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TWO OPTIMAL CONSTANT WEIGHT CODES

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Two optimal constant weight codes

by

A.E. Brouwer & A. Schrijver

ABSTRACT

We write the complete graph  $K_{22}$  as a line-disjoint union of 37 copies of  $K_4$  and one copy of  $K_{3,3}$  and show that it is impossible to pack 38 copies of  $K_4$  in  $K_{22}$ . This implies  $A(22,6,4) = 37$  and  $A(21,6,4) = 31$  if  $A(n,d,w)$  denotes the maximum number of code words of weight  $w$  possible in a binary code of word length  $n$  and minimum distance  $d$ .

KEY WORDS & PHRASES: *packing, constant weight code, Room square.*

## 0. INTRODUCTION

A (binary) *constant weight code* with *word length*  $n$  and *weight*  $w$  is a collection of vectors in  $\text{GF}(2)^n$  such that each of the vectors has exactly  $w$  components equal to 1.

The *distance* between two vectors is the number of coordinates at which they differ.

Let  $A(n,d,w)$  denote the maximum number of vectors possible in a constant weight code with word length  $n$  and weight  $w$  such that the distance between any two distinct vectors is at least  $d$ . In this note we determine the values of  $A(n,6,4)$  for  $n = 21$  and  $n = 22$ .

Another terminology is possible: if we consider each code word as a characteristic function, then a constant weight code can be identified with a collection of  $w$ -subsets of a given  $n$ -set. The requirement on the distance can be translated as follows: Let  $\lambda$  be the maximum cardinality of an intersection of two  $w$ -subsets in the collection; then  $2(w-\lambda) \geq d$ .

This leads to a problem on hypergraphs. In our case ( $d=6, w=4$ ) we find  $\lambda \leq 1$ , i.e. no two 4-subsets can have a pair in common. So this is the case of a graph and we must try to find as many edge-disjoint  $K_4$ 's as possible in a  $K_n$ .

From the work of HANANI it follows that  $K_n$  can be partitioned into disjoint  $K_4$ 's iff  $n \equiv 1$  or  $4 \pmod{12}$ . Shortening these codes we get optimal codes for  $n \equiv 0$  or  $3 \pmod{12}$ . In the remaining cases, only for small  $n$  were some values of  $A(n,6,4)$  known. Recently we proved the existence of a partition of  $K_n$  into a matching and  $\frac{1}{12} n(n-2)$   $K_4$ 's for each  $n \equiv 2 \pmod{12}$ , thus proving that  $A(n,6,4) = \frac{1}{12} n(n-2)$  for these  $n$  (see [2]).

As far as we know the current state of knowledge concerning the value of  $A(n,6,4)$  for  $n \leq 25$  is reflected by the following table:-

n	0-3	4-6	7-8	9	10	11	12	13	14	15	16
A(n,6,4)	0	1	2	3	5 <sup>B</sup>	6	9 <sup>BS</sup>	13 <sup>BCS</sup>	14 <sup>C</sup>	15 <sup>CS</sup>	20 <sup>S</sup>

n	17	18	19	20	21	22	23	24	25
A(n,6,4)	20 <sup>X</sup>	20-22	24 <sup>T</sup> -27	29 <sup>T</sup> -30	31 <sup>*</sup>	37 <sup>*</sup>	37 <sup>*</sup> -40	42 <sup>S</sup>	50 <sup>S</sup>

B = transpose of a block design

C = circulant on (11001010...0)

S = Steiner system or shortened Steiner system

T = using a transversal design TD (5,1;20)

X = see [1]

\* = this note

## 1. UPPER BOUNDS

Upper bounds for  $A(n,d,w)$  are given by JOHNSON [4] as follows:

$$A(n,d,w) \leq \frac{n}{w} A(n-1,d,w-1)$$

and

$$A(n,d,w) \leq \frac{n}{n-w} A(n-1,d,w).$$

Since obviously  $A(n,6,3) = \lfloor \frac{n}{3} \rfloor$  we get from the first inequality

$$A(n,6,4) \leq \frac{n}{4} \lfloor \frac{n-1}{3} \rfloor,$$

i.e.,

$$A(n,6,4) \leq \begin{cases} \frac{1}{12} n(n-3) & \text{if } n \equiv 0 \pmod{3}, \\ \frac{1}{12} n(n-1) & \text{if } n \equiv 1 \pmod{3}, \\ \frac{1}{12} n(n-2) & \text{if } n \equiv 2 \pmod{3}. \end{cases}$$

Substituting these bounds in the second inequality,

$$A(n,6,4) \leq \frac{n}{n-4} A(n-1,6,4),$$

we obtain better bounds when  $n \equiv 7$  or  $10 \pmod{12}$ :

$$A(n,6,4) \leq \frac{1}{12} (n(n-1)-18) \quad \text{if } n \equiv 7 \text{ or } 10 \pmod{12}$$

and the same bound as before when  $n \equiv 1$  or  $4 \pmod{12}$ .

Therefore define the Johnson bound  $J(n,6,4)$  to be

$$J(n,6,4) = \begin{cases} \left\lfloor \frac{n}{4} \left\lfloor \frac{n-1}{3} \right\rfloor \right\rfloor - 1 & \text{if } n \equiv 7 \text{ or } 10 \pmod{12}, \\ \left\lfloor \frac{n}{4} \left\lfloor \frac{n-1}{3} \right\rfloor \right\rfloor & \text{otherwise.} \end{cases}$$

This means that if  $A(3m+1,6,4)$  attains the Johnson bound, then so does  $A(3m,6,4)$ .

In particular, when we have shown that  $A(22,6,4) = \frac{1}{12} (22 \cdot 21 - 18) = 37$ , then it immediately follows that  $A(21,6,4) = \left\lfloor \frac{1}{12} \cdot 21 \cdot 18 \right\rfloor = 31$ .

## 2. THE CONSTRUCTION

We now construct a collection of 37 4-tuples on 22 points such that each pair occurs together in at most one of the 4-tuples, thus showing that  $A(22,6,4) = 37$ .

Let  $V_1 = \{a, b, c, d, e, f, g\}$ ,  $V_2 = \{A, B, C, D, E, F, G\}$  and  $V_3 = \{0, 1, 2, 3, 4, 5, 6, 7\}$ .

Consider the following Room square, indexed by  $V_1 \times V_2$ :

	A	B	C	D	E	F	G
a	01	-	45	67	-	-	23
b	57	02	-	-	-	13	46
c	-	56	03	12	-	47	-
d	-	37	-	04	26	-	15
e	36	14	27	-	05	-	-
f	24	-	-	35	17	06	-
g	-	-	16	-	34	25	07

This Room square has the following properties (verified by inspection):

- (i) In each row each digit occurs exactly once.
- (ii) In each column each digit occurs exactly once.
- (iii) Each of the  $\binom{8}{2}$  unordered pairs from  $V_3$  occurs exactly once in the square.

[These three properties are just the definition of a Room square.]

- (iv)  $\begin{pmatrix} - & - \\ - & - \end{pmatrix}$  is not a minor of this matrix.

[That is, the positions of the minus signs form the incidence matrix of  $PG(2,2)$ .]

Let  $V = V_1 \cup V_2 \cup V_3$ ; then  $|V| = 22$ . Choose the following 4-tuples from  $V$ :

$\{xXij\}$ , where  $\{ij\}$  is the unordered pair at position  $xX$  of the Room square,

$\{xXYZ\}$ , where  $xX$ ,  $xY$  and  $xZ$  are the unoccupied positions in row  $x$  of the square,

$\{abcd\}$  and  $\{aefg\}$ .

This gives  $28 + 7 + 2 = 37$  4-tuples, and it is easily verified that no pair occurs twice. The  $\binom{22}{2} - 37\binom{4}{2} = 9$  pairs which have not been used form a  $K_{3,3}$  on the set  $\{b,c,d,e,f,g\}$ .

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