'5;TTAG[X-5]:=BLANK;TRYSTATE:='TRUE';'GOTO'999;5;'G
OTO'0;'END';3:'IF'-(Q=BLANK)'THEN''GOTO'7;'IF'-(S=APPLIED)'THEN''GOTO'7;TRYSTATE:='TRUE';'GOTO'999;7;0:TRYSTATE:='FALSE';999:'END';
119 'BOCLEAN''PROCEDURE'TRYSORT(X,P);'INTEGER'P,X;'BEGIN''INTEGER'Q;Q:=TTAG[X-4];'IF'-(Q=0)'THEN''GOTO'1;TTAG[X-4]:=P;'BEGIN''IF'-(P=LOCAL)'
THEN''GOTO'2;ENTERLOCAL(X);TRYSORT:='TRUE';'GOTO'999;2;'IF'-(P=GLOBAL)'THEN''GOTO'3;ENTERGLOBAL(X);TRYSORT:='TRUE';'GOTO'999;3;'IF'-(P=MACRO)'TH
EN''GOTO'4;ENTERGLOBAL(X);TRYSORT:='TRUE';'GOTO'999;4;'GOTO'0;'END';1:'IF'-(P=Q)'THEN''GOTO'6;TRYSORT:='TRUE';'GOTO'999;6;0:TRYSORT:='FALSE';999
:'END';
120 'BOCLEAN''PROCEDURE'TRYTYPE(X,P);'INTEGER'P,X;'BEGIN''INTEGER'Q;Q:=TTAG[X-3];'IF'-(Q=0)'THEN''GOTO'1;TTAG[X-3]:=P;TRYTYPE:='TRUE';'GOTO'
999;1;'IF'-(P=Q)'THEN''GOTO'2;TRYTYPE:='TRUE';'GOTO'999;2;TRYTYPE:='FALSE';999:'END';
121 'PROCEDURE'NEWPLACE(X);'INTEGER'X;'BEGIN''INTEGER'SORT;SORT:=TTAG[X-4];'IF'-(SORT=LOCAL)'THEN''GOTO'1;'GOTO'999;1;ADDPLACE(X);999:'END';
122 'PROCEDURE'NEWNPARS(X,P);'INTEGER'P,X;'BEGIN''INTEGER'Q;Q:=TTAG[X-8];'IF'-(Q=P)'THEN''GOTO'1;'GOTO'999;1;'IF'-(Q=999)'THEN''GOTO'2;TTAG[
X-8]:=P;'GOTO'999;2;ERROR(WRONGNUMBEROFFPARAMETERS,X);999:'END';
123 'BOCLEAN''PROCEDURE'WASACTION(HEAD);'INTEGER'HEAD;'BEGIN''INTEGER'T;T:=TTAG[HEAD-3];'IF'-(T=ACTION)'THEN''GOTO'1;WASACTION:='TRUE';'GOTO
999;1;WASACTION:='FALSE';999:'END';
124 'BOCLEAN''PROCEDURE'WASAFFIX(HEAD);'INTEGER'HEAD;'BEGIN''INTEGER'T;T:=TTAG[HEAD-3];'IF'-(T=POINTER)'THEN''GOTO'1;WASAFFIX:='TRUE';'GOTO'
999;1;'IF'-(T=FLAG)'THEN''GOTO'2;WASAFFIX:='TRUE';'GOTO'999;2;'IF'-(T=LIST)'THEN''GOTO'3;WASAFFIX:='TRUE';'GOTO'999;3;WASAFFIX:='FALSE';999:'END
';
125 'BOCLEAN''PROCEDURE'WASMACRC(HEAD);'INTEGER'HEAD;'BEGIN''INTEGER'S;S:=TTAG[HEAD-4];'IF'-(S=MACRO)'THEN''GOTO'1;WASMACRC:='TRUE';'GOTO'99
9;1;WASMACRC:='FALSE';999:'END';
126 'BOCLEAN''PROCEDURE'WASPARAMLESSMACRO(HEAD);'INTEGER'HEAD;'BEGIN''IF'-(100001'LE'HEAD^HEAD<PTAG)'THEN''GOTO'1;'IF'-'WASMACRO(HEAD)'THEN''
GOTO'1;'IF'-'WASAFFIX(HEAD)'THEN''GOTO'0;WASPARAMLESSMACRO:='TRUE';'GOTO'999;1;0:WASPARAMLESSMACRO:='FALSE';999:'END';
127 'COMMENT'4,4,3SPECIFICATIONS;
128 'BOCLEAN''PROCEDURE'SPECIFICATION;'BEGIN''IF'-'EXTERNALSPECIFICATION'THEN''GOTO'1;SPECIFICATION:='TRUE';'GOTO'999;1;'IF'-'INTERNALSPECIFIC
ATION'THEN''GOTO'2;SPECIFICATION:='TRUE';'GOTO'999;2;'IF'-'MACROSPECIFICATION'THEN''GOTO'3;SPECIFICATION:='TRUE';'GOTO'999;3;SPECIFICATION:='FALS
E';999:'END';
129 'BOCLEAN''PROCEDURE'EXTERNALSPECIFICATION;'BEGIN''INTEGER'TYPE;'IF'-'R(EXTERNAL)'THEN''GOTO'1;'IF'-'ISTYPE(TYPE)'THEN''GOTO'0;TREATSPECLIS
T(DEFINED,GLOBAL,TYPE);EXTERNALSPECIFICATION:='TRUE';'GOTO'999;1;0:EXTERNALSPECIFICATION:='FALSE';999:'END';
130 'BOCLEAN''PROCEDURE'ISTYPE(X);'INTEGER'X;'BEGIN'X:=INPT;'IF'-'R(ACTION)'THEN''GOTO'1;ISTYPE:='TRUE';'GOTO'999;1;'IF'-'R(PREDICATE)'THEN''G
OTO'2;ISTYPE:='TRUE';'GOTO'999;2;'IF'-'R(POINTER)'THEN''GOTO'3;ISTYPE:='TRUE';'GOTO'999;3;'IF'-'R(FLAG)'THEN''GOTO'4;ISTYPE:='TRUE';'GOTO'999;4;IS
TYPE:='FALSE';999:'END';
131 'PROCEDURE'TREATSPECLIST(STATE, SORT, TYPE);'INTEGER'TYPE, SORT, STATE;'BEGIN''INTEGER'HEAD;NXT;'IF'-'ISTAG(HEAD)'THEN''GOTO'1;REDEFINE(HEAD,
STATE, SORT, TYPE);'BEGIN''IF'-'R(COMMA)'THEN''GOTO'2;'GOTO'NXT;2:SHIFT2LINES;'IF'-'R(POINT)'THEN''GOTO'3;'GOTO'999;3;'GOTO'0;'END';1;0:999:'END';
132 'BOCLEAN''PROCEDURE'INTERNALSPECIFICATION;'BEGIN''INTEGER'TYPE;'IF'-'ISTYPEOFLOCAL(TYPE)'THEN''GOTO'1;TREATSPECLIST(APPLIED,GLOBAL,TYPE);
INTERNALSPECIFICATION:='TRUE';'GOTO'999;1;INTERNALSPECIFICATION:='FALSE';999:'END';
133 'BOCLEAN''PROCEDURE'ISTYPEOFLOCAL(X);'INTEGER'X;'BEGIN'X:=INPT;'IF'-'R(ACTION)'THEN''GOTO'1;ISTYPEOFLOCAL:='TRUE';'GOTO'999;1;'IF'-'R(PRED
ICATE)'THEN''GOTO'2;ISTYPEOFLOCAL:='TRUE';'GOTO'999;2;ISTYPEOFLOCAL:='FALSE';999:'END';
134 'BOCLEAN''PROCEDURE'MACROSPECIFICATION;'BEGIN''INTEGER'TYPE,HEAD;'IF'-'R(MACRO)'THEN''GOTO'1;'IF'-'ISTYPE(TYPE)'THEN''GOTO'0;NXT;'IF'-'ISTA
G(HEAD)'THEN''GOTO'0;REDEFINE(HEAD,DEFINED,MACRO,TYPE);READMACRO(HEAD);'BEGIN''IF'-'R(COMMA)'THEN''GOTO'2;'GOTO'NXT;2:SHIFT2LINES;'IF'-'R(POINT)'T
HEN''GOTO'3;MACROSPECIFICATION:='TRUE';'GOTO'999;3;'GOTO'0;'END';1;0:MACROSPECIFICATION:='FALSE';999:'END';
135 'PROCEDURE'READMACRO(HEAD);'INTEGER'HEAD;'BEGIN''INTEGER'X;TTAG[HEAD-7]:=PMACR;NXT;'BEGIN''IF'-'AHEAD(POINT)'THEN''GOTO'2;ENTERMACRO(POIN
T);'GOTO'999;2;'IF'-'AHEAD(COMMA)'THEN''GOTO'3;ENTERMACRO(POINT);'GOTO'999;3;'IF'-'ISTAG(X)'THEN''GOTO'4;'IF'-'TRYSTATE(X,APPLIED)'THEN''GOTO'0;ENT
ERMACRO(X);'GOTO'NXT;4:ENTERMACRO(INPT);NEXTSYMBOL;'GOTO'NXT;'END';0:999:'END';
136 'COMMENT'4,4,4DECLARATIONS;
137 'BOCLEAN''PROCEDURE'DECLARATION;'BEGIN''IF'-'POINTERDECLARATION'THEN''GOTO'1;DECLARATION:='TRUE';'GOTO'999;1;'IF'-'FLAGDECLARATION'THEN''G
OTO'2;DECLARATION:='TRUE';'GOTO'999;2;'IF'-'LISTDECLARATION'THEN''GOTO'3;DECLARATION:='TRUE';'GOTO'999;3;DECLARATION:='FALSE';999:'END';
138 'BOCLEAN''PROCEDURE'POINTERDECLARATION;'BEGIN''IF'-'R(POINTER)'THEN''GOTO'1;TREATDECLLIST(QINTEGER,POINTER);POINTERDECLARATION:='TRUE';'G
OTO'999;1;PCINTERDECLARATION:='FALSE';999:'END';
139 'PROCEDURE'TREATDECLLIST(ALGOLTYPE, TYPE);'INTEGER'TYPE,ALGOLTYPE;'BEGIN''INTEGER'HEAD;BLANKLINE;G(ALGOLTYPE);NXT;'IF'-'ISTAG(HEAD)'THEN''

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GOTC'1;REDEFINE(HEAD,DEFINED,GLOBAL,TYPE);G(HEAD);'BEGIN''IF'-R(COMMA)'THEN''GOTO'2;G(COMMA);'GOTO'NXT;2;SHIFT2LINES;'IF'-R(POINT)'THEN''GOTO'3;
G(SEMICOLON);'GOTO'999;3;'GOTO'0;'END';1;0:999;'END';
140 'EOCLEAN''PROCEDURE'FLAGDECLARATION;'BEGIN''IF'-R(FLAG)'THEN''GOTO'1;TREATDECLLIST(QBOOLEAN,FLAG);FLAGDECLARATION:='TRUE';'GOTO'999;1;FL
AGDECLARATION:='FALSE';999;'END';
141 'EOCLEAN''PROCEDURE'LISTDECLARATION;'BEGIN''INTEGER'HEAD;'IF'-R(LIST)'THEN''GOTO'1;BLANKLINE;G(QINTEGER);G(QARRAY);U;NXT;'IF'-ISTAG(HEAD
)'THEN''GOTO'0;G(HEAD);REDEFINE(HEAD,DEFINED,GLOBAL,LIST);'IF'-R(SUB)'THEN''GOTO'0;G(SUB);CONSTANTTEXT;'IF'-R(COLON)'THEN''GOTO'0;G(COLON);CONST
ANTTEXT;'IF'-R(BUS)'THEN''GOTO'0;G(BUS);'BEGIN''IF'-R(COMMA)'THEN''GOTO'2;G(COMMA);NEWLINE;'GOTO'NXT;2;SHIFT2LINES;'IF'-R(POINT)'THEN''GOTO'3;G(
SEMICOLON);LISTDECLARATION:='TRUE';'GOTO'999;3;'GOTO'0;'END';1;0:LISTDECLARATION:='FALSE';999;'END';
142 'PROCEDURE'CONSTANTTEXT;'BEGIN''INTEGER'SYMB;NXT;'IF'-ISTAG(SYMB)'THEN''GOTO'1;APPLY(SYMB);'BEGIN''IF'-WASPARAMLESSMACRO(SYMB)'THEN''GOT
O'2;PUTPARAMLESSMACRO(SYMB);'GOTO'NXT;2;G(SYMB);'GOTO'NXT;'END';1;'IF'-ISCONS(SYMB)'THEN''GOTO'4;G(SYMB);'GOTO'NXT;4;'IF'-R(PLUS)'THEN''GOTO'5;G
(PLUS);'GOTO'NXT;5;'IF'-R(MINUS)'THEN''GOTO'6;G(MINUS);'GOTO'NXT;6;'END';
143 'COMMENT'4,4,5AFFIXEXPRESSIONS;
144 'PROCEDURE'PUTAFFIXEXPRESSION(HEAD);'INTEGER'HEAD;'BEGIN''IF'-TRYSORT(HEAD,GLOBAL)'THEN''GOTO'1;PUTDIRECT(HEAD);'GOTO'999;1;'IF'-TRYSORT
(HEAD,MACRO)'THEN''GOTO'2;'BEGIN''IF'-WASACTION(HEAD)'THEN''GOTO'3;PUTMACRO(HEAD);'GOTO'999;3;G(OPEN);PUTMACRO(HEAD);G(CLOSE);'GOTO'999;'END';2;
999;'END';
145 'PROCEDURE'PUTDIRECT(HEAD);'INTEGER'HEAD;'BEGIN''INTEGER'NPARS;NPARS:=0;G(HEAD);AFFIXPACKOPTION(NPARS);NEWNPARS(HEAD,NPARS);'END';
146 'PROCEDURE'AFFIXPACKOPTION(NPARS);'INTEGER'NPARS;'BEGIN''IF'-R(PLUS)'THEN''GOTO'1;G(OPEN);TREATAFFIX;RST;'BEGIN'NPARS:=NPARS+1;'IF'-R(PL
US)'THEN''GOTO'2;G(COMMA);TREATAFFIX;'GOTO'RST;2;G(CLOSE);'GOTO'999;'END';1:999;'END';
147 'PROCEDURE'TREATAFFIX;'BEGIN''INTEGER'AFFX;'IF'-ISTAG(AFFX)'THEN''GOTO'1;PUTSINGLE(AFFX);APPLY(AFFX);'GOTO'999;1;'IF'-ISBOLD(AFFX)'THEN''
GOTO'2;G(STRINGQUOTE);G(AFFX);G(STRINGQUOTE);'GOTO'999;2;'IF'-ISCONS(AFFX)'THEN''GOTO'3;G(AFFX);'GOTO'999;3;ERROR(WRONGAFFX,AFFX);999;'END';
148 'PROCEDURE'PUTSINGLE(X);'INTEGER'X;'BEGIN''IF'-WASAFFIX(X)'THEN''GOTO'1;'BEGIN''IF'-WASMACRO(X)'THEN''GOTO'2;PUTPARAMLESSMACRO(X);'GOTO'
999;2;G(X);'GOTO'999;'END';1;REDEFINE(X,APPLIED,GLOBAL,POINTER);G(X);999;'END';
149 'PROCEDURE'PUTPARAMLESSMACRO(Y);'INTEGER'Y;'BEGIN''INTEGER'X,M;X:=TTAG[Y-7];NEWNPARS(Y,0);NXT:X=X+1;M:=LMACR[X];'BEGIN''IF'-M=POINT)'T
HEN''GOTO'2;'GOTO'999;2;G(M);'GOTO'NXT;'END';999;'END';
150 'PROCEDURE'PUTMACRO(X);'INTEGER'X;'BEGIN''INTEGER'P1,P2,P3,P4,P5,Q,M,NPARS;NPARS:=0;Q:=TTAG[X-7];GETPARAMETERS(NPARS,P1,P2,P3,P4,P5);NEW
NPARS(X,NPARS);NXT:Q=Q+1;M:=LMACR[Q];'BEGIN''IF'-M=POINT)'THEN''GOTO'2;'GOTO'999;2;'IF'-M=ONE)'THEN''GOTO'3;PUTPAR(P1,1);'GOTO'NXT;3;'IF'-M=
TWO)'THEN''GOTO'4;PUTPAR(P2,2);'GOTO'NXT;4;'IF'-M=THREE)'THEN''GOTO'5;PUTPAR(P3,3);'GOTO'NXT;5;'IF'-M=FOUR)'THEN''GOTO'6;PUTPAR(P4,4);'GOTO'NX
T;6;'IF'-M=FIVE)'THEN''GOTO'7;PUTPAR(P5,5);'GOTO'NXT;7;'IF'-WASPARAMLESSMACRO(M)'THEN''GOTO'8;PUTPARAMLESSMACRO(M);'GOTO'NXT;8;G(M);'GOTO'NXT;'
END';999;'END';
151 'PROCEDURE'GETPARAMETERS(P,P1,P2,P3,P4,P5);'INTEGER'P5,P4,P3,P2,P1,P;'BEGIN'GETPAR(P,P1);GETPAR(P,P2);GETPAR(P,P3);GETPAR(P,P4);GETPAR(P
,P5);'END';
152 'PROCEDURE'GETPAR(P,X);'INTEGER'X,P;'BEGIN''IF'-R(PLUS)'THEN''GOTO'1;P:=P+1;'BEGIN''IF'-ISTAG(X)'THEN''GOTO'2;APPLY(X);'GOTO'999;2;'IF'-
ISCONS(X)'THEN''GOTO'3;'GOTO'999;3;'GOTO'0;'END';1;X:=XPARAMETERERROR;'GOTO'999;0:999;'END';
153 'PROCEDURE'PUTPAR(X,Y);'INTEGER'Y,X;'BEGIN''IF'-R(40001'LE'X<X<PCONS)'THEN''GOTO'1;G(X);'GOTO'999;1;'IF'-X=XPARAMETERERROR)'THEN''GOTO'
2;ERROR(X,Y);'GOTO'999;2;'IF'-R(20001'LE'X<X<PBOLD)'THEN''GOTO'3;G(STRINGQUOTE);G(X);G(STRINGQUOTE);'GOTO'999;3;PUTSINGLE(X);999;'END';
154 'COMMENT'4,4,6BUILDINGSTONESOFALB;
155 'PROCEDURE'PUTCONNECTAND;'BEGIN''IF'-GIVERSTO'THEN''GOTO'1;PUTJUMP(CURLAB);TONEXT:=TONEXT+1;'GOTO'999;1;'IF'-ADVANCED'THEN''GOTO'2;ERROR
(POSSIBLYBACKTRACKNECESSARY,0);PUTJUMP(0);TOLOST:=TOLOST+1;'GOTO'999;2;PUTJUMP(CURLAB);TONEXT:=TONEXT+1;999;'END';
156 'PROCEDURE'PUTCONNECTOR;'BEGIN''IF'-AHEAD(SEMICOLON)'THEN''GOTO'1;PUTRESTORE;FRESHLABEL;'BEGIN''IF'-R(TONEXT=0)'THEN''GOTO'2;ERRCR(ALTERN
ATIVEEVERREACHED,0);'GOTO'999;2;'GOTO'999;'END';1;'IF'-AHEAD(CLOSE)'THEN''GOTO'4;'BEGIN''IF'-R(TONEXT=0)'THEN''GOTO'5;'GOTO'999;5;PUTLABEL(CURLA
B);FRESHLABEL;'BEGIN''IF'-GIVERSTO'THEN''GOTO'7;'GOTO'999;7;PUTJUMP(0);TOLOST:=TOLOST+1;'GOTO'999;'END';'END';4;PUTRESTORE;LOSTLABEL;999;'END';
157 'PROCEDURE'LOSTLABEL;'BEGIN''IF'-R(TOLOST=0)'THEN''GOTO'1;'BEGIN''IF'-R(TONEXT=0)'THEN''GOTO'2;'GOTO'999;2;FALSERESULT;'GOTO'999;'END';1;P
UTLABEL(0);FALSERESULT;999;'END';
158 'PROCEDURE'ENDBLABEL;'BEGIN''IF'-R(TOEND=0)'THEN''GOTO'1;'GOTO'999;1;PUTLABEL(999);999;'END';
159 'PROCEDURE'ENDJUMP;'BEGIN''IF'-R(INPT=POINT)'THEN''GOTO'1;'IF'-R(TOLOST=0)'THEN''GOTO'1;'IF'-R(TONEXT=0)'THEN''GOTO'1;'GOTO'999;1;PUTJUMP(9
99);TCEND:=TOEND+1;999;'END';
160 'PROCEDURE'PUTINIT;'BEGIN''IF'-GIVERSTO'THEN''GOTO'1;G(INITRESTORE);'GOTO'999;1;999;'END';
161 'PROCEDURE'PUTRESTORE;'BEGIN''IF'-R(TONEXT=0)'THEN''GOTO'1;'GOTO'999;1;PUTLABEL(CURLAB);'BEGIN''IF'-GIVERSTO'THEN''GOTO'3;G(DORESTORE);'G
OTO'999;3;'GOTO'999;'END';999;'END';
162 'PROCEDURE'FRESHLABEL;'BEGIN'CURLAB:=NEWLAB;NEWLAB:=NEWLAB+1;'END';
163 'PROCEDURE'TRUESRESULT;'BEGIN''IF'-WASACTION(HANDLE)'THEN''GOTO'1;'GOTO'999;1;G(HANDLE);BECOMES;G(QTRUE);SEMICOL;999;'END';
164 'PROCEDURE'FALSERESULT;'BEGIN''IF'-WASACTION(HANDLE)'THEN''GOTO'1;'GOTO'999;1;G(HANDLE);BECOMES;G(QFALSE);SEMICOL;999;'END';
165 'PROCEDURE'PUTJUMP(LAB);'INTEGER'LAB;'BEGIN'G(QGOTO);G(LAB);G(SEMICOLON);NEWLINE;APPLYLABEL(LAB);'END';
166 'PROCEDURE'PUTLABEL(LAB);'INTEGER'LAB;'BEGIN'NEWLINE;L;G(LAB);G(COLON);U;DEFINELABEL(LAB);TABLINE;'END';
167 'PROCEDURE'SEMICOL;'BEGIN'G(SEMICOLON);'END';
168 'PROCEDURE'BECOMES;'BEGIN'G(COLON);G(EQUALS);'END';
169 'PROCEDURE'BEGIN;'BEGIN'NEWLINE;G(QBEGIN);U;TABLINE;'END';
170 'PROCEDURE'ENDBR;'BEGIN'NEWLINE;L;G(QEND);'END';
171 'COMMENT'4,4,7GENERALFORMOFARULE;
172 'EOCLEAN''PROCEDURE'RULE;'BEGIN''IF'-LEFTHANDSIDE'THEN''GOTO'1;'IF'-MIDDLE'THEN''GOTO'0;'IF'-RIGHTHANDSIDE'THEN''GOTO'0;FORGETLCCALS;SMI

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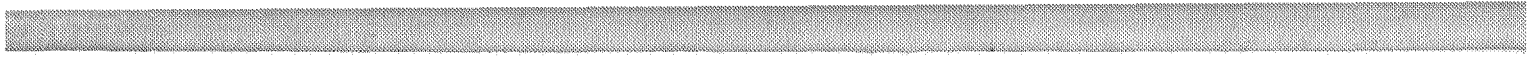
FT2LINES: 'IF'-R(POINT)' THEN 'GOTO' 0; RULE:='TRUE'; 'GOTO' 999; 1:0: RULE:='FALSE'; 999: 'END';
173 'EOCLEAN' 'PROCEDURE' 'LEFTHANDSIDE'; 'BEGIN' 'IF'-R(STAG(HANDLE)) ' THEN 'GOTO' 1; BLANKLINE; BOUNDX:=0; SPECX:=0; 'BEGIN' 'IF'-WASACTION(HANDLE) ' THEN '
'GOTO' 2; DEFINEACTION(HANDLE); RESTLHS; LEFTHANDSIDE:='TRUE'; 'GOTO' 999; 2:G(QBOOLEAN); DEFINEPREDICATE(HANDLE); RESTLHS; LEFTHANDSIDE:='TRUE'; 'GOTO' 999
; 'END'; 1:LEFTHANDSIDE:='FALSE'; 999: 'END';
174 'PROCEDURE' 'RESTLHS'; 'BEGIN' G(QPROCEDURE); G(HANDLE); OPTIONALBOUNDAFFIXES; NEWNPARS(HANDLE, BOUNDX); SEMICOLON; 'END';
175 'PROCEDURE' 'OPTIONALBOUNDAFFIXES'; 'BEGIN' 'INTEGER' AFFX; 'IF'-BOUNDAFFIX1(AFFX) ' THEN 'GOTO' 1; G(OPEN); G(AFFX); RESTBOUNDAFFIXES; 'GOTO' 999; 1: 'I
F'-BOUNDAFFIX2(AFFX) ' THEN 'GOTO' 2; G(OPEN); G(AFFX); RESTBOUNDAFFIXES; G(AFFX); 'GOTO' 999; 2:999: 'END';
176 'PROCEDURE' 'RESTBOUNDAFFIXES'; 'BEGIN' 'INTEGER' AFFX; 'IF'-BOUNDAFFIX1(AFFX) ' THEN 'GOTO' 1; G(COMMA); G(AFFX); RESTBOUNDAFFIXES; 'GOTO' 999; 1: 'IF'-
BOUNDAFFIX2(AFFX) ' THEN 'GOTO' 2; G(COMMA); G(AFFX); RESTBOUNDAFFIXES; G(AFFX); SPECX:=SPECX-1; 'BEGIN' 'IF'-R(SPECX=0) ' THEN 'GOTO' 3; 'GOTO' 999; 3:G(COMMA);
'GOTO' 999; 'END'; 2:G(CLOSE); 'BEGIN' 'IF'-R(SPECX=0) ' THEN 'GOTO' 6; 'GOTO' 999; 6:SEMICOLON; G(QINTEGER); 'GOTO' 999; 'END'; 999: 'END';
177 'EOCLEAN' 'PROCEDURE' 'BOUNDAFFIX1(AFFX); 'INTEGER' AFFX; 'BEGIN' 'IF'-R(TIMES) ' THEN 'GOTO' 1; 'IF'-R(STAG(AFFX)) ' THEN 'GOTO' 0; DEFINEAFFIX(AFFX); BO
UNDX:=BOUNDX+1; BOUNDAFFIX1:='TRUE'; 'GOTO' 999; 1:0: BOUNDAFFIX1:='FALSE'; 999: 'END';
178 'EOCLEAN' 'PROCEDURE' 'BOUNDAFFIX2(AFFX); 'INTEGER' AFFX; 'BEGIN' 'IF'-R(PLUS) ' THEN 'GOTO' 1; 'IF'-R(STAG(AFFX)) ' THEN 'GOTO' 0; DEFINEAFFIX(AFFX); BOU
NDX:=BOUNDX+1; SPECX:=SPECX+1; BOUNDAFFIX2:='TRUE'; 'GOTO' 999; 1:0: BOUNDAFFIX2:='FALSE'; 999: 'END';
179 'EOCLEAN' 'PROCEDURE' 'MIDDLE'; 'BEGIN' 'BEGIN'; OPTIONALFREEAFFIXES; 'IF'-R(COLON) ' THEN 'GOTO' 1; PUTINIT; MIDDLE:='TRUE'; 'GOTO' 999; 1: MIDDLE:='FALSE
'; 999: 'END';
180 'PROCEDURE' 'OPTIONALFREEAFFIXES'; 'BEGIN' 'INTEGER' AFFX; 'IF'-R(MINUS) ' THEN 'GOTO' 1; G(QINTEGER); FFX: 'IF'-R(STAG(AFFX)) ' THEN 'GOTO' 0; G(AFFX); DEF
INEAFFIX(AFFX); 'BEGIN' 'IF'-R(MINUS) ' THEN 'GOTO' 2; G(COMMA); 'GOTO' FFX; 2:G(SEMICOLON); 'GOTO' 999; 'END'; 1: 'GOTO' 999; 0:999: 'END';
181 'EOCLEAN' 'PROCEDURE' 'RIGHTHANDSIDE'; 'BEGIN' CURLAB:=1; NEWLAB:=2; TOLOST:=0; TOEND:=0; NXT: 'IF'-ALTERNATIVE ' THEN 'GOTO' 1; PUTCONNECTOR; 'BEGIN' 'I
F'-R(SEM. COLON) ' THEN 'GOTO' 2; 'GOTO' NXT; 2: ENDLABEL; ENDBR; SEMICOLON; RIGHTHANDSIDE:='TRUE'; 'GOTO' 999; 'END'; 1: RIGHTHANDSIDE:='FALSE'; 999: 'END';
182 'EOCLEAN' 'PROCEDURE' 'ALTERNATIVE'; 'BEGIN' 'INTEGER' LAB, OLJNEXT; TONEXT:=0; ADVANCED:='FALSE'; NXT: 'IF'-R(COMMA) ' THEN 'GOTO' 1; 'GOTO' NXT; 1: 'IF'-
R(COLON) ' THEN 'GOTO' 2; 'IF'-R(STAG(LAB)) ' THEN 'GOTO' 0; PUTJUMP(LAB); ALTERNATIVE:='TRUE'; 'GOTO' 999; 2: 'IF'-R(OPEN) ' THEN 'GOTO' 3; OLDNEXT:=TONEXT; 'IF'-C
HOICE ' THEN 'GOTO' 0; 'IF'-R(CLOSE) ' THEN 'GOTO' 0; 'BEGIN' 'IF'-G(VERSTO) ' THEN 'GOTO' 4; TONEXT:=1; ALTERNATIVE:='TRUE'; 'GOTO' 999; 4: TONEXT:=OLDNEXT; ALTERN
ATIVE:='TRUE'; 'GOTO' 999; 'END'; 3: 'IF'-MEMBER ' THEN 'GOTO' 6; 'GOTO' NXT; 6: TRUERESULT; ENDJUMP; ALTERNATIVE:='TRUE'; 'GOTO' 999; 0: ALTERNATIVE:='FALSE'; 999
: 'END';
183 'EOCLEAN' 'PROCEDURE' 'CHOICE'; 'BEGIN' 'INTEGER' LAB; LAB:=CURLAB; 'BEGIN' PUTINIT; FRESHLABEL; NXT: 'IF'-ALTERNATIVE ' THEN 'GOTO' 1; PUTCONNECTOR; 'BEGI
N' 'IF'-R(SEMICOLON) ' THEN 'GOTO' 2; 'GOTO' NXT; 2: ENDBR; SEMICOLON; CURLAB:=LAB; CHOICE:='TRUE'; 'GOTO' 999; 'END'; 1: CHOICE:='FALSE'; 999: 'END';
184 'EOCLEAN' 'PROCEDURE' 'MEMBER'; 'BEGIN' 'INTEGER' HEAD; 'IF'-R(STAG(HEAD)) ' THEN 'GOTO' 1; APPLY(HEAD); 'BEGIN' 'IF'-R(COLON) ' THEN 'GOTO' 2; PUTLABEL(HEA
D); MEMBER:='TRUE'; 'GOTO' 999; 2: 'IF'-TRYTYPE(HEAD, PREDICATE) ' THEN 'GOTO' 3; G(QIF); G(QNOT); PUTAFFIXEXPRESSION(HEAD); G(QTHEN); PUTCONNECTAND; ADVANCED:='
TRUE'; MEMBER:='TRUE'; 'GOTO' 999; 3: 'IF'-TRYTYPE(HEAD, FLAG) ' THEN 'GOTO' 4; G(QIF); G(QNOT); PUTAFFIXEXPRESSION(HEAD); G(QTHEN); PUTCONNECTAND; MEMBER:='TR
UE'; 'GOTO' 999; 4: 'IF'-TRYTYPE(HEAD, ACTION) ' THEN 'GOTO' 5; PUTAFFIXEXPRESSION(HEAD); SEMICOLON; MEMBER:='TRUE'; 'GOTO' 999; 5: 'GOTO' 0; 'END'; 1:0: MEMBER:='FA
LSE'; 999: 'END';
185 'COMMENT' '4.4.8 OTHER BUILDING STONES OF A GRAMMAR';
186 'EOCLEAN' 'PROCEDURE' 'COMMAND'; 'BEGIN' 'IF'-R(RSTOON) ' THEN 'GOTO' 1; GIVERSTO:='TRUE'; COMMAND:='TRUE'; 'GOTO' 999; 1: 'IF'-R(RSTCOFF) ' THEN 'GOTO' 2
; GIVERSTO:='FALSE'; COMMAND:='TRUE'; 'GOTO' 999; 2: 'IF'-R(SHORT) ' THEN 'GOTO' 3; LEGIBLE:='FALSE'; COMMAND:='TRUE'; 'GOTO' 999; 3: 'IF'-R(LONG) ' THEN 'GOTO' 4
; LEGIBLE:='TRUE'; COMMAND:='TRUE'; 'GOTO' 999; 4: 'IF'-R(TRACEON) ' THEN 'GOTO' 5; GIVE-RACE:='TRUE'; COMMAND:='TRUE'; 'GOTO' 999; 5: 'IF'-R(TRACEOFF) ' THEN 'G
OTO' 6; GIVETRACE:='FALSE'; COMMAND:='TRUE'; 'GOTO' 999; 6: COMMAND:='FALSE'; 999: 'END';
187 'EOCLEAN' 'PROCEDURE' 'COMMENT'; 'BEGIN' 'IF'-R(SUB) ' THEN 'GOTO' 1; BLANKLINE; G(QCOMMENT); RST: 'BEGIN' 'IF'-R(BUS) ' THEN 'GOTO' 2; G(SEMICOLON); COMME
NT:='TRUE'; 'GOTO' 999; 2:G(INPT); READ(INPT); 'GOTO' RST; 'END'; 1: COMMENT:='FALSE'; 999: 'END';
188 'EOCLEAN' 'PROCEDURE' 'STARTINGSYMBOL'; 'BEGIN' 'INTEGER' HEAD; 'IF'-R(RESULT) ' THEN 'GOTO' 1; 'IF'-R(STAG(HEAD)) ' THEN 'GOTO' 0; NEWPAGE; 'BEGIN' 'IF'-R
YTYPE(HEAD, ACTION) ' THEN 'GOTO' 2; APPLYACTION(HEAD); TERMINALS; G(HEAD); STARTINGSYMBOL:='TRUE'; 'GOTO' 999; 2: APPLYPREDICATE(HEAD); TERMINALS; G(HEAD); ST
ARTINGSYMBOL:='TRUE'; 'GOTO' 999; 'END'; 1:0: STARTINGSYMBOL:='FALSE'; 999: 'END';
189 'PROCEDURE' 'SKIP UNTIL POINT'; 'BEGIN' NXT: ERROR(SKIPPED, INPT); 'BEGIN' 'IF'-R(POINT) ' THEN 'GOTO' 2; 'GOTO' 999; 2: NEXTSYMBOL; 'GOTO' NXT; 'END'; 999: 'E
ND';
190 'COMMENT' '4.4.9 TERMINALS';
191 'EOCLEAN' 'PROCEDURE' 'TERMINALS'; 'BEGIN' BLANKLINE; FORGETGLOBALS; NEWLINE; 'IF'-R(PLOC=500001) ' THEN 'GOTO' 1; 'GOTO' 999; 1: APPLYACTION(XREAD); APPLYACTION(X
OUT); APPLYACTION(INITREAD); BLANKLINE; G(INITREAD); G(SEMICOLON); NEWLINE; BLANKLINE; FORGETTERMINALS; NEWLINE; BLANKLINE; 999: 'END';
192 'PROCEDURE' 'FORGETGLOBALS'; 'BEGIN' 'INTEGER' P, TERM; P:=600001; PLOC:=500001; NXT: 'IF'-R(P=PGLOB) ' THEN 'GOTO' 1; 'GOTO' 999; 1: TERM:=LGLOBAL[P]; MAYBET
ERMINAL(TERM); P:=P+1; 'GOTO' NXT; 999: 'END';
193 'PROCEDURE' 'MAYBETERMINAL(X); 'INTEGER' X; 'BEGIN' 'IF'-WASTERMINAL(X) ' THEN 'GOTO' 1; PUTDECL(X); ENTERLOCAL(X); 'GOTO' 999; 1: WARN(X); 999: 'END';
194 'EOCLEAN' 'PROCEDURE' 'WASTERMINAL(X); 'INTEGER' X; 'BEGIN' 'INTEGER' S; S:=TTAG[X-5]; 'IF'-R(S=APPLIED) ' THEN 'GOTO' 1; S:=TTAG[X-4]; 'IF'-R(S=GLOBAL) '
THEN 'GOTO' 1; S:=TTAG[X-3]; 'IF'-R(S=PCINTER) ' THEN 'GOTO' 1; WASTERMINAL:='TRUE'; 'GOTO' 999; 1: WASTERMINAL:='FALSE'; 999: 'END';
195 'PROCEDURE' 'PUTDECL(X); 'INTEGER' X; 'BEGIN' NEWLINE; G(QINTEGER); G(X); G(SEMICOLON); POSITION(64); INFORM(X); 'END';
196 'PROCEDURE' 'FORGETTERMINALS'; 'BEGIN' 'INTEGER' P, TERM; P:=500001; NXT: 'IF'-R(P=PLOC) ' THEN 'GOTO' 1; 'GOTO' 999; 1: TERM:=LLOC[P]; PUTCALL(XREAD, TERM)
; P:=P+1; 'GOTO' NXT; 999: 'END';
197 'PROCEDURE' 'PUTCALL(PROC, X); 'INTEGER' X, PROC; 'BEGIN' NEWLINE; G(PROC); G(OPEN); G(X); G(CLOSE); G(SEMICOLON); 'END';
198 'COMMENT' '4.4.10 POSTMORTEM';
199 'PROCEDURE' 'POSTMORTEM'; 'BEGIN' NLCR; OUTINT(100001); OUTINT(PTAG); OUTINT(110000); NLCR; OUTINT(200001); OUTINT(PBOLD); OUTINT(200600); NLCR; OUTIN
T(300001); OUTINT(PSPEC); OUTINT(300100); NLCR; OUTINT(400001); OUTINT(PCONS); OUTINT(400300); NLCR; OUTINT(500001); OUTINT(XLOC); OUTINT(500200); NLCR; OUT
INT(60000); OUTINT(PGLOB); OUTINT(601000); NLCR; OUTINT(700001); OUTINT(PMACR); OUTINT(701100); NLCR; OUTINT(800001); OUTINT(8TEXT); OUTINT(804000); NLCR;

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OUTINT(LINE);CUTINT(CARD);NLCR;NLCR;NLCR;DICTIONARY;'END';
200 'PROCEDURE'DICTIONARY;'BEGIN'LISTTAGS(FIRSTTAG);'END';
201 'PROCEDURE'LISTTAGS(X);'INTEGER'X;'BEGIN''INTEGER'P,STATE, SORT,TYPE;'IF'-(X=0)'THEN''GOTO'1;'GOTO'999;1:P:=TTAG[X-2];LISTTAGS(P);STATE:=
TTAG[X-5];SCRT:=TTAG[X-4];TYPE:=TTAG[X-3];'IF'-(SORT=0)'THEN''GOTO'2;'GOTO'SKP;2:NLCR;OUT(X);POSITION(32);OUT(STATE);OUT(SORT);OUT(TYPE);POSITIO
N(59);ENTRIES(X);SKP:P:=TTAG[X-1];LISTTAGS(P);999:'END';
202 'PROCEDURE'ENTRIES(P);'INTEGER'P;'BEGIN''INTEGER'Q,H,K;Q:=TTAG[P-6];'IF'-(Q=0)'THEN''GOTO'1;NLCR;'GOTO'999;1:LINE:H:=LTEXT[Q];Q:=H/'8192
;K:=H-8192*Q;CUTINT(K);'BEGIN''IF'-(Q=0)'THEN''GOTO'3;NLCR;'GOTO'999;3:Q:=Q+800001;Q:=Q-1;'GOTO'LINE;'END';999:'END';
203 'COMMENT'4.5.HEARTOFCOMPILERCOMPILER;
204 'INTEGER'INITRESTORE,DORESTORE,XREAD,XOUT,INITREAD,TITLE;
205 'PROCEDURE'SENTENCE;'BEGIN'STARTSYSTEM;BLANKLINE;BEGIN;COMPILERDESCRIPTION;ENDBR;BLANKLINE;NEWPAGE;POSTMORTEM;'END';
206 'PROCEDURE'STARTSYSTEM;'BEGIN'READ(INITRESTORE);READ(DOESTORE);READ(XREAD);READ(XOUT);READ(INITREAD);READ(TITLE);GIVERSTO='FALSE';GIVE
TRACE='FALSE';GIVETEXT='TRUE';PLOC:=500001;PLOB:=600001;PTEXT:=800001;PMACR:=700001;CARD:=0;CPOS:=1;XLOC:=0;LINEDUP='FALSE';LEGIBLE='TRUE';
INDENTATION:=0;LINE:=0;POS:=0;OUT(TITLE);READ(INPT);'END';
207 'PROCEDURE'COMPILERDESCRIPTION;'BEGIN'NXT;'IF'-'SPECIFICATION'THEN''GOTO'1;'GOTO'NXT;1:'IF'-'DECLARATION'THEN''GOTO'2;'GOTO'NXT;2:'IF'-'COM
MAND'THEN''GOTO'3;'GOTO'NXT;3:'IF'-'COMMENT'THEN''GOTO'4;'GOTO'NXT;4:'IF'-'STARTINGSYMBOL'THEN''GOTO'5;'GOTO'999;5:'IF'-'RULE'THEN''GOTO'6;'GOTO'NX
T;6:SKIPUNTILPOINT;'GOTO'NXT;999:'END';
208 'INTEGER'TAGFULL;'INTEGER'BCDFULL;'INTEGER'SPECFULL;'INTEGER'CONSFULL;'INTEGER'WRONGINPUTCODE;'INTEGER'BLANK;'INTEGER'LOCFULL;'INTEGER'
GLOBFULL;'INTEGER'MACRFULL;'INTEGER'TEXTFULL;'INTEGER'DEFINED;'INTEGER'GLOBAL;'INTEGER'ACTION;'INTEGER'PREDICATE;'INTEGER'LOCAL;'INTEGER'POINTER
;'INTEGER'LABEL;'INTEGER'APPLIED;'INTEGER'XIMPOSSIBLE;'INTEGER'MACRO;'INTEGER'WRONGNUMBEROFPARAMETERS;'INTEGER'FLAG;'INTEGER'LIST;'INTEGER'EXTER
NAL;'INTEGER'COMMA;'INTEGER'POINT;'INTEGER'QINTEGER;'INTEGER'SEMICOLON;'INTEGER'QBOOLEAN;'INTEGER'QARRAY;'INTEGER'SUB;'INTEGER'COLON;'INTEGER'BU
S;'INTEGER'PLUS;'INTEGER'MINUS;'INTEGER'OPEN;'INTEGER'CLOSE;'INTEGER'STRINGQUOTE;'INTEGER'WRONGAFFIX;'INTEGER'ONE;'INTEGER'TWO;'INTEGER'THREE;'I
NTEGER'FOUR;'INTEGER'FIVE;'INTEGER'XPARAMETERERROR;'INTEGER'POSSIBLYBACKTRACKNECESSARY;'INTEGER'ALTERNATIVEEVERREACHED;'INTEGER'QTRUE;'INTEGER'
QFALSE;'INTEGER'QGOTO;'INTEGER'EQUALS;'INTEGER'QBEGIN;'INTEGER'QEND;'INTEGER'QPROCEDURE;'INTEGER'TIMES;'INTEGER'QIF;'INTEGER'QNOT;'INTEGER'QTHEN;'
INTEGER'QSTOON;'INTEGER'QSTOOFF;'INTEGER'SHORT;'INTEGER'LONG;'INTEGER'TRACEON;'INTEGER'TRACEOFF;'INTEGER'QCOMMENT;'INTEGER'RESULT;'INTEGER'SKIP
PED;
209 INITIALIZEFORREADING;
210 READ(TAGFULL);READ(BOLDFULL);READ(SPECFULL);READ(CONSFULL);READ(WRONGINPUTCODE);READ(BLANK);READ(LOCFULL);READ(GLOBFULL);READ(MACRFULL);
READ(TEXTFULL);READ(DEFINED);READ(GLOBAL);READ(ACTION);READ(PREDICATE);READ(LOCAL);READ(POINTER);READ(LABEL);READ(APPLIED);READ(XIMPOSSIBLE);REA
D(MACRO);READ(WRONGNUMBEROFPARAMETERS);READ(FLAG);READ(LIST);READ(EXTERNAL);READ(COMMA);READ(POINT);READ(QINTEGER);READ(SEMICOLON);READ(QBOOLEAN
);READ(QARRAY);READ(SUB);READ(COLON);READ(BUS);READ(PLUS);READ(MINUS);READ(OPEN);READ(CLOSE);READ(STRINGQUOTE);READ(WRONGAFFIX);READ(ONE);READ(T
WO);READ(THREE);READ(FOUR);READ(FIVE);READ(XPARAMETERERROR);READ(POSSIBLYBACKTRACKNECESSARY);READ(ALTERNATIVEEVERREACHED);READ(QTRUE);READ(QFAL
SE);READ(QGOTO);READ(QEQUALS);READ(QBEGIN);READ(QEND);READ(QPROCEDURE);READ(TIMES);READ(QIF);READ(QNOT);READ(QTHEN);READ(QSTOON);READ(QSTOOFF);REA
D(QSHORT);READ(QLONG);READ(QTRACEON);READ(QTRACEOFF);READ(QCOMMENT);READ(RESULT);READ(SKIPPED);
211 SENTENCE'END'

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References.

- [1] J. Feldman, D. Gries, "Translator Writing Systems", Communications of the ACM, Volume 11 number 2 (1968).
- [2] C.H.A. Koster, "Affix Grammars", in: J.E.L. Peck (Editor), "ALGOL 68 implementation", North-Holland Publishing Company (1971).
- [3] A. van Wijngaarden (Editor), "Report on the Algorithmic Language ALGOL 68", Numerische Mathematik, 14 (1969).
- [4] P. Naur (Editor), "Revised Report on the Algorithmic Language ALGOL 60", Regnecentralen, Copenhagen (1962).
- [5] D.E. Knuth, "Semantics of Context Free Languages", Math. Systems Theory 2, (1968).
- [6] D.E. Knuth, "Top-down syntax analysis", International Summer School on Computer Programming, Copenhagen (1967).
- [7] J. Earley, H. Sturgis, "A Formalism for Translator Interactions", Communications of the ACM, Volume 13 number 10 (1970).



The differences between versions 17 and 19 of the CDL-compiler; first update to MR 127/71.

1. p. 23, bottom line, ignored, add:
; moreover any sequence of symbols, not containing a
\$-character, enclosed between \$-characters is also ignored
2. p. 24, l. 16, equals =, add:
not ¬
3. p. 24, l. 25, list ..., add:
comment symbol 'comment'
4. p. 25, l. 17, flag symbol, add:
; list symbol
5. p. 26, l. 2 from bottom, list (global, add:
or macro
6. p. 27, l. 13, 'external' ... flag, add:
'external' 'list' defined global list
7. p. 27, l. 19, 'macro'.... pointer, add:
'macro' 'list' defined macro list
8. p. 30, l. 8, 2.6. Comments. *input*; , replace by
2.6. Comments.
comment: skipped comment; translated comment.
skipped comment: sub, rest skipped comment.
rest skipped comment: bus; nonbus, rest skipped comment.
translated comment: comment symbol, rest translated comment.
rest translated comment: point; nonpointnonsemicolon,
rest translated comment.

(Here, nonbus stands for any symbol except] and nonpoint
nonsemicolon for any symbol except . or ; .)

A skipped comment is skipped and a translated comment is translated into an ALGOL 60 <comment>; e.g. [input] is skipped while 'comment' input. is translated into '*comment*' *input*; . The text of a comment may only consist of allowed CDL-symbols (see 2.2), as opposed to the sequence of characters between two \$-characters (see 2.2).

9. p. 30, bottom line, replace by:

(semicolon, right hand side;) .

10. p. 31, lines 1-8, replace by:

alternative:

label, colon, alternative;
 affix expression, (comma, alternative;);
 group; jump; .

affix expression:

handle, optional affixes;
 not, tag.

11. p. 32, line 10, affix expression, add:

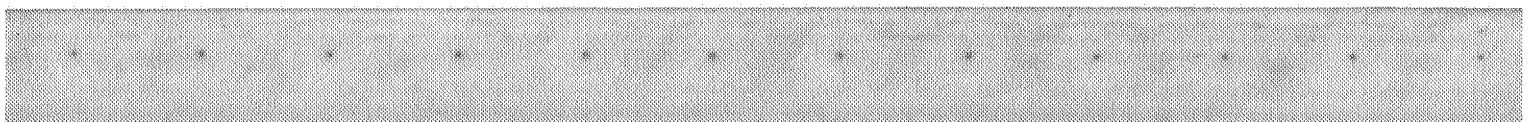
that starts with a handle

12. p. 33, after line 20, α ; , add:

The translation of an affix expression that starts with a not is:

if tag then goto λ ;

where λ is determined by means of the algorithm on the previous page.



13. p. 33, bottom, , add:

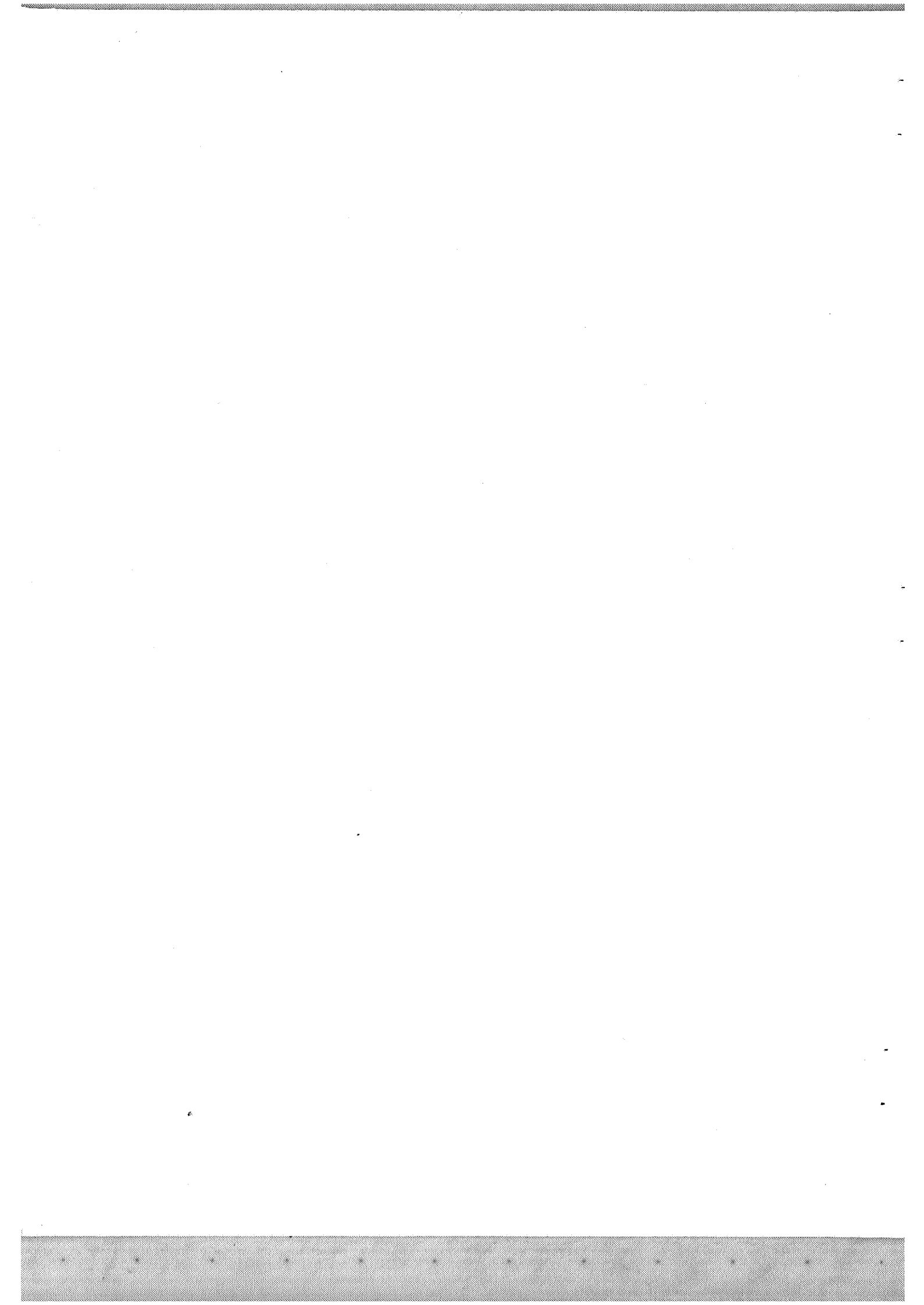
The above translation scheme makes use of integer labels, which, for some applications, are undesirable. Therefore a symbol is prefixed to every label in the translation; this symbol is the terminal "label prefix" of the CDL-compiler, which is read in during the initialization phase of the compiler. The prefix should be chosen so that the resulting names will not clash with existing names. The present value of "label prefix" is "l"; in the examples in this report its value is "" (i.e. the empty symbol).

14. p. 36, l. 4, delete: Various themselves.

15. p. 36, l. 25, occurrence. , add:

Line numbers of specification are preceded by S:,
line numbers of definition are preceded by D:.

16. after page 36 add the following pages:



2.10. Error detection and recovery.

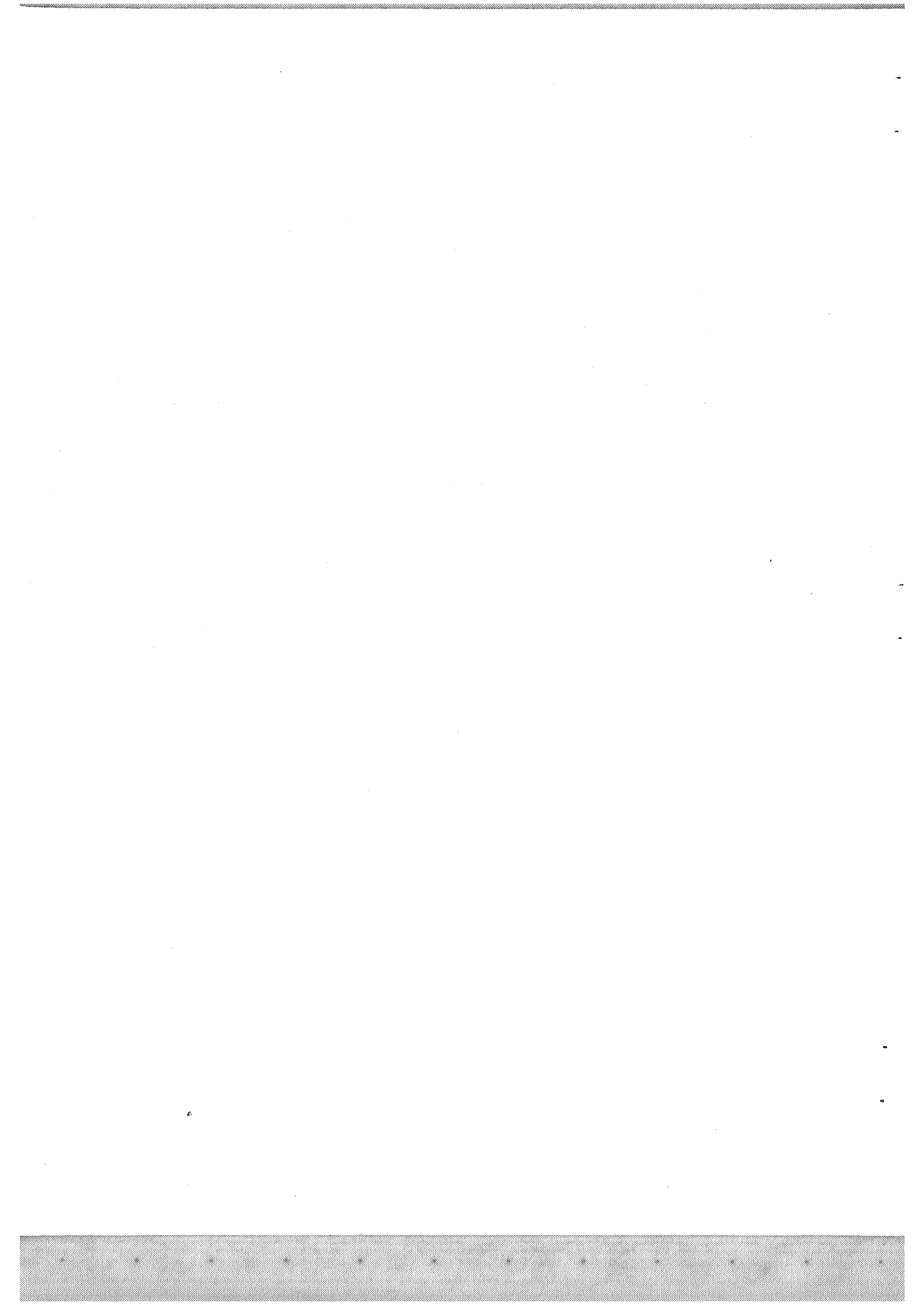
All errors fall in one of the following four classes.

1) Warnings:

backtrack?	An alternative may fail although one of its predicates may have succeeded.
alternative never reached	One of the previous alternatives in this right hand side will always succeed.
nonfalse	This predicate will always succeed, and is in effect an action.
may be false	Possibly none of the alternatives of this action will succeed.

2) Unexpected symbols:

wrong input code	Some unreadable item is met on the input medium.
symbol missing, assumed: symbol	<ol style="list-style-type: none"> 1. The grammar requires at this point the presence of the indicated symbol, which, however, was not found in the input. The symbol is assumed to be present, so that e.g. "'list' a[1:10." will yield a correct translation. 2. The grammar requires at this point the presence of a tag, which, however, was not found in the input. The tag "dummy" is inserted. 3. The grammar requires at this point the presence of one of the symbols: <ul style="list-style-type: none"> 'action' 'predicate'



'pointer'

'flag'

'list'

None of these, however, was found. The symbol 'predicate' was inserted.

symbol not allowed inside comment symbol

The indicated symbol violates the restrictions on the symbols allowed inside a comment. In this version of the compiler the symbol can only be a semicolon.

affix missing

The grammar requires at this point the presence of an affix, which, however, was not found. The tag "dummy" was inserted.

skipped --- state, sort, type, symbol

The indicated symbol was skipped because it was unacceptable as the first symbol of a compiler description.

3) Conflict of attributes

impossible redefinition state, sort, type, tag, new state, new sort,
new type

The indicated tag with the indicated state, sort and type is specified, defined or applied with the conflicting attributes new state, new sort and new type.

wrong number of parameters state, sort, type, tag

The number of affixed following the indicated tag is incompatible with that number in the definition of the tag.

wrong sort state, sort, type, tag

The indicated local tag is applied as a handle.

