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A NOTE ON AUTOMATIC STORAGE OF
ARBITRARY TREES IN ALGOL 60

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0. Summary

In this note a method is demonstrated for making an elegant ALGOL 60 program which stores a tree in a one-dimensional array.

1. The problem

Consider a set s of trees t ; each tree being either an atom (representable by an integer) or the root of an arbitrary number n ($n > 0$) of other trees t_i , $i = 1, \dots, n$, $t_i \in s$.

The trees are assumed to be pure in the sense that no two trees are subtrees of each other.

It is supposed that the trees have to be stored once, have to be inspected often, but will not be changed or deleted.

In that case the following storage organization may be chosen.

Declare the integer array $INF[1: bound]$.

If t is an atom, then store the integer representing this atom in $INF[t]$, where t is some appropriate value, called the location of t ; otherwise, if t is the root of n trees t_i , with locations t_i , $i = 1, \dots, n$, fill $n+1$ consecutive array elements $INF[t], \dots, INF[t+n]$ as follows:

$INF[t] := t+n+1$ and

for $i = 1, \dots, n$: if t_i is a root then $INF[t+i] := t_i$, otherwise

$t+i = t_i$ and $INF[t+i]$ contains the integer representing the atom t_i ,

where t is again some appropriate value, called the location of t .

A tree consisting of N atoms and M roots will occupy $N+2M$ places in INF . For the storage of this tree a program has to be written which performs the $N+2M$ necessary assignments. For any given tree it is not too difficult to write down these assignment statements; the order of which has to be chosen judiciously however.

The problem is to construct auxiliary procedures, which perform the assignment statements, by means of which for every arbitrary tree a program can be made with a structure reflecting directly the structure of this tree.

In section 3 the solution the solution of the problem will be given in the form of the procedures *STORE TREE* and *L*, which store an arbitrary tree consisting of at least one root.

Before doing this it is necessary to introduce, in the next section, means for treating lists with a variable number of elements as actual parameters in procedure statements.

2. Lists of variable length as parameter

Consider a set of lists, each list consisting of a variable number n of elements e_i , $i = 1, \dots, n$; these elements have to be treated by a standard process P in the ordering e_1, \dots, e_n . Before treating the elements, an initialization process P_0 has to be executed which uses the number n .

A way of programming this in ALGOL 60, where it is not allowed to have a variable number of actual parameters in a procedure statement, is the following:

Declare the procedure L :

```

Boolean procedure  $L(\textit{first}, \textit{ei})$ ; Boolean  $\textit{first}$ ; integer  $\textit{ei}$ ;
begin comment  $\textit{count the number of elements}$ ;  $\textit{cnt} := \textit{cnt} + 1$ ;
    if  $\textit{first}$  then  $P_0$ ;
     $P$ ;  $L := \textit{false}$ 
end  $L$ 

```

The list (e_1, e_2, \dots, e_n) is then treated by the statement:

```

 $L(\dots L(L(\textit{true},$ 
     $e_1),$ 
     $e_2),$ 
    ...
     $e_n)$ .

```

3. Storage of trees

A procedure *STORE TREE* and a procedure L are given now, by means of which the tree of the following example is stored in the array *INF*. Note that *STORE TREE* stores trees consisting of at least one root.

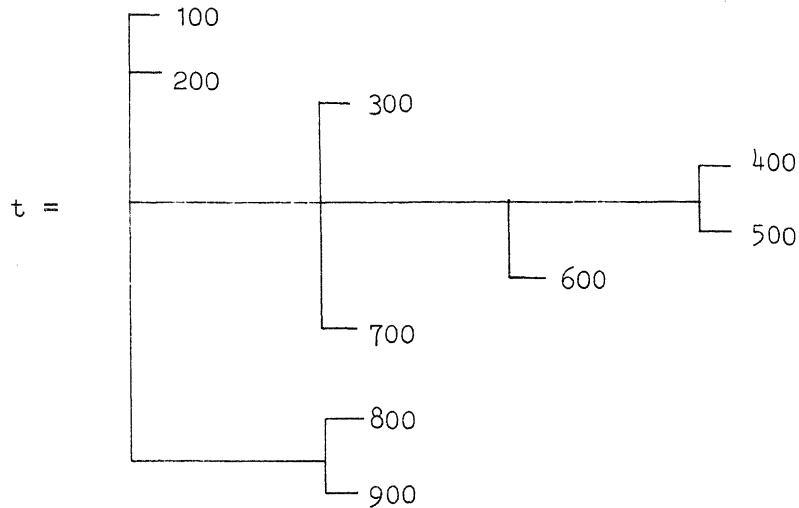
3.1. Example

In the ALGOL 60 program of section 3.2. the following tree t is stored.

In bracketed form:

$$t = (100, 200, (300, ((400, 500), 600), 700), (800, 900)).$$

Two-dimensionally t looks like:



3.2. The ALGOL 60 program

```

begin integer ptr, fill ptr, t;
  integer array INF[1:1000];
  integer procedure STORE TREE(list); Boolean list;
  begin STORE TREE := fill ptr := ptr := ptr + 1;
  if list then;
  end STORE TREE;

  Boolean procedure L(first, inf); Boolean first; integer inf;
  begin integer loc fill ptr;
  ptr := ptr + 1;
  if first then
  begin INF[fill ptr] := ptr + 1;
  fill ptr := fill ptr + 1;
  if ptr > 1000 then EXIT
  end;
  loc fill ptr := fill ptr;
  INF[loc fill ptr] := inf;
  fill ptr := loc fill ptr + 1;
  L := false
  end L;

```

```

ptr:= 0;
t:= STORE TREE(L(L(L(L(true,
    100),
    200),
    STORE TREE(L(L(L(true,
        300),
        STORE TREE(L(L(true,
            STORE TREE(L(L(true,
                400),
                500))),
            600))),
        700))),
    STORE TREE(L(L(true,
        800),
        900))));
NLCR; PRINTTEXT(the contents of INF is displayed.
Each line contains the information of one root);
t:= 0; fill ptr:= 1; for t:= t + 1 while t < ptr do
begin if t = fill ptr then
    begin NLCR; fill ptr:= INF[fill ptr] end;
    ABSFIXT(2,0,t); FIXT(3,0,INF[t]); SPACE(5)
end end

```

the contents of INF is displayed.
Each line contains the information of one root

1	+6	2	+100	3	+200	4	+6	5	+16
6	+10	7	+300	8	+10	9	+700		
10	+13	11	+13	12	+600				
13	+16	14	+400	15	+500				
16	+19	17	+800	18	+900				