|  | stichting <br> mathematisch centrum |  |
| :---: | :---: | :---: |
| AFDELING MATHEMATISCHE STATISTIEK (DEPARTMENT OF MATHEMATICAL STATISTICS) |  | JANUARI |
| A.J. VAN ES, R.D. GILL \& C. VAN PUTTEN |  |  |
| Preprint |  |  |
|  | kruislaan 4131098 SJ | amsterdam |
| mbuotheex | MATHEMATISC:- (EvNEM aMSTERDAM $\qquad$ |  |

## Printed at the Mathematical Centre, Kruislaan 413, Amsterdam, The Netherlands.

The Mathematical Centre, founded 11th February 1946, is a non-profit institution for the promotion of pure and applied mathematics and computer science. It is sponsored by the Netherlands Government through the Netherlands Organization for the Advancement of Pure Research (Z.W.O.).

[^0]Random number generators for a pocket calculator *)
by
A.J. van Es, R.D. Gill \& C. Van Putten ${ }^{* *) ~}$

## ABSTRACT

Some theory of linear congruential pseudo random number generators $x_{n+1}=\left(a x_{n}+c\right) \bmod m$ is summarized for the case in which the modulus $m$ is a power of 10. These generators are especially suitable for implementation on both programmable and non-programmable pocket calculators. Results are presented of extensive statistical testing of two specific generators.

KEY WORDS \& PHRASES: Random number generators, pocket calculators

[^1]
## 1. INTRODUCTION

It has recently become practicable to carry out complex statistical sampling procedures and large simulation experiments using a (programmable) pocket calculator. This calls for pseudo random number generators which are quick, simple and reliable, especially in the sense that a long stream of numbers can be safely produced for a single sample or simulation. In a recent consultation project we came across this need and discovered that, on the one hand, the ready made generators in the computing literature are generally aimed at large computers and binary arithmetic (pocket calculators use decimal arithmetic), while, on the other hand, the random number generators in manufacturers' statistical packages are sometimes of surprisingly bad quality. In fact one of the recommended generators would produce a stream of numbers in which a very short subsequence continually repeated itself if the seed (the first number in the stream) had been chosen unfortunately.

No new theory is presented in this paper. We simply use the standard theory, to be found in KNUTH (1981), to motivate the choice of two specific generators from the well-known class of linear congruential generators. We present the results of extensive statistical testing of the two generators. The simpler one, which produces 5-figure random numbers, turns out not to be suitable for widescale use; however the more complicated 9-figure generator seems to be quite acceptable. We hope that this generator will be used in practice, or better still, that our paper will stimulate the development of better generators, which will also be reported on.

Methods of producing random numbers from distributions other than the uniform are described by VAN PUTTEN \& VAN DER TWEEL (1979). Two important bibliographies are given in SOWEY (1972) and SAHAI (1980).
2. LINEAR CONGRUENTIAL GENERATORS WITH MODULUS A POWER OF TEN

Linear congruential generators are defined by the formula

$$
\begin{equation*}
x_{n+1}=\left(a x_{n}+c\right) \bmod m \tag{1}
\end{equation*}
$$

where $a, c$ and $m$ are positive integers, $a$ and $c$ less than $m$. If a seed
$x_{0}$ is chosen from the set $\{0,1, \ldots, m-1\}$ then the formula produces a sequence $x_{1}, x_{2}, \ldots$ of numbers in the same set, called pseudo random numbers. We shall only consider generators of the form (1) with $m=10^{k}$ for some positive integer $k$. If we define $u_{n}=x_{n} / m, b=c / m$, then an equivalent generator is

$$
\begin{equation*}
u_{n+1}=\left(a u_{n}+b\right) \bmod 1=\left\{a u_{n}+b\right\} \tag{2}
\end{equation*}
$$

where $\{\cdot\}$ denotes the fractional part of a number. We shall also use the notation [•] to denote the entier or whole part of a number. In (2), $b$ and the stream of numbers $u_{0}, u_{1}, \ldots$ are $k$-figure decimal fractions; $a$ is an at most $k$-figure integer. The form (2) is very suitable for a pocket calculator since many calculators include keys for the operations \{•\} and $[\cdot]$; and ctherwise the operations are easy to carry out mentally. Suppose a calculator has an $r$-figure display and r-figure accuracy, so that multiplication or addition of two positive integers whose product or sum is smaller than $10^{r}$ can be carried out exactly (i.e. without rounding errors). Then (2) is particularly easy to implement with $k=[r / 2]$. With a l0-figure display this means that 5-figure random numbers can be easily produced and a programable calculator is not needed at all. However the specific generator of this type which we describe later turns out only to be suitable for very rough-andready applications. If more accuracy or a large stream of numbers is needed, or if a programmable calculator is available, we suggest taking $k=r-1$ and carrying out (2) without rounding errors by writing $u_{n}$ as the sum of two numbers containing its $\ell=[r / 2]$ most significant and $k-\ell$ least significant digits, and $a$ as the sum of two numbers containing its $\ell$ least significant and $k-\ell$ most significant digits. The product in (2) becomes a sum of 4 products, one of which is an integer and so may be discarded immediately, while the others can be calculated exactly and their fractional parts taken separately.

To state this more concisely, our recommendation is to rewrite (2) as

$$
\begin{equation*}
u_{n+1}=\left\{\left\{a^{(1)} u_{n}^{(2)}\right\}+\left\{a^{(2)} u_{n}^{(1)}\right\}+a^{(2)} u_{n}^{(2)}+b\right\} \tag{3}
\end{equation*}
$$

where
congruential generators (1), the so called spectral test. This test is based on the fact that the $m$ points $\left(x_{i}, x_{i+1} \ldots, x_{i+t-1}\right), i=0, \ldots, m-1$, 1ie in the hypercube $[0, m]^{t}$ on various sets of parallel equidistant hyperplanes (according to Wallace Givens, the random numbers "stay mainly in the planes"). The random numbers are of good quality if the distance between consecutive hyperplanes,for the set of hyperplanes which makes this distance maximal, is as small as possible. The spectral test simply consists in computing these maximal distances $v_{t}^{-1}$ for a number of values of $t$; the transformation $\mu_{t}=\pi^{t / 2} \nu_{t}^{t} /((t / 2)!m)$ yields an index of quality which is roughly comparable over different values of $t$ and $m$ (see KNUTH (1981)).

We carried out the spectral test for our two generators (see VAN ES \& VAN PUTTEN (1979) for a discussion of the difficulties involved in this computation) obtaining the results in Table 1.

|  | $\mu_{2}$ | $\mu_{3}$ | $\mu_{4}$ | $\mu_{5}$ | $\mu_{6}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| gen. I | 0.11 | 1.52 | 0.91 | 1.24 | 0.21 |
| gen. II | 0.81 | 2.15 | 0.56 | 2.21 | 3.43 |

Table 1: Results of the spectral test

Both generators "pass the spectral test" according to Knuth's rule of thumb: all of $\mu_{2}$ to $\mu_{6}$ are larger than 0.1 . Gen. II is quite close to "passing the spectral test with flying colours", for which $\mu_{2}$ to $\mu_{6}$ must all exceed the value 1.0 .

## 4. TESTING THE GENERATORS: STATISTICAL TESTS

We finally subjected streams of numbers from generators $I$ and II to a number of standard statistical tests of (various aspects of the null hypothesis: " $x_{1}, x_{2}, \ldots, x_{N t}$ are realizations of $N t$ independent random variables, each uniformly distributed over $\{0,1, \ldots, m-1\}$ ". The following tests, described more fully in KNUTH (1981), were used ${ }^{(*)}$. The notation is also his. Unless otherwise stated $t=1$.

[^2]| (i) | Kolmogorov-Smirnov test | d.f. | $\begin{gathered} N \\ 100 \end{gathered}$ |
| :---: | :---: | :---: | :---: |
| (ii) | Frequency test ( $d=51, t=2$ ) | 50 | 1000 |
| (iii) | Serial test ( $d=10$ ) | 99 | 1000 |
| (iv) | Gap test $\quad\left(\alpha=0, \beta=\frac{1}{2}, t=7\right)$ | 7 | 1000 |
| (v) | Gap test $\quad\left(\alpha=\frac{1}{4}, \beta=\frac{3}{4}, t=7\right)$ | 7 | 1000 |
| (vi) | Gap test $\quad\left(\alpha=\frac{1}{2}, \beta=1, t=7\right)$ | 7 | 1000 |
| (vii) | Partition test ( $d=5, t=4$ ) | 3 | 1000 |
| (viii) | Coupon Collector's test ( $d=5, t=10$ ) | 5 | 500 |
| (ix) | Permutation test ( $t=4$ ) | 23 | 1000 |
| (x) | Run test (up) | 6 | 5000 |
| (xi) | Run test (down) | 6 | 5000 |

A11 but the first test (Kolmogorov-Smirnov) yield statistics which, under the null hypothesis, are approximately $\chi^{2}$-distributed for large $N$. The column "d.f." contains the corresponding number of degrees of freedom. Each test was carried out 40 times for each generator, except that we did not carry out the run tests for gen. I since these tests use more than the whole cycle for this generator ${ }^{(*)}$. For gen. II, except for the KolmogorovSmirnov test, all tests were carried out on disjunct subsequences of length $N t$ of the sequence $\left\{x_{0}, x_{1}, \ldots, x_{m-1}\right\}$. These test results are therefore, under the null hypothesis, independent, and may be combined in various ways.

We have combined the results for gen. II in the following ways:
(i) For each test we have counted the number of significant values at the $\alpha=.05$ leve1.
(ii) Apart from test (i), for each test we have computed the sum of the 40 realizations which, under the null hypothesis, is also a realization of an approximately $\chi^{2}$-distributed random variable.
(iii) We have computed the Fisher combination of the 40 p-values (right tail probabilities) for each test (comparing $-2 \Sigma_{i=1}^{40} \log p_{i}$ with the $x^{2}$-distribution with 80 degrees of freedom).
(iv) A Fisher combined test has also been carried out for all $10 \times 40$ pvalues for tests (ii) to (xi).
${ }^{(*)}$ Admittedly this also applies to tests (iv) to (ix).
(v) We have added the values of all $10 \times 40$ test statistics for tests (ii) to (xi) obtaining another approximately $\chi^{2}$-distributed test statistic. These results are summarized in Table 2. In the table we also give the number of significant values, again at the $5 \%$ level, for gen. I. An asterisk draws attention to a p-value less than $5 \%$

|  | gen. II |  |  |  |  | gen. I |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | sum | p-value | Fisher comb. | p-value | \# sign. | \# sign. |
| (i) K.-S. | - | - | - | - | 1 | 6 |
| (ii) freq. | 1952 | . 775 | 73 | . 711 | 1 | 17 |
| (iii) serial | 4018 | . 259 | 95 | . 127 | 4 | 4 |
| (iv) gap (0, ${ }^{2}$ ) | 251 | . 893 | 65 | . 886 | 0 | 1 |
| (v) gap $\left(\frac{1}{4}, \frac{3}{4}\right)$ | 267 | . 671 | 75 | . 635 | 3 | 1 |
| (vi) gap ( $\left.\frac{1}{2}, 1\right)$ | 258 | . 826 | 66 | . 870 | 1 | 1 |
| (vii) part. | 112 | . 689 | 73 | . 696 | 2 | 0 |
| (viii) coup. coll. | 234 | . 045 * | 100 | . 061 | 2 | 2 |
| (ix) perm. | 899 | . 692 | 74 | . 656 | 4 | 0 |
| (x) run(up) | 209 | . 924 | 63 | . 919 | 2 | - |
| (xi) run(down) | 291 | .010* | 106 | . 027 * | 3 | - |
| $\begin{aligned} & \text { combined } \\ & \text { (excluding (i)) } \end{aligned}$ | 8493 | . 417 | 790 | . 405 | - | - |

## Table 2. Results of statistical tests.

In the Appendix we also present the actual values of the test statistics for those who would like to use other criteria.

We have also tested gen. II by computing Pearson correlations between successive random numbers and between random numbers which are two, three, four or five places apart in the sequence. Among 40 correlations based on 100 pairs of numbers, two were significant at the $5 \%$ level. These results are also included in the appendix.

## 5. CONCLUSIONS

Generator I. This generator is unfortunately of rather poor quality. Both the Kolmogorov-Smirnov test and the frequency test give bad results; these two tests are tests of uniformity over $\{0,1, \ldots, m-1\}$ which is perhaps the most crucial requirement for a pseudo random number generator. Generator II. In general the results for this generator are very satisfactory. Especially the behaviour under Kolmogorov-Smirnov and frequency tests is excellent. The only test which suggests a bad aspect of the generator is the run test (down).
Advice for use. If at all possible use gen. II. In general, when using either generator, we suggest the following rule of thumb : if an application needs $N$ pseudo random numbers, then do not rely on the $l_{\text {ast }} \log _{10} N+2$ digits of the decimal fractions $u_{1}, u_{2}, \ldots, u_{N}$.

The reason for this advice is that the last figures of the numbers from a generator (1) show very clear systematic behaviour (the last figure follows a cycle of length 10 , etc.). The more numbers are required, the more serious this is.

## REFERENCES

VAN ES, A.J. \& C. VAN PUTTEN (1979), The Statal random number generator, report $\mathrm{SN} 8 / 79$, Mathematical Centre, Amsterdam.

KNUTH, D.E. (1981), The art of computer programming, Vol. II, Seminumerical algorithms (2nd edition), Addison-Wesley, Reading, Massachusetts.

VAN PUTTEN, C. \& I. VAN DER TWEEL (1979), On generating random variables, report SN 9/79, Mathematical Centre, Amsterdam.

SAHAI, H. (1980), A supplement to Sowey's bibliography on random number generation and related topics, Biom. J. 22 447-461.

SOWEY, E.R. (1972), A chronological and classified bibliography on random number generation and testing, Int. Stat. Rev. 40 355-371.

## APPENDIX

Table 3. Results of Kolmogorov-Smirnov test (* denotes significance at $5 \%$ 1eve1)


Table 4.
Results of statistical tests (ii)-(ix) for generator $I$ (* denotes significance at $5 \%$ level; note subcycles in columns (iii) and (vii))

|  | (ii) | (iii) | (iv) |  | (v) | (vi) | (vii) | (viii) |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | (ix)

Table 5.

Results of statistical tests (ii)-(xi) for generator II. (* denotes significance at 5\% leve1)

| (ii) | (iii) | (iv) | (v) | (vi) | (vii) | viii) | (ix) | (x) | (xi) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 57.4 | 100.6 | 9.56 | 6.62 | 4.20 | 2.03 | 8.24 | 32.8 | 3.92 | 5.47 |
| 46.3 | 111.4 | 10.86 | 0.82 | * 18.27 | 5.55 | 9.53 | 27.7 | 2.43 | 4.50 |
| 39.3 | 99.4 | 3.31 | 7.76 | 10.95 | 1.20 | 6.39 | 28.6 | 5.32 | 7.33 |
| 52.5 | 91.2 | 3.52 | 7.96 | 7.58 | 2.87 | 2.55 | 15.8 | 1.87 | 9.26 |
| 58.1 | 77.8 | 7.75 | 3.52 | 4.52 | 1.86 | 7.66 | 16.5 | * 16.35 | *15.24 |
| 48.7 | 83.4 | 4.09 | 10.53 | 0.99 | 3.53 | 4.32 | 18.3 | 6.18 | 6.98 |
| 47.7 | *125.8 | 10.00 | 6.81 | 5.72 | 0.93 | 5.89 | 26.4 | 8.33 | 8.99 |
| 62.2 | 76.0 | 8.91 | 4.69 | 7.92 | 6.85 | 7.98 | 18.2 | *1.05 | *17.19 |
| 45.1 | 86.4 | 8.01 | 1.80 | 8.37 | 0.69 | 4.26 | *38.9 | 1.40 | 5.59 |
| 40.2 | 81.0 | 10.60 | 3.57 | 5.76 | 3.22 | 2.14 | 15.5 | 7.22 | 5.69 |
| 47.3 | 83.8 | 12.99 | 5.10 | 4.13 | 3.34 | 3.76 | 22.1 | 4.76 | 6.10 |
| 48.7 | 95.6 | 6.38 | 7.14 | 7.57 | *8.88 | 6.22 | 27.8 | 9.17 | 6.21 |
| 44.4 | *136.2 | 0.56 | 7.10 | 7.43 | 1.93 | 9.34 | 11.7 | 2.03 | 8.72 |
| 43.3 | 96.2 | 7.07 | 6.87 | 1.86 | 1.29 | 7.39 | 21.5 | 4.08 | 3.32 |
| 42.4 | 97.4 | 10.49 | 7.68 | 4.15 | 0.79 | 7.00 | 19.2 | 3.30 | 10.16 |
| 52.7 | 92.4 | 5.57 | * 14.37 | 7.75 | 3.84 | 4.84 | 19.5 | 2.44 | 6.42 |
| 43.2 | 84.0 | 5.00 | 5.60 | 5.31 | 2.14 | 9.55 | 26.0 | 7.01 | 6.57 |
| 34.5 | 90.4 | 11.87 | 6.37 | 3.84 | 1.37 | 3.91 | 23.1 | 3.23 | 5.77 |
| 46.4 | 87.6 | 5.87 | 7.62 | 7.71 | 0.76 | 4.09 | 18.7 | 2.77 | 8.35 |
| 51.7 | 107.0 | 0.53 | 6.65 | 3.94 | 2.16 | 3.59 | 15.2 | 7.94 | 11.27 |
| *85.0 | 122.6 | 6.19 | 8.49 | 4.62 | 4.57 | 2.94 | 19.3 | 2.66 | 1.9 |
| 38.8 | 96.2 | 5.96 | 4.07 | 7.01 | 4.11 | 2.19 | 20.6 | 6.70 | 9.68 |
| 56.4 | *143.8 | 5.24 | 4.73 | 4.65 | 0.37 | 4.04 | 13.2 | 7.36 | 3.6 |
| 44.6 | 112.4 | 6.32 | 3.41 | 4.69 | 0.51 | * 12.04 | 22.3 | 3.16 | 2.4 |
| 23.7 | 88.4 | 3.23 | 9.50 | 6.52 | 5.18 | 9.66 | 18.4 | 4.42 | 5.7 |
| 38.6 | 110.6 | 8.85 | 2.42 | 4.86 | 1.22 | 2.56 | 22.3 | 4.34 | 3.3 |
| 50.3 | 98.8 | 1.75 | *21.56 | 5.72 | 2.44 | 7.50 | 19.8 | 8.63 | 4.2 |
| 52.2 | 114.8 | 4.31 | 13.22 | 9.37 | 5.27 | 1.80 | *35.9 | 1.40 | 11. |
| 27.6 | 120.6 | 1.35 | 4.04 | 8.81 | 1.42 | 1.07 | 16.1 | 7.00 | *15.73 |
| 46.1 | 84.2 | 1.94 | 4.84 | 11.29 | 1.01 | 1.20 | 13.9 | 4.87 | 10.3 |
| 41.9 | 98.4 | 9.80 | 4.75 | 10.86 | 1.80 | 3.99 | *36.4 | 5.62 | 3. |
| 58.4 | 116.0 | 4.76 | 5.4 | 5.12 | 3.02 | 2.94 | 27.0 | 8.39 | 5. |
| 56.0 | 110.4 | 5.61 | 5.34 | 4.14 | 3.00 | 3.36 | 26.8 | 5.85 | 6. |
| 49.5 | 83.6 | 9.44 | 10.12 | 3.80 | 1.82 | 4.86 | *37.2 | 5.29 | 6. |
| 53.6 | 101.4 | 9.97 | 5.17 | 5.00 | 0.91 | 5.87 | 16.8 | 12.72 | 6.5 |
| 64.1 | 86.8 | 5.23 | 2.08 | 4.70 | 4.91 | *20.04 | 21.6 | 3.07 | 3. |
| 53.5 | 93.0 | 3.66 | 5.92 | 10.65 | 1.32 | 5.36 | 23.7 | 3.86 | 12.28 |
| 46.8 | 104.2 | 4.67 | 5.20 | 8.64 | *8.38 | 6.68 | 24.6 | 6.56 | 6. |
| 62.4 | * 130.0 | 4.13 | * 14.35 | 4.28 | 4.66 | 9.92 | 18.2 | 2.58 | 6. |
| 50.8 | 98.0 | 5.25 | 6.35 | 4.90 | 0.95 | 7.26 | 21.0 | 3.36 | 6.36 |

The following tot Pearson correlation matrices were computed from disjunct subsequences of size $100 t$ from generator II ; i.e. we computed the correlation matrix for the observations $\left(u_{t i+1}, u_{t i+2}, \ldots, u_{t i+t}\right), i=0, \ldots, 99$. The critical value at the $5 \%$ level for a Pearson correlation between two independent uniform [0,1] variables with 100 observations is $\pm 0.195$. For $t=2,3,4$ and 5 we computed two correlation matrices.
$t=2$
Table 6. Pearson correlations.

$t=5$



[^0]:    1980 Mathematics subject classification: Primary 65C10;
    Secondary 62-04, 65U05

[^1]:    *) This report will be submitted for publication elsewhere.
    **) N.V. Philips gloeilampenfabrieken, Centre for quantitative methods, Eindhoven; formerly at Mathematical Centre, Amsterdam

[^2]:     uses subsequences of fixed length.

